

Developing Reactive Systems using Statecharts

Modelling of Software-Intensive Systems

Simon Van Mierlo
University of Antwerp
Belgium
simon.vanmierlo@uantwerpen.be

Hans Vangheluwe
University of Antwerp
Belgium
hans.vangheluwe@uantwerpen.be

Axel Terfloth
itemis AG
Germany
terfloth@itemis.de

Introduction

Reactive Systems



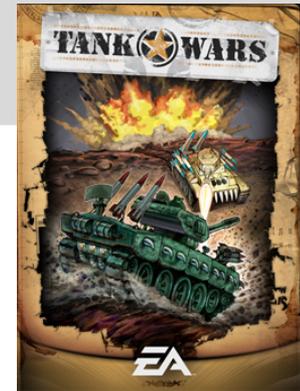
- Com
- In contrast to *transformational* systems, which take input and, eventually, produce output

Modelling Reactive Systems

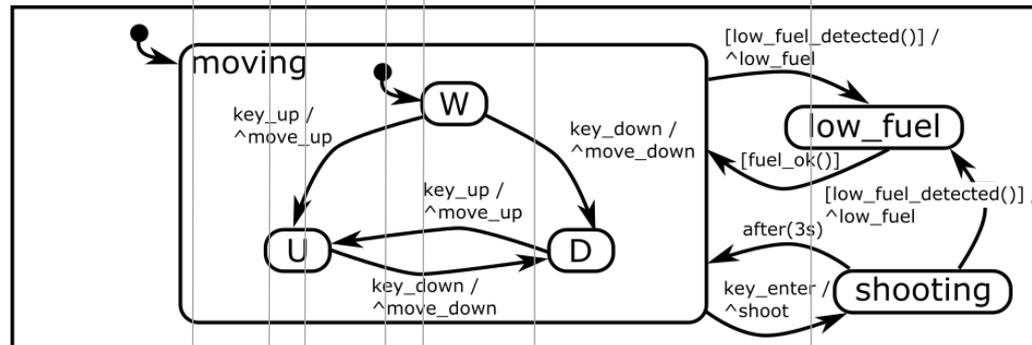
- Interaction with the environment: reactive to *events*
 - Autonomous behaviour: *timeouts* + *spontaneous* transitions
 - System behaviour: *modes* (hierarchical) + *concurrent* units
 - Use programming language + threads and mutexes (OS)
- Programming language (and OS) is too low-level
-> most appropriate formalism: "what" vs. "how"**
- "Minimal software written with threads, semaphores, and mutexes are incomprehensible to humans"*



Discrete-Event Abstraction



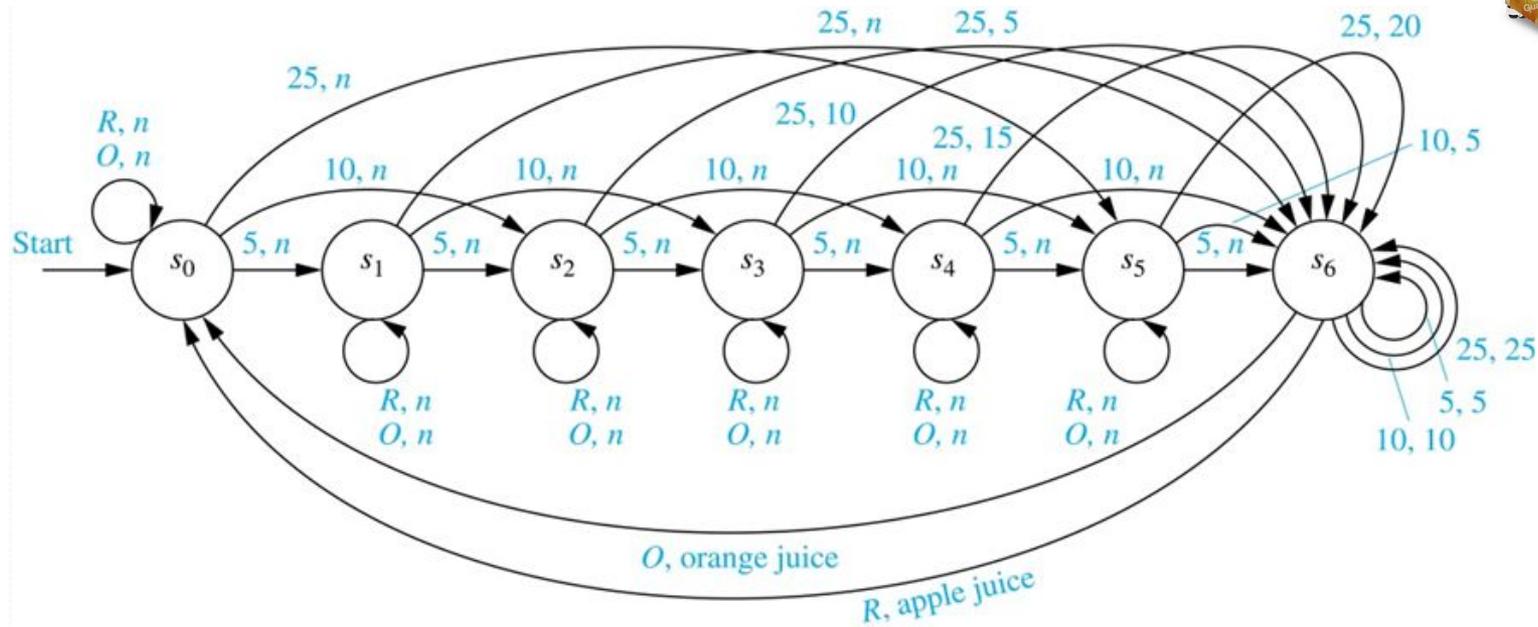
behavioural
model



State Diagrams

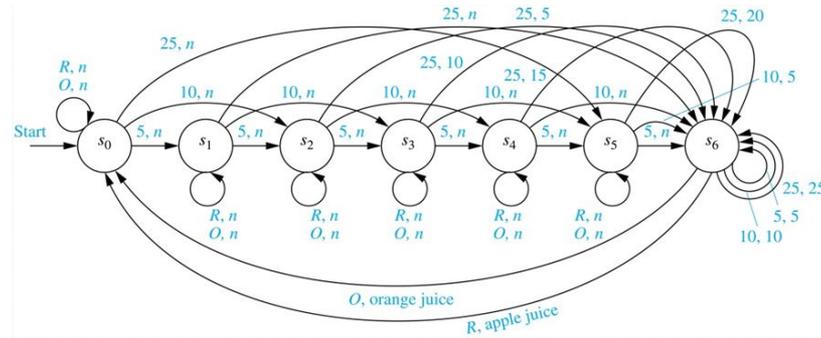


VB1-2



- All states are explicitly represented (unlike Petrinets, for example)
- Flat representation (no hierarchy)
- Does not scale well: becomes too large too quickly to be usable (by humans)

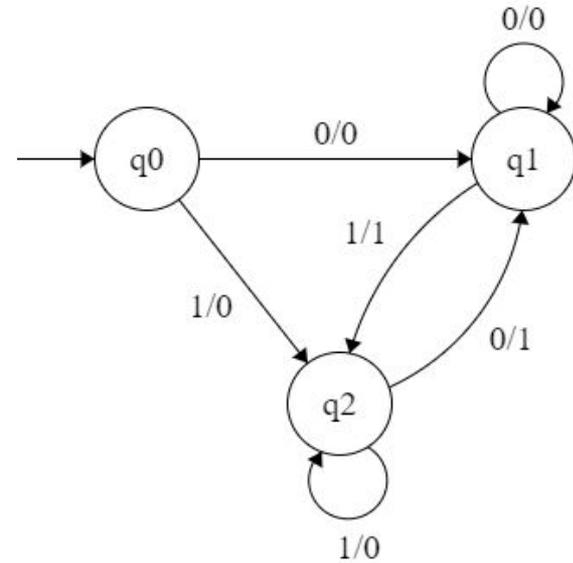
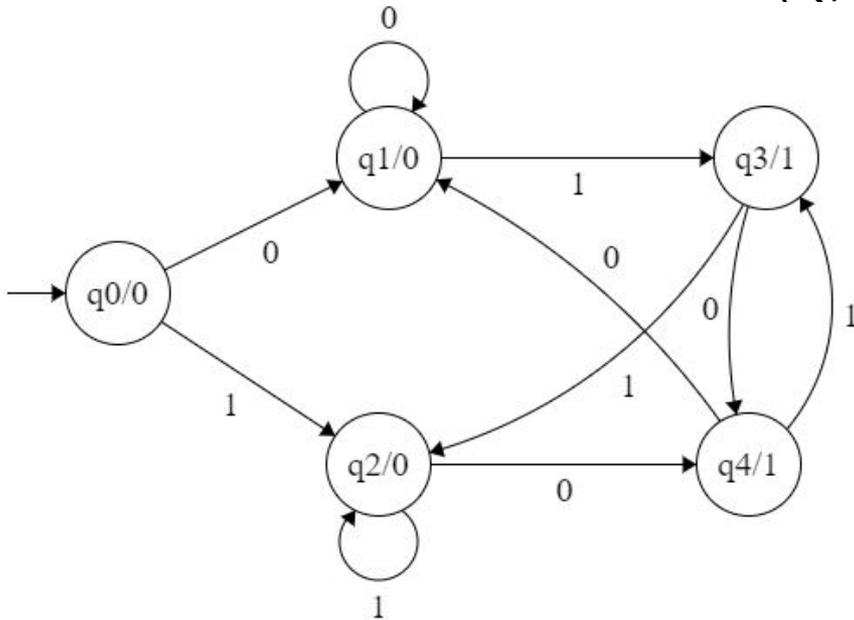
Alternative Representation: Parnas Tables



event \ state	s_0	s_1	s_2	s_3	s_4	s_5	s_6
5	s_1, n	s_2, n	s_3, n	s_4, n	s_5, n	s_6, n	$s_6, 5$
10	s_2, n	s_3, n	s_4, n	s_5, n	s_6, n	$s_6, 5$	$s_6, 10$
25	s_5, n	s_6, n	$s_6, 5$	$s_6, 10$	$s_6, 15$	$s_6, 20$	$s_6, 25$
O	s_0, n	s_1, n	s_2, n	s_3, n	s_4, n	s_5, n	$s_0, \text{orange juice}$
R	s_0, n	s_1, n	s_2, n	s_3, n	s_4, n	s_5, n	$s_0, \text{apple juice}$

Mealy and Moore Machines

FSA: $(Q, q_0, \Sigma, O, \delta, \omega)$



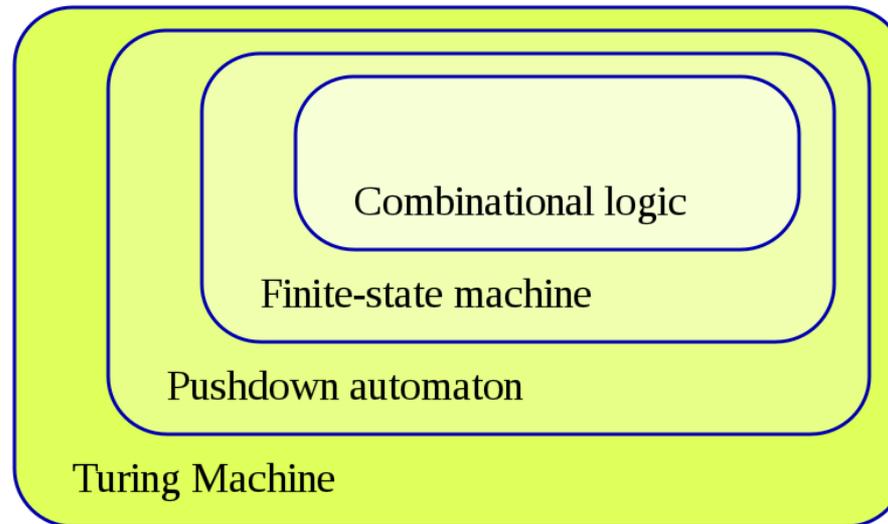
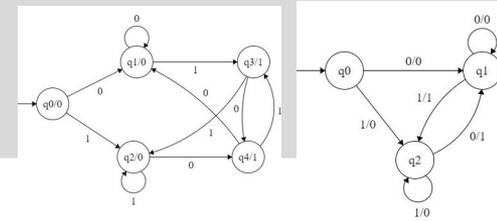
Moore Machines

- Output only depends on current state.
: $Q \rightarrow O$
- Input stream: 10
Output stream: 00

Mealy Machines

- Output depends on current state and on current input.
: $Q \times \Sigma \rightarrow O$
- Input stream: 10
Output stream: 01

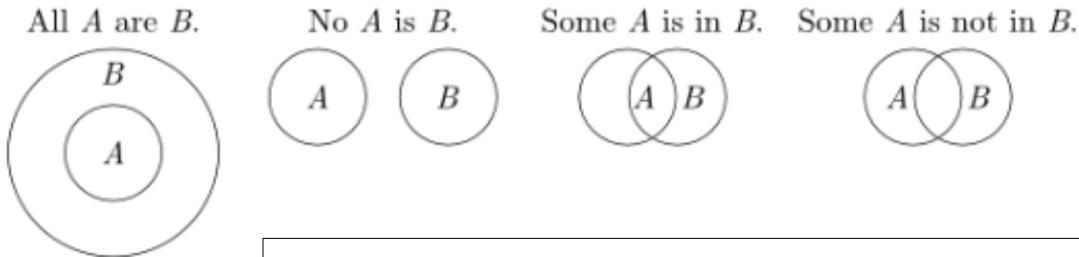
FSAs: Expressiveness



- Can be made Turing-complete
data memory, control flow, branching
- Extend FSAs
borrow semantics from Mealy and Moore machines

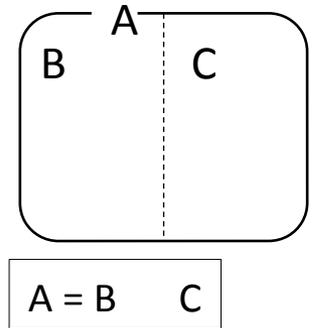
Higraphs

Euler Diagrams

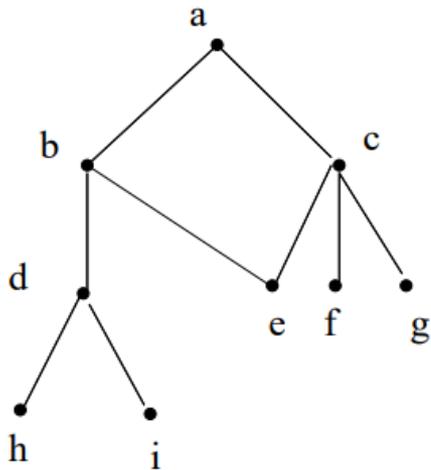


topological notions for set union, difference, intersection

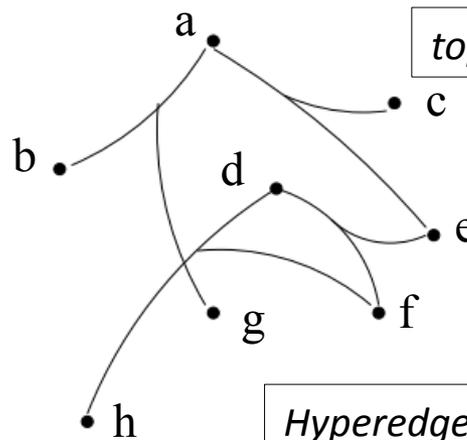
Unordered Cartesian Product



Hypergraphs



a graph



a hypergraph

topological notion (syntax): connectedness

Hyperedges: 2^X (undirected), $2^X \times 2^X$ (directed).

$$X = \{a, b, \dots, h\}$$

Higraphs

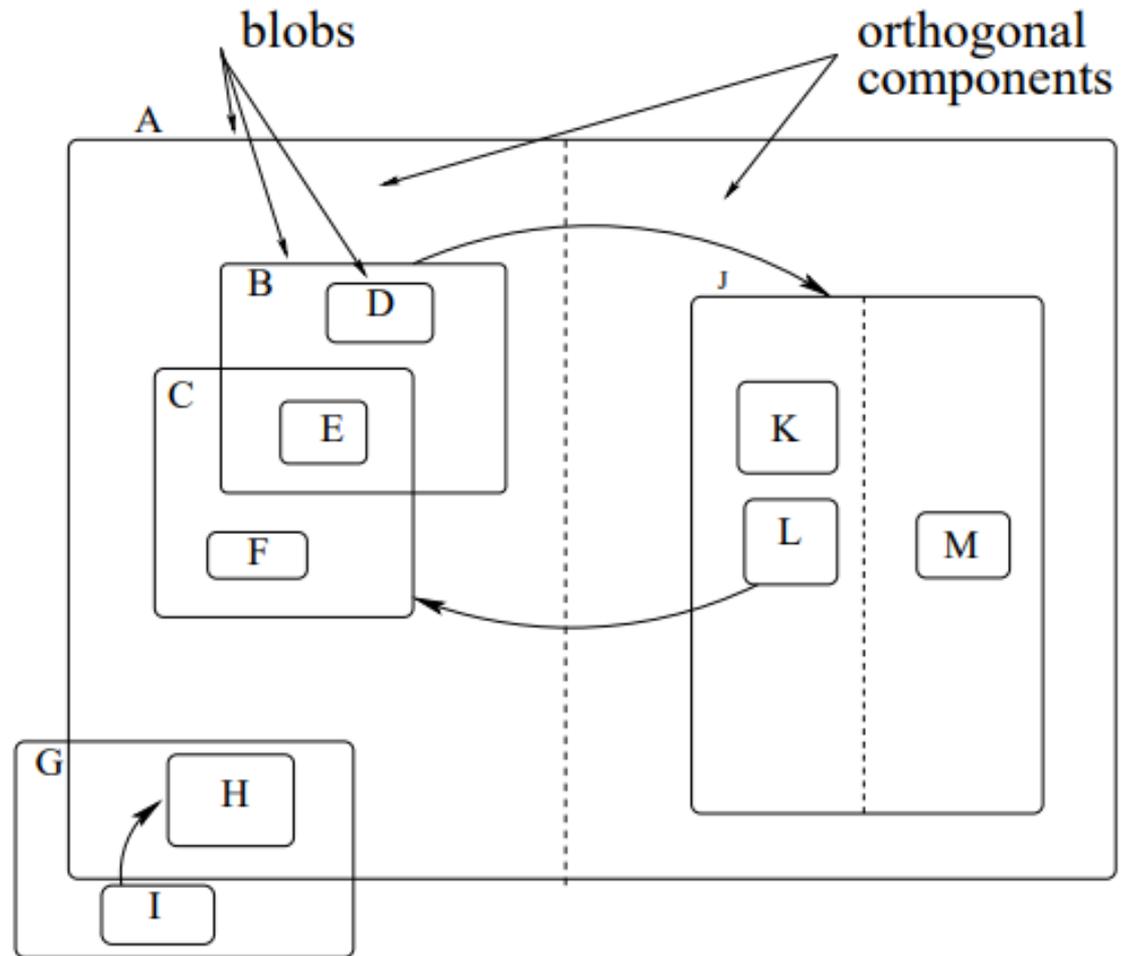
Euler Diagrams

+

Hypergraphs

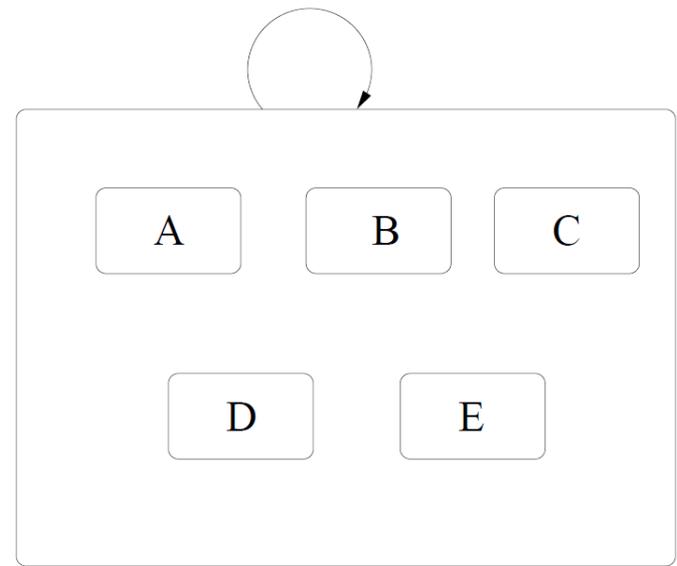
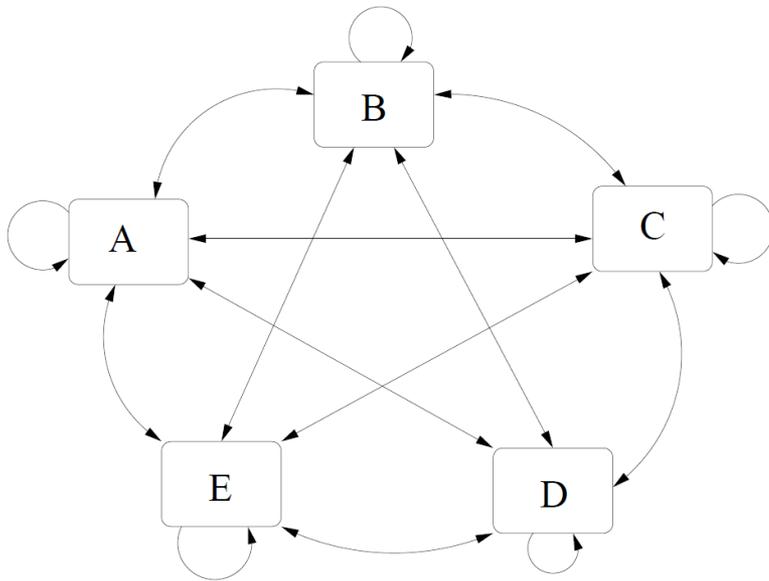
+

Unordered Cartesian Product



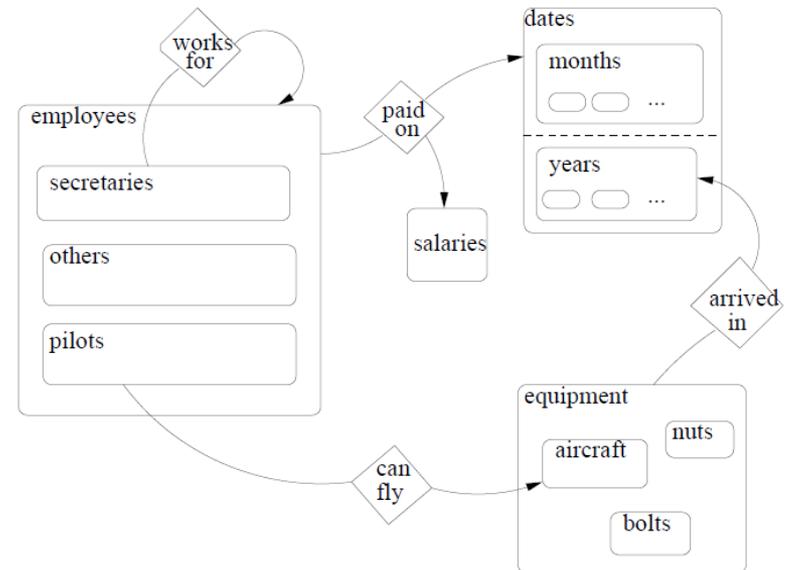
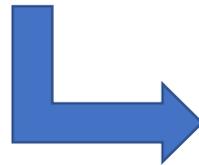
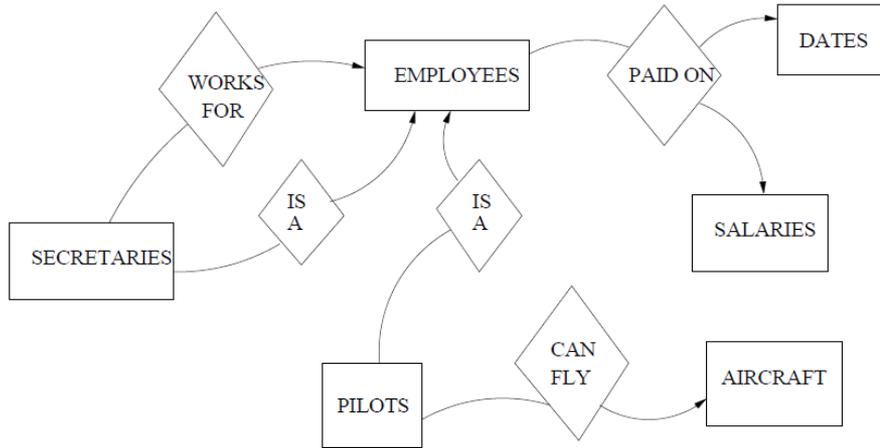
Higraph: Examples

- Clique

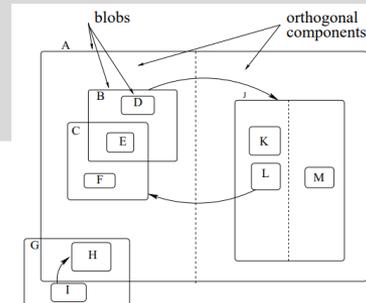


Higraphs: Examples

- ER-Diagrams



Higraphs: Formal Definition



- A higraph H is a quadruple

$$H = (B, E,$$

- B is a finite set of all unique *blobs*

- E is a set of *hyperedges*

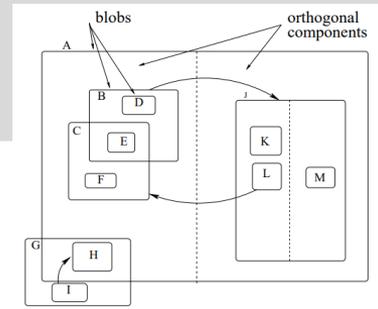
$$2^B \times 2^B$$

- The subblob function

$$: B \rightarrow 2^B$$

$$\sigma^{i+1}(x) = \bigcup_{y \in \sigma^i(x)} \sigma(y) \quad \sigma^+(x) = \bigcup_{i=1} \sigma^i(x)$$

Higraphs: Formal Definition



- Subblobs relation cycle-free

$$x \quad (x)$$

- The partitioning function π associates an *equivalence relationship* with x

$$B \quad 2^{B \times B}$$

- Equivalence classes π_i are orthogonal components of x

$$\pi_1(x), \pi_2(x), \dots, \pi_{k_x}(x)$$

- $k_x = 1$ means a single orthogonal component

- Blobs in different orthogonal components of x are disjoint

$$y, z \quad (x) : \quad +(y) \quad +(z) =$$

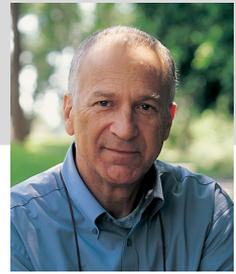
Higraphs Applications

- Apply syntactic constructs to an existing modelling language.
- Add specific meaning to these constructs.
- Examples:
 - E-R diagrams
 - Dataflow/Activity Diagrams
 - Inheritance
 - **Statecharts**

Statecharts

- Visual (topological, not geometric) formalism
- Precisely defined syntax and semantics
- Many uses:
 - Documentation (for human communication)
 - ~~Analysis (of behavioural properties)~~
 - Simulation
 - Code synthesis
 - ... and derived, such as testing, optimization, ...

Statecharts History



- Introduced by David Harel in 1987
- Notation based on higraphs = hypergraphs + Euler diagrams + unordered Cartesian product
- Semantics extends deterministic finite state automata with:
 - Depth (Hierarchy)
 - Orthogonality
 - Broadcast Communication
 - Time
 - History
 - Syntactic sugar, such as enter/exit actions

Statecharts History

- Incorporated in UML: State Machines (1995)
- More recent: xUML for semantics of UML subset (2002)
- W3 Recommendation: State Chart XML (SCXML) (2015)
<https://www.w3.org/TR/scxml/>
- Standard: Precise Semantics for State Machines (2019)
<https://www.omg.org/spec/PSSM/>

Statechart (Variants) Tools

STATEMATE: A Working Environment for the Development of Complex Reactive Systems

Rational software

<https://www.ibm.com/us-en/marketplace/systems-design-rhapsody>

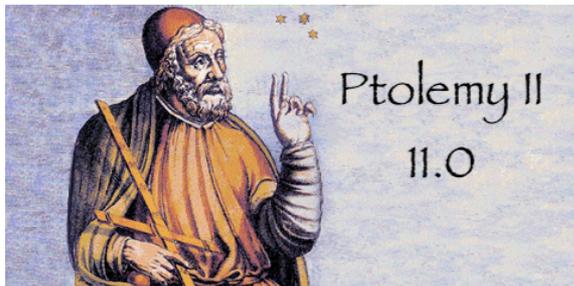


Matlab

Simulink

Stateflow

<https://www.mathworks.com/products/stateflow.html>



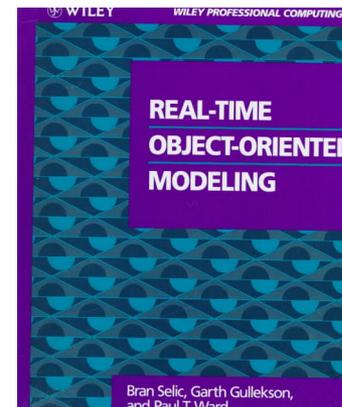
<https://ptolemy.berkeley.edu/ptolemyII/ptII11.0/index.htm>



<https://www.itemis.com/en/yakindu/state-machine/>

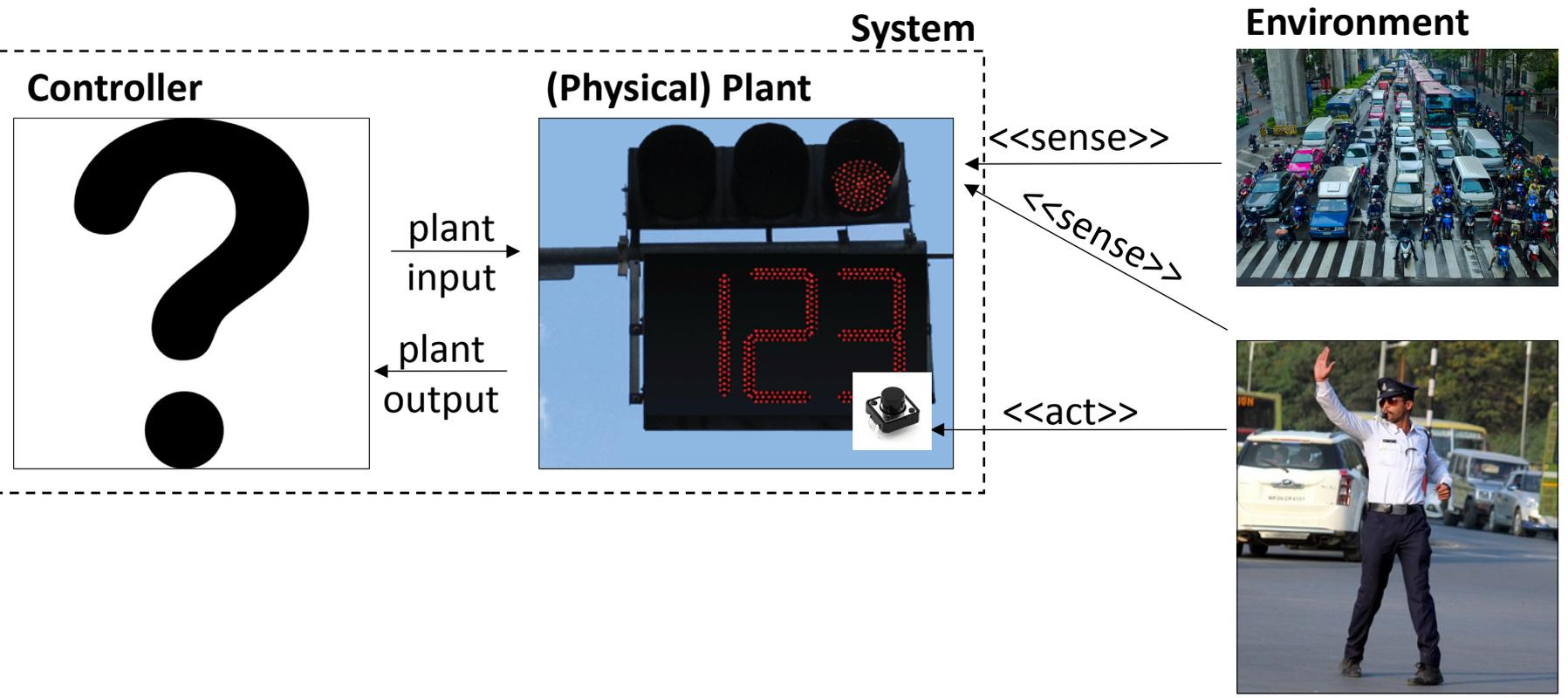


<https://www.eclipse.org/papyrus-rt/>

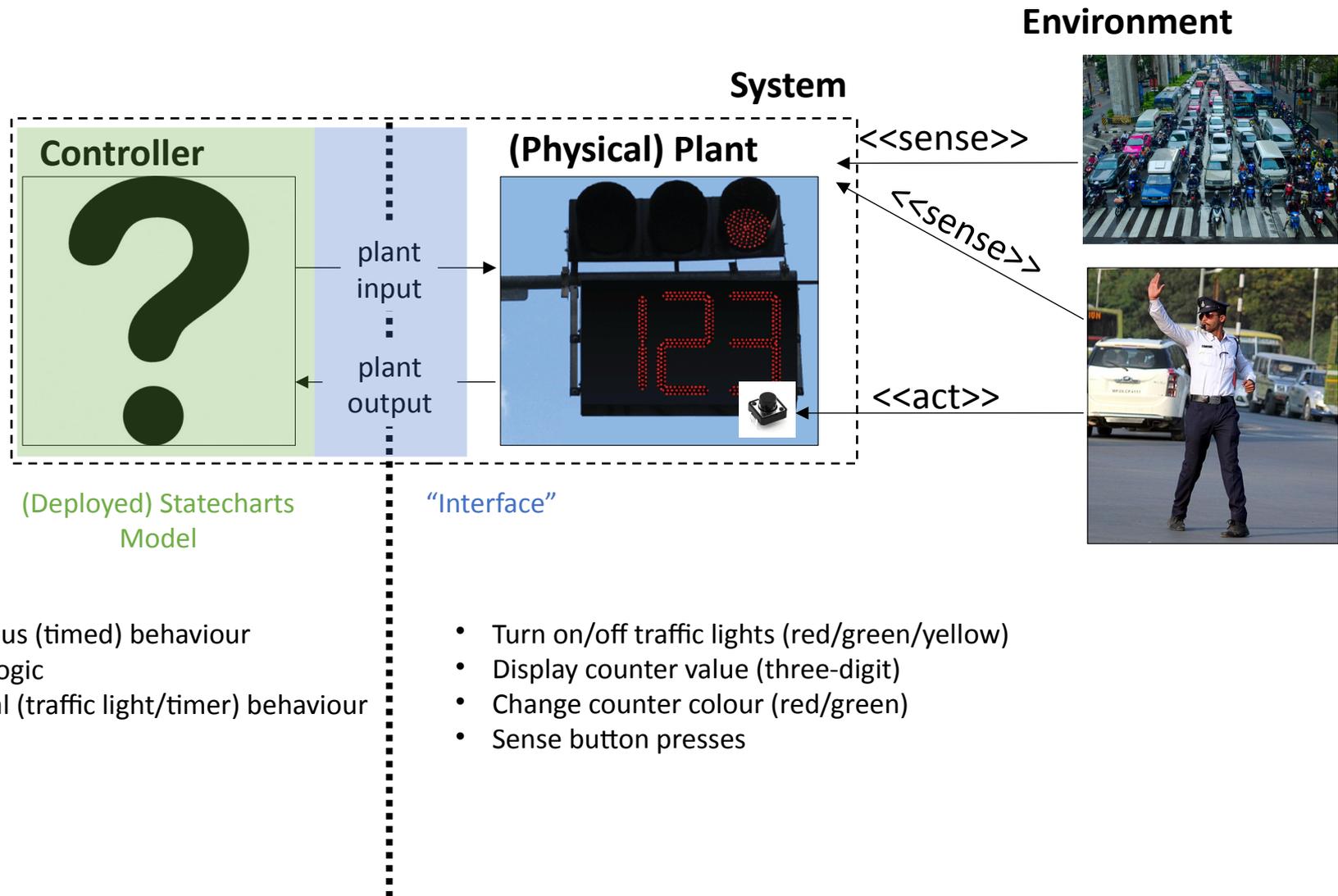


<https://www.eclipse.org/etrice/>

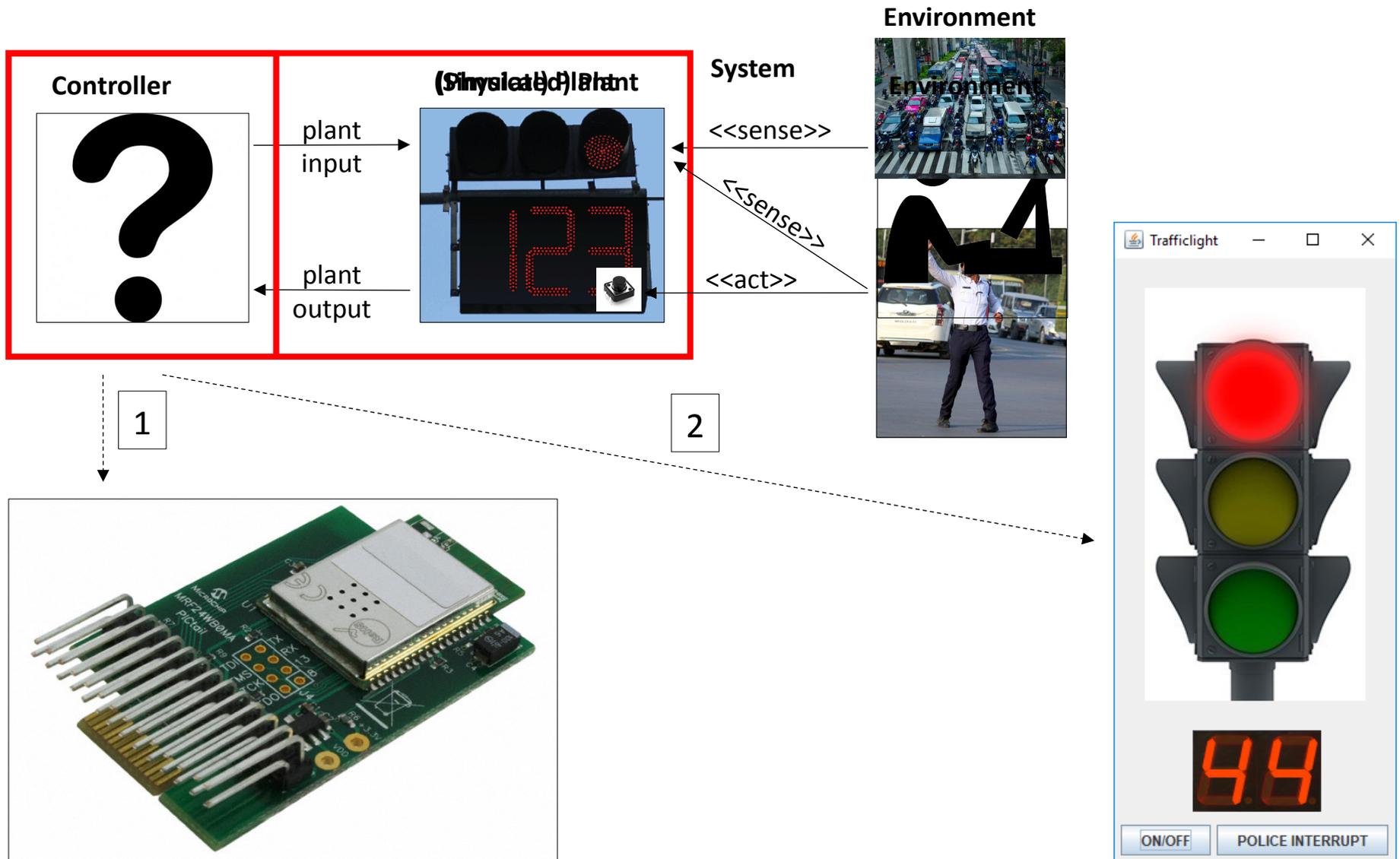
Running Example



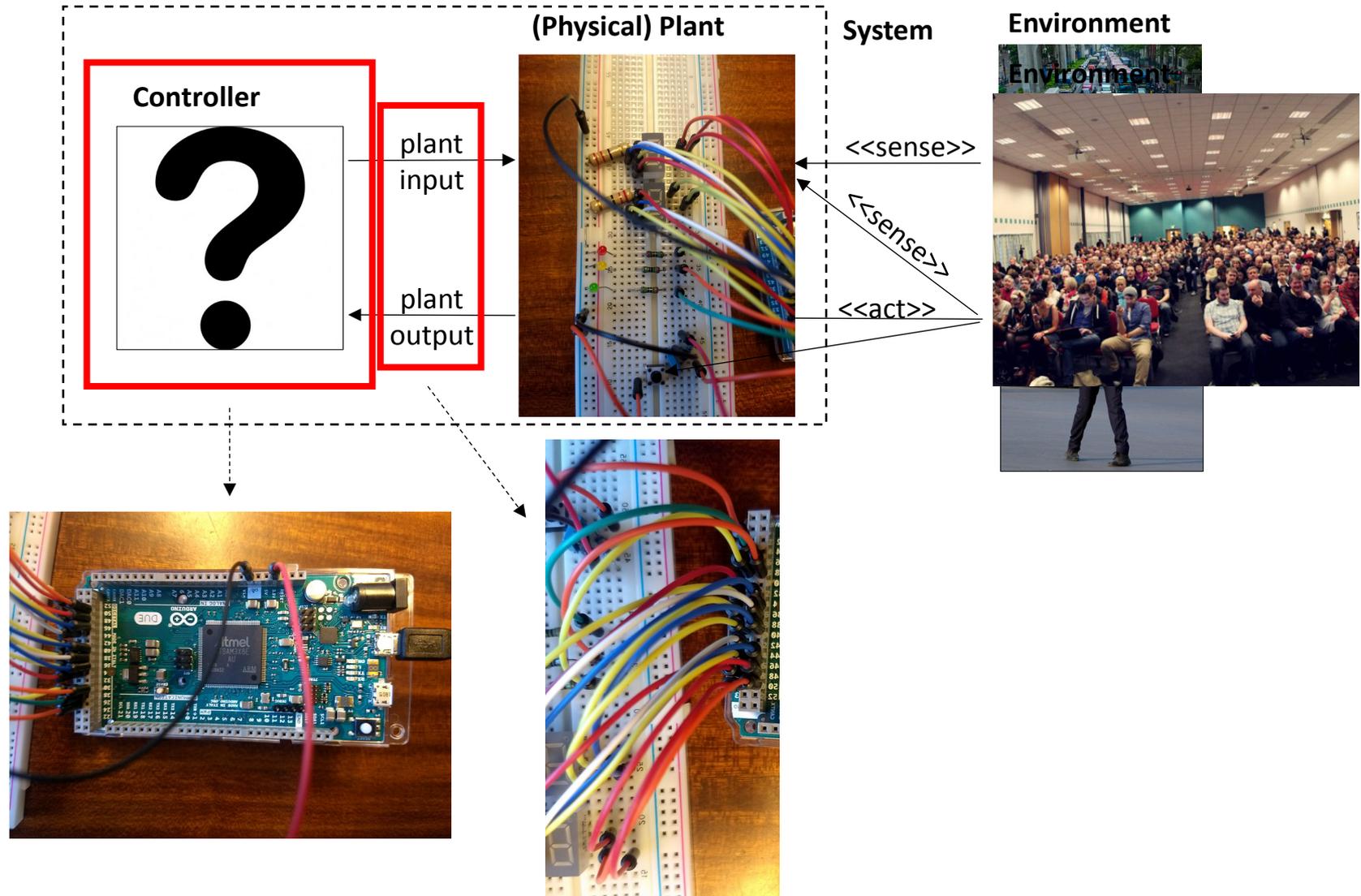
What are we developing?



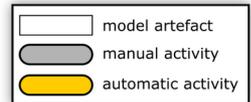
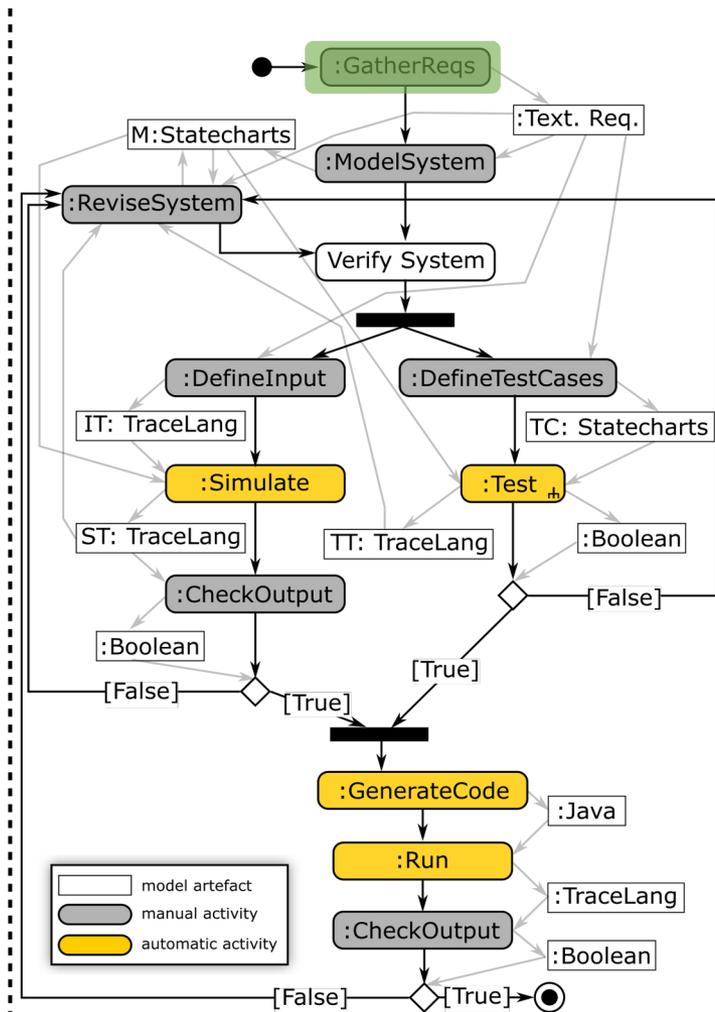
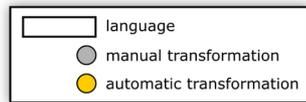
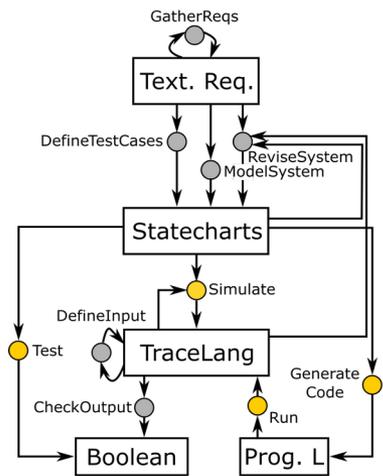
Deployment (Simulation)



Deployment (Hardware)



Workflow

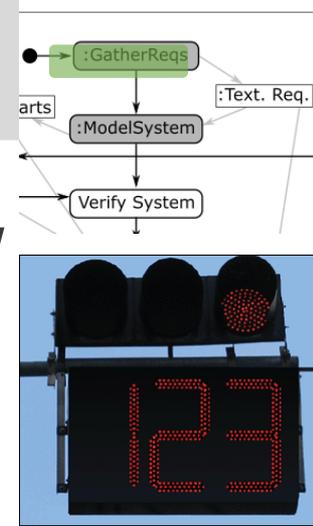


Hans Vangheluwe and Ghislain C. Vansteenkiste. A multi-paradigm modeling and simulation methodology: Formalisms and languages. In European Simulation Symposium (ESS), pages 168-172. Society for Computer Simulation International (SCS), October 1996. Genoa, Italy.

Levi Lúcio, Sadaf Mustafiz, Joachim Denil, Hans Vangheluwe, Maris Jukss. FTG+PM: An Integrated Framework for Investigating Model Transformation Chains. System Design Languages Forum (SDL) 2013, Montreal, Quebec. LNCS Volume 7916, pp 182-202, 2013.

Requirements

- R1: three differently coloured lights: red, green, yellow
- R2: at most one light is on at any point in time
- R3: at system start-up, the red light is on
- R4: cycles through red on, green on, and yellow on
- R5: red is on for 60s, green is on for 55s, yellow is on for 5s
- R6: time periods of different phases are configurable.
- R7: police can interrupt autonomous operation
 - Result = blinking yellow light (on -> 1s, off -> 1s)
- R8: police can resume an interrupted traffic light
 - Result = light which was on at time of interrupt is turned on again
- R9: traffic light can be switched on and off and restores its state
- R10: a timer displays the remaining time while the light is red or green; this timer decreases and displays its value every second. The colour of the timer reflects the colour of the traffic light.





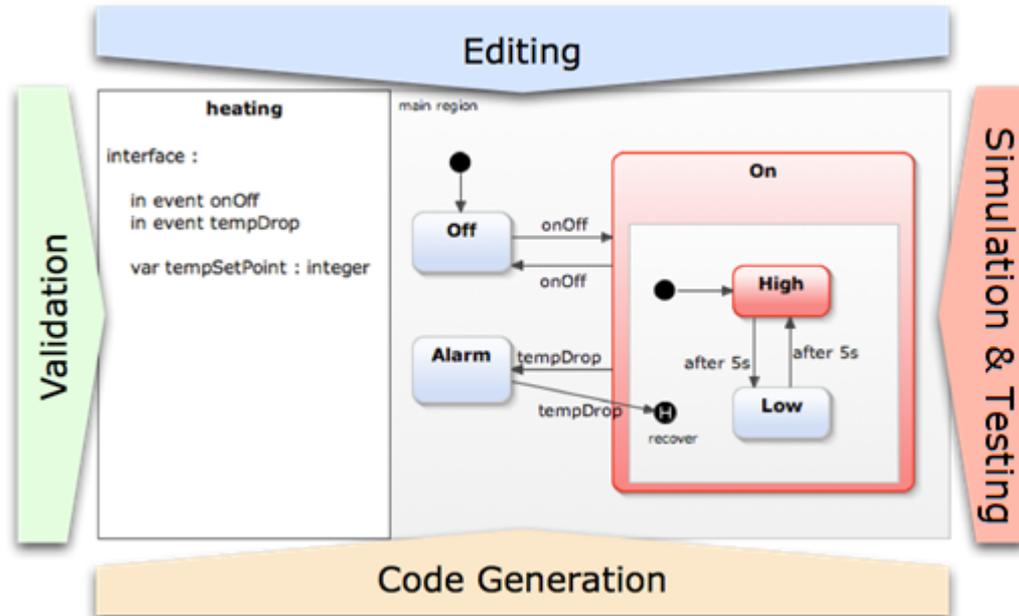
YAKINDU Statechart Tools

Statecharts made easy...



What are YAKINDU Statechart Tools?

YAKINDU Statechart Tools provides an **integrated modeling environment** for the specification and development of **reactive, event-driven systems** based on the concept of statecharts.



The Statecharts Language

States

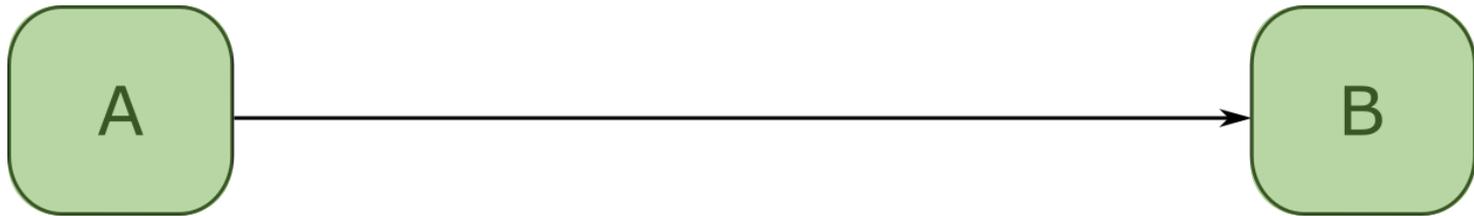


being **in** a state
= state <<*name*>> is
active
= the system is in
configuration <<*name*>>



initial state
exactly one per
model
“entry point”

Transitions



- Model the **dynamics** of the system:
 - *when*
 - the system is **in state A** and
 - the **event** is **received**
 - *then*
 1. **output_action** is evaluated and
 2. the new **active state** becomes B

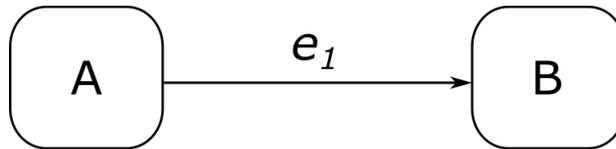
Transitions: Events

event(in_params) / output_action(out_params)

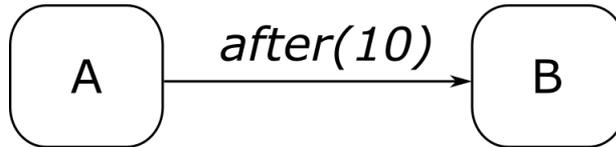
- Spontaneous



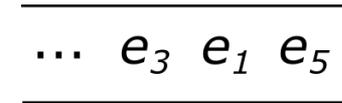
- Input Event



- After Event



queue of event notices



processing

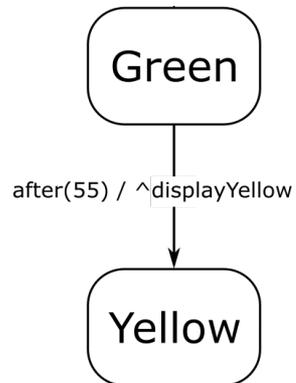
queue of event notices



processing

Transitions: Raising Output Events

`event(in_params) / output_action(out_params)`

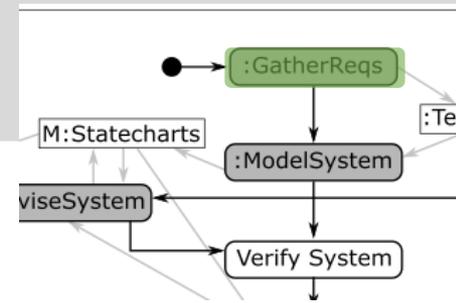


Syntax for output action:

^output_event

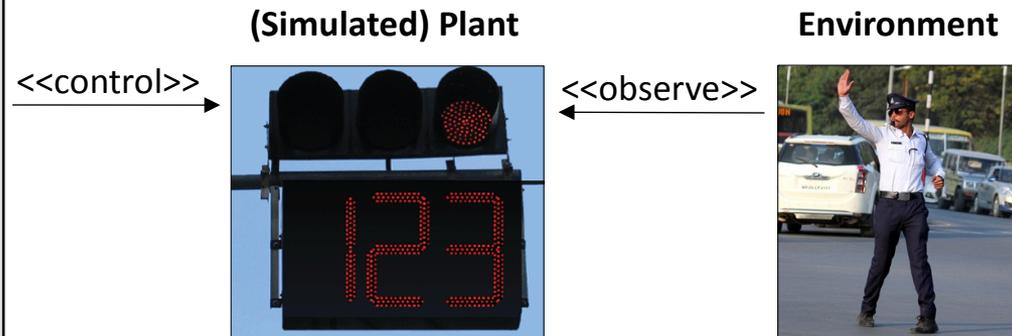
means “raise the event *output_event* (to the environment)”

Exercise 1 - Requirements

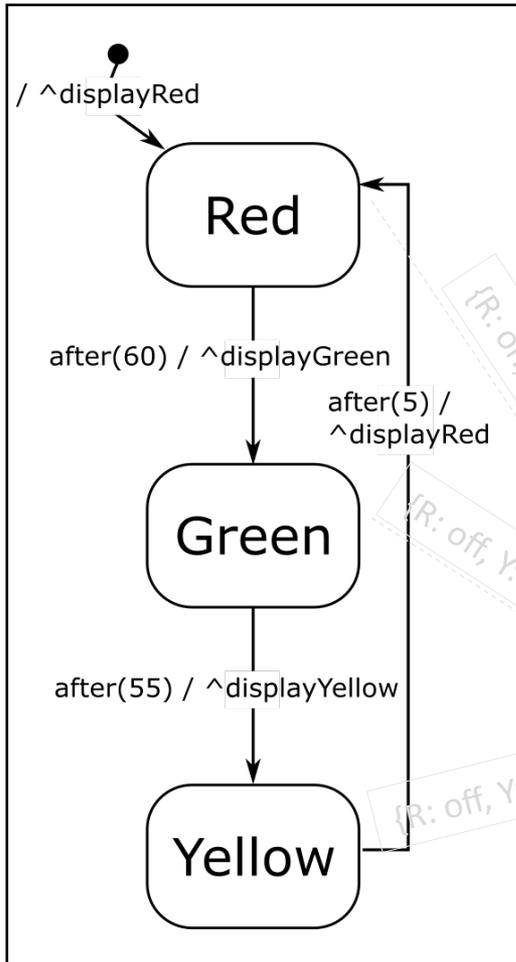
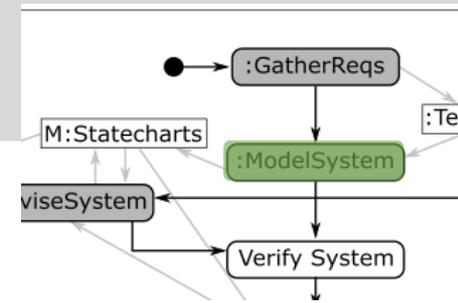


Your model here.

- R1: three differently coloured lights: red (R), green (G), yellow (Y)
- R2: at most one light is on at any point in time
- R3: at system start-up, the red light is on
- R4: cycles through red on, green on, and yellow on
- R5: red is on for 60s, green is on for 55s, yellow is on for 5s



Exercise 1 - Solution



- R1: three differently coloured **lights**: red (R), green (G), yellow (Y)
- R2: at most one light is on at any point in time
- R3: at system start-up, the **red** light is on
- R4: cycles through red on, green on, and yellow on
- R5: **red** is on for 60s, **green** is on for 55s, **yellow** is on for 5s

{R: on, Y: off, G: on}
{R: off, Y: off, G: on}
{R: off, Y: on, G: off}

(Simulated) Plant



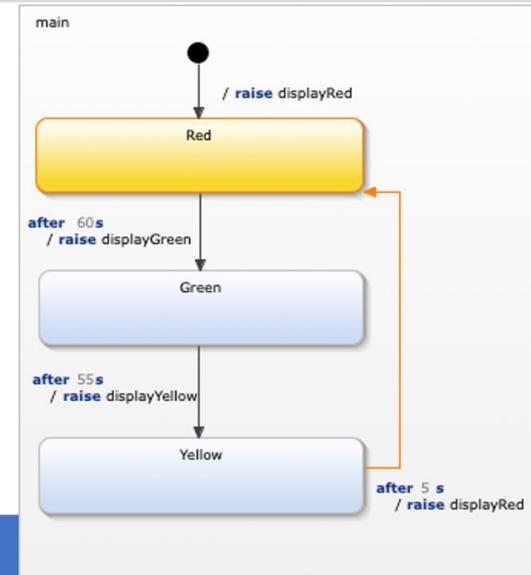
Environment



`<<observe>>`



Exercise 1 - Solution



requirement

modelling approach

R1: three differently coloured **lights**: red (R), green (G), yellow (Y)

For each colour a **state** is defined. Transitions that lead to a state raise the proper out event which interacts with the plant.

R2: at most one light is on at any point in time

The states are all contained in a single region and thus are exclusive to each other ("or" states).

R3: at system start-up, the red light is on

The entry node points to state Red and the entry transition raises the event displayRed.

R4: cycles through red on, green on, and yellow on

The transitions define this cycle.

R5: red is on for 60s, green is on for 55s, yellow is on for 5s

Time events are specified on the transitions.

Data Store

Full System State

<<name>>

+

DataStore
- var ₁ : t ₁ = val ₁
- var ₂ : t ₂ = val ₂
...
- var _n : t _n = val _n

being **in** a state
= state <<*name*>> is
active
= the system is in
configuration <<*name*>>

data store **snapshot**
= variable values

=

full system state

Full System State: Initialization



DataStore
- var ₁ : t ₁ = val ₁
- var ₂ : t ₂ = val ₂
...
- var _n : t _n = val _n

```
1 int main() {  
2  
3 }
```

initial state
exactly one per
model
“entry point”

provide **default
value** for each
variable
“initial snapshot”

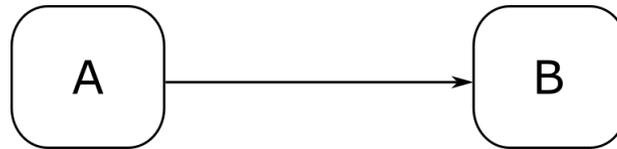
Compare:
C++ initialization
implicit state
(program counter)
+ **data store**

Transitions: Guards

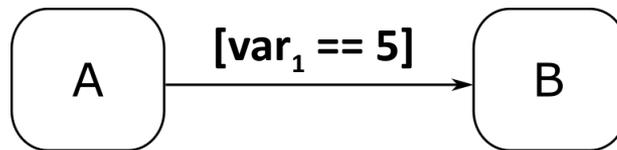
`event(in_params) [guard] / output_action(out_params)`

Modelled by “guard expression” (evaluates to Boolean) in some appropriate language

- Spontaneous [True]

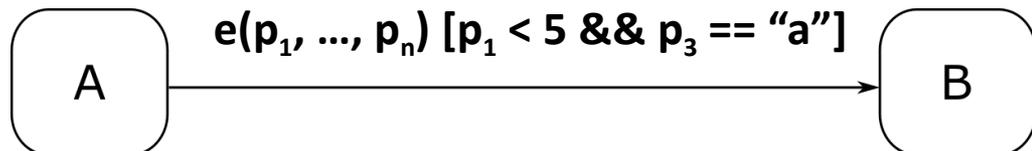


- Data Store Variable Value



DataStore
- var ₁ : t ₁ = val ₁
- var ₂ : t ₂ = val ₂
...
- var _n : t _n = val _n

- Parameter Value

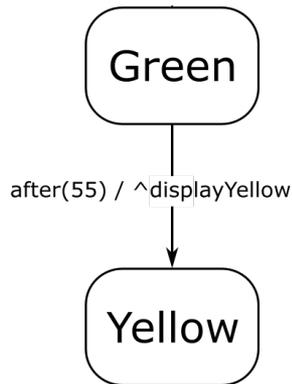


Transitions: Output Actions

`event(params) [guard] / output_action(params)`

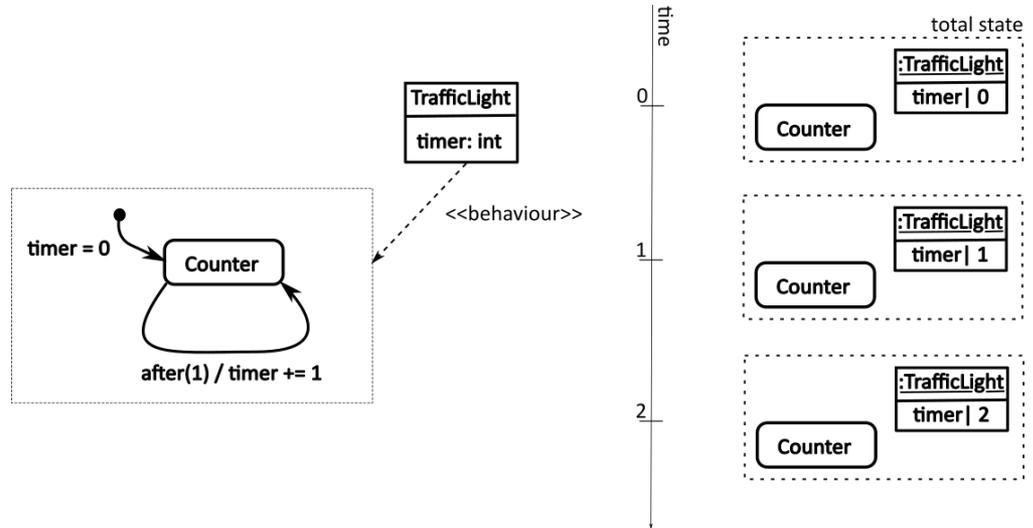
Output Event

$\wedge output_event(p_1, p_2, \dots, p_n)$

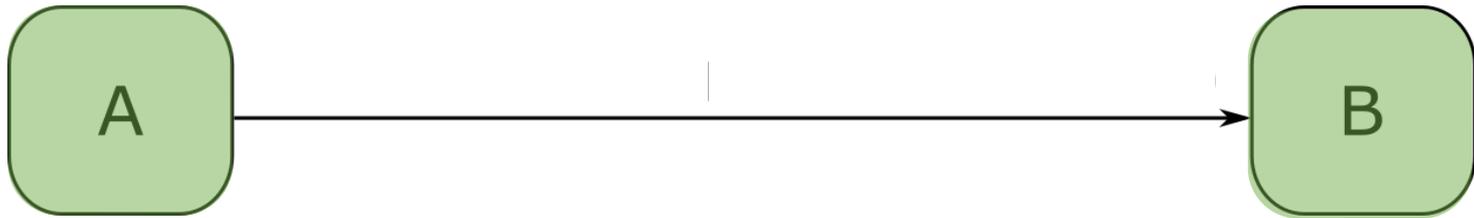


Assignment (to the non-modal part of the state)

- by action code in some appropriate language



Transitions

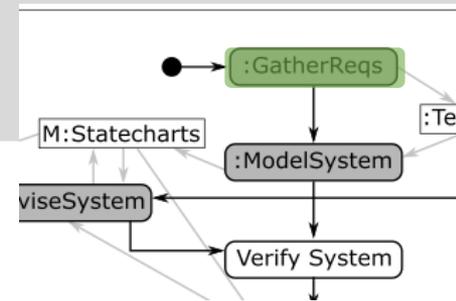


- Model the **dynamics** of the system:
 - *when*
 - the system is **in state A** and
 - **event is received** and
 - **guard** evaluates to **True**
 - *then*
 1. **output_action** is evaluated and
 2. the new **active state** becomes B

Exercise 2

Add data stores

Exercise 2 - Requirements



Your model here.

- R6': During the last 6 seconds of red being on, the traffic light announces to go to green by blinking its yellow light (1s on, 1s off) while leaving its red light on.
- R6: The time period of the different phases should be configurable.

TrafficLight

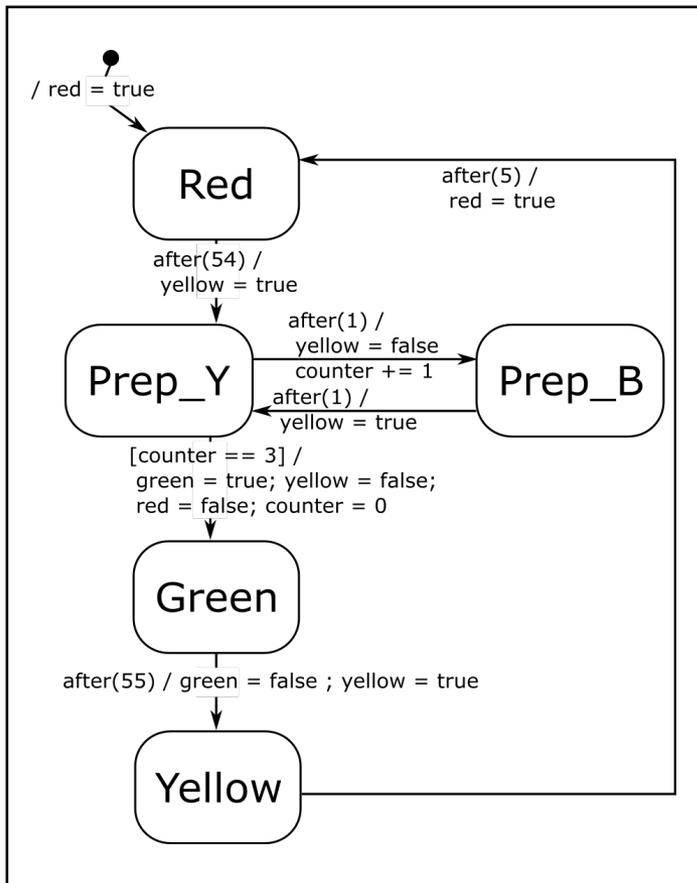
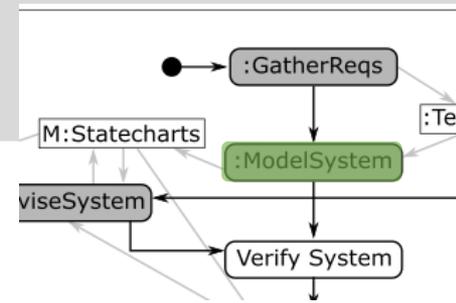
- counter: Integer = 0
- green: Boolean = false
- red: Boolean = false
- yellow: Boolean = false

Make sure that:

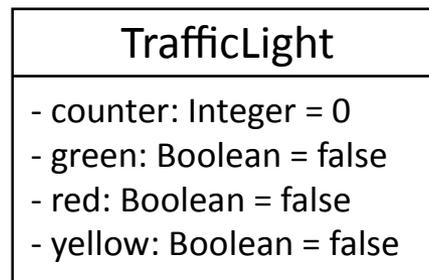
- **the values of the variables reflect which lights are on/off**
- **you use at least one conditional transition**

<<behavior>>

Exercise 2: Solution



- R6': During the last 6 seconds of red being on, the traffic light announces to go to green by blinking its yellow light (1s on, 1s off) while leaving its red light on.
- R6: The time period of the different phases should be configurable.



<<behavior>>

Statechart Execution

Run-To-Completion Step

- A Run-To-Completion (RTC) step is an *atomic execution step* of a state machine.
- It transitions the state machine from a *valid state configuration* into the next *valid state configuration*.
- RTC steps are executed one after the other - they must *not interleave*.
- New incoming events *cannot interrupt* the processing of the current event and must be stored in an *event queue*

Flat Statecharts: Simulation Algorithm (1)

```
1 simulate(sc: Statechart) {
```

```
18 }
```

Flat Statecharts: Simulation Algorithm (2)

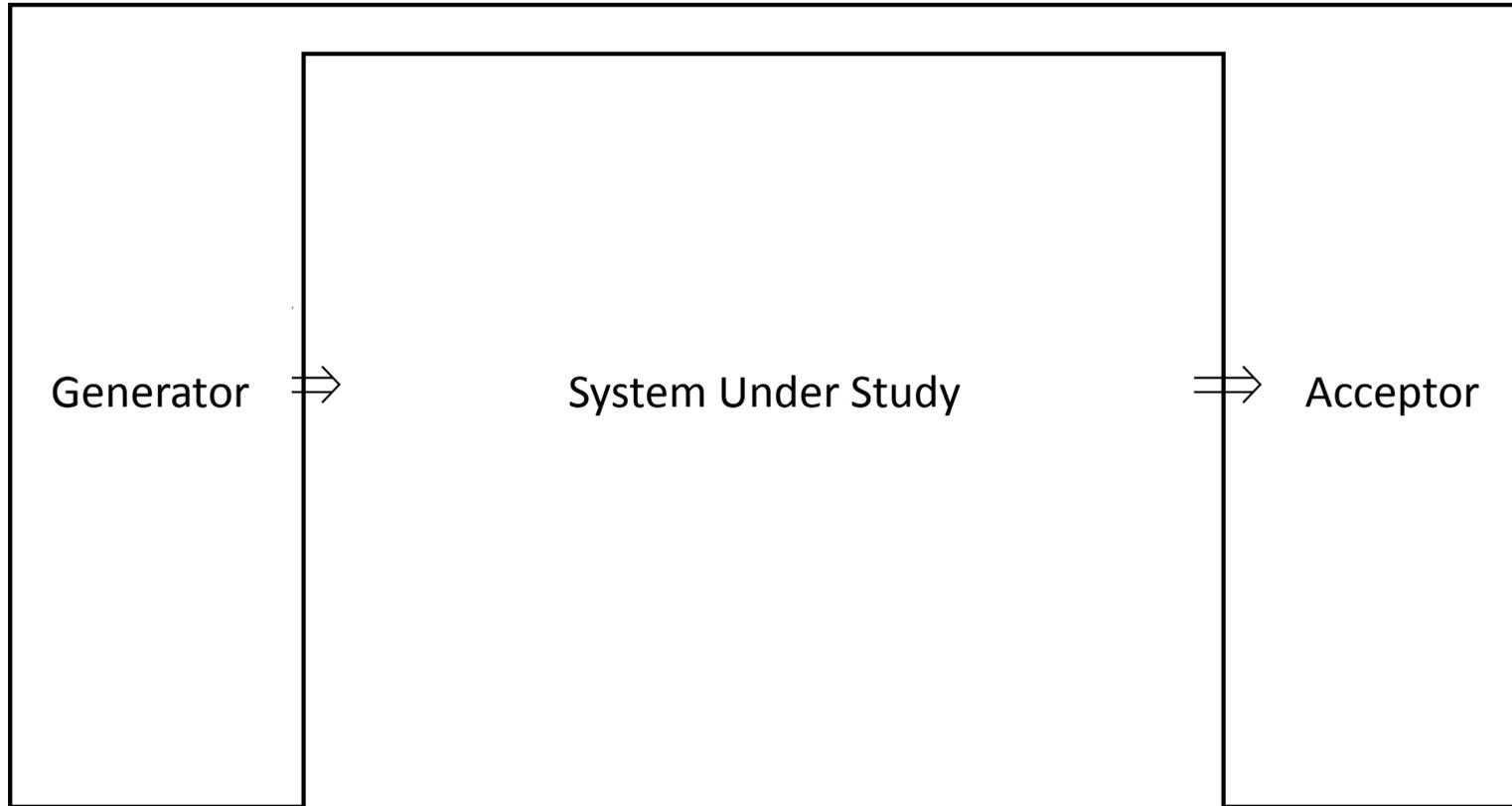
```
1 simulate(sc: Statechart) {
2     input_events = initialize_queue()
3     output_events = initialize_queue()
4     timers = initialize_set()
5     curr_state = sc.initial_state
6     for (var in sc.variables) {
7         var.value = var.initial_value
8     }
9     while (not finished()) {
10        curr_event = input_events.get()
11        while (not quiescent()) {
12            enabled_transitions = find_enabled_transitions(curr_state, curr_event, sc.variables)
13            chosen_transition = choose_one_transition(enabled_transitions)
14            cancel_timers(curr_state, timers)
15            curr_state = chosen_transition.target
16            chosen_transition.action.execute(sc.variables, output_events)
17            start_timers(curr_state, timers)
18        }
19    }
20 }
```

Flat Statecharts: Simulation Algorithm (3)

```
1  simulate(sc: Statechart) {
2      input_events = initialize_queue()
3      output_events = initialize_queue()
4      timers = initialize_set()
5      curr_state = sc.initial_state
6      for (var in sc.variables) {
7          var.value = var.initial_value
8      }
9      while (not finished()) {
10         curr_event = input_events.get()
11         enabled_transitions = find_enabled_transitions(curr_state, curr_event, sc.variables)
12         while (not quiescent()) {
13             chosen_transition = choose_one_transition(enabled_transitions)
14             cancel_timers(curr_state, timers)
15             curr_state = chosen_transition.target
16             chosen_transition.action.execute(sc.variables, output_events)
17             start_timers(curr_state, timers)
18             enabled_transitions = find_enabled_transitions(curr_state, sc.variables)
19         }
20     }
21 }
```

Testing Statecharts

Testing Statecharts



Zeigler BP. Theory of modelling and simulation. New York: Wiley-Interscience, 1976.

Mamadou K. Traoré, Alexandre Muzy, Capturing the dual relationship between simulation models and their context, Simulation Modelling Practice and Theory, Volume 14, Issue 2, February 2006, Pages 126-142.

- X-unit testing framework for YAKINDU Statechart Tools
- Test-driven development of Statechart models
- Test generation for various platforms
- Executable in YAKINDU Statechart Tools
- Virtual Time

Finished after 0,013 seconds

Runs: 1/1 Errors: 0 Failures: 0

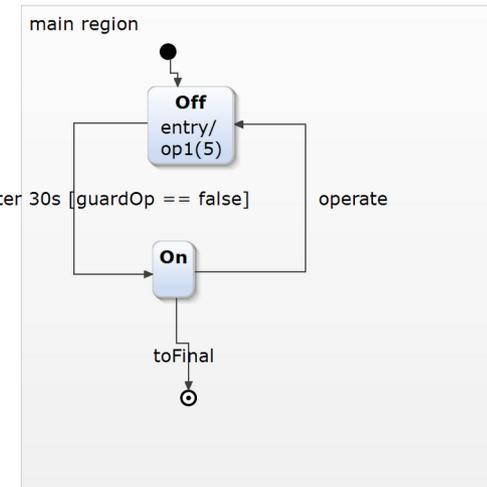
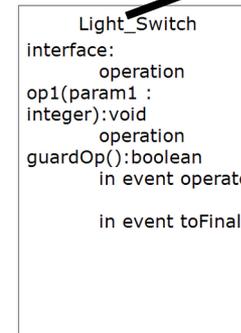
✓ org.yakindu.sct.LightSwitchTest [Runner: JUnit 4] (0,001 s)
 ✓ initialStateIsOff (0,001 s)

```
testclass LightSwitchTest for statechart Light_Switch{
    @Test
    operation initialStateIsOff(){
        enter
        assert active(Light_Switch.main_region.Off)
    }
}
```



```
testclass someTestClass for statechart Light_Switch {  
  
}
```

- Has a unique name
- Has a reference to a statechart
- Contains one or more operations





```
testsuite SomeTestSuite {  
    someTestClass  
}
```

- Has a unique name
- A testsuite contains at least one reference to a testclass



```
testclass someTestClass for statechart Light_Switch {  
    @Test  
    operation test(): void{  
        enter  
    }  
}
```

- May have @Test or @Run annotation
- Has a unique name
- May have 0..n parameters
- Has a return type (standard is void)
- Contains 0..n statements



// entering / exiting the statechart

enter, exit

// raising an event

raise event : value

// proceeding time or cycles

proceed 2 **cycle**

proceed 200 **ms**

// asserting an expression, expression must evaluate to boolean

assert expression

// is a state active

active(someStatechart.someRegion.someState)



SCTUnit allows to

- mock operations defined in the statechart model
- verify that an operation was called with certain values

// mocking the return value of an operation

mock mockOperation **returns** (20)

mock mockOperation(5) **returns** (30)

// verifying the call of an operation

assert called verifyOperation

assert called verifyOperation **with** (5, 10)



Control Structures

// if expression

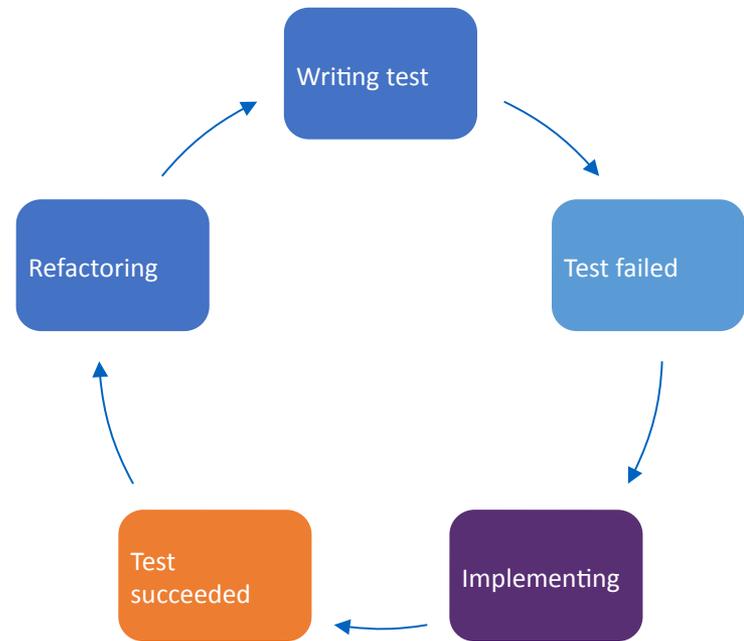
```
if (x==5) {  
    doSomething()  
} else {  
    doSomethingelse()  
}
```

// while expression

```
while (x==5) {  
    doSomething()  
}
```

Test-Driven Development

- Software development process, where software is developed driven by tests
- Test-first-approach
- 3 steps you do repeatedly:
 - writing a test
 - implementing the logic
 - refactoring



Exercise 3

Testing Models

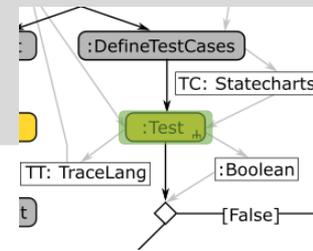
Exercise 3 – Unit testing Statecharts

The statechart editor displays a statechart for TrafficLightCtrl with the following structure:

- Initial state: main
- State Red (green box):
 - Transition: / raise displayRed
 - After redPeriod s: / raise displayGreer
- State Green (yellow box):
 - Transition: / raise displayYellow
 - After greenPeriod s: / raise displayYellow
- State Yellow (red box):
 - Transition: / raise displayRed
 - After yellowPeriod s: / raise displayRed

The JUnit runner shows the following coverage table:

Model Element	Coverage
Statechart TrafficLightCtrl	50 % (6)
Region main	50 % (6)
State Red	100 % (2)
→ Red -> Green (after	100 % (1)
State Green	50 % (2)
→ Green -> Yellow (at	0 % (1)
State Yellow	0 % (2)
→ Yellow -> Red (after	0 % (1)



- Create a test that checks the following requirements:
 - R3: at system start-up, the red light is on
 - R4: cycles through red on, green on, and yellow on
 - R5: red is on for 60s, green is on for 55s, yellow is on for 5s

Exercise 3 – Solution

```
package trafficlight.test
```

```
testclass TrafficLightTests for statechart TrafficLightCtrl {
```

```
  @Test operation switchTrafficLightOn () {
```

```
    // given the traffic light is inactive
    assert !is_active
    // when
    enter
    // then traffic light is off which means no color was switched on
    assert displayRed
    assert !displayGreen
    assert !displayYellow
  }
```

```
  @Test operation switchLightFromRedToGreen () {
```

```
    // given
    switchTrafficLightOn
    // when
    proceed 60s
    // then
    assert displayGreen
  }
```

```
  @Test operation switchLightFromGreenToYellow () {
```

```
    // given
    switchLightFromRedToGreen
    // when
    proceed 55s
    // then
    assert displayYellow
  }
```

```
  @Test operation switchLightFromYellowToRed () {
```

```
    // given
    switchLightFromGreenToYellow
    // when
    proceed 5s
    // then
    assert displayRed
  }
```

```
  @Test operation lightCycles () {
```

```
    // given
    switchLightFromYellowToRed
```

```
    var i : integer = 10
```

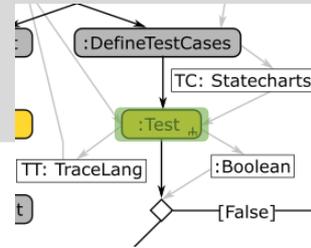
```
    while (i > 0) {
      i=i-1
```

```
      //when
      proceed 60 s
      // then
      assert displayGreen
```

```
      //when
      proceed 55 s
      // then
      assert displayYellow
```

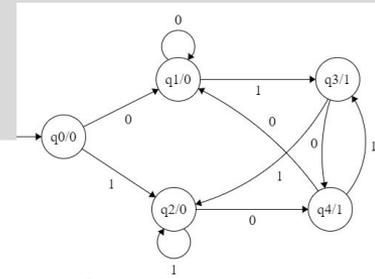
```
      //when
      proceed 5 s
      // then
      assert displayRed
    }
```

```
  }
```

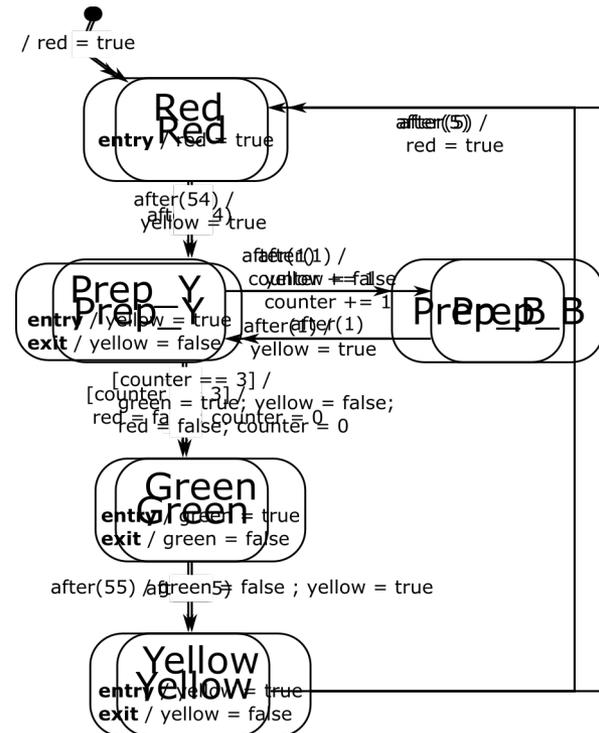
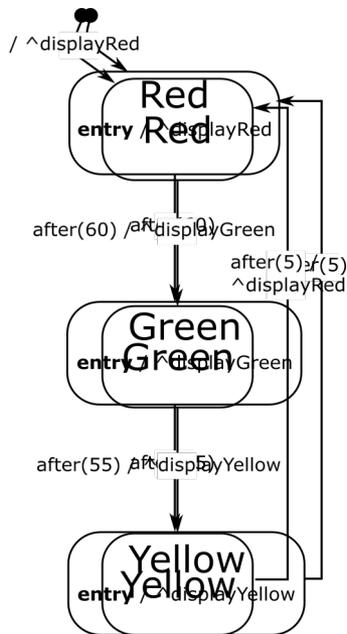


Hierarchy

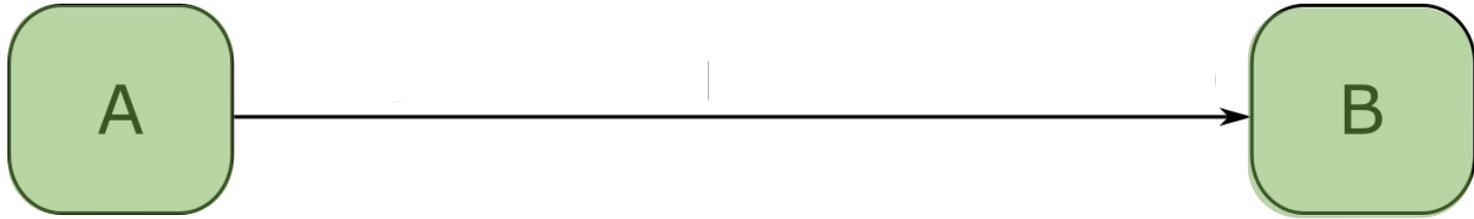
Entry/Exit Actions



- A state can have entry and exit actions.
- An *entry action* is executed whenever a state is entered (made active).
- An *exit action* is executed whenever a state is exited (made *inactive*).
- Same expressiveness as *transition actions* (i.e., syntactic sugar).



Transitions



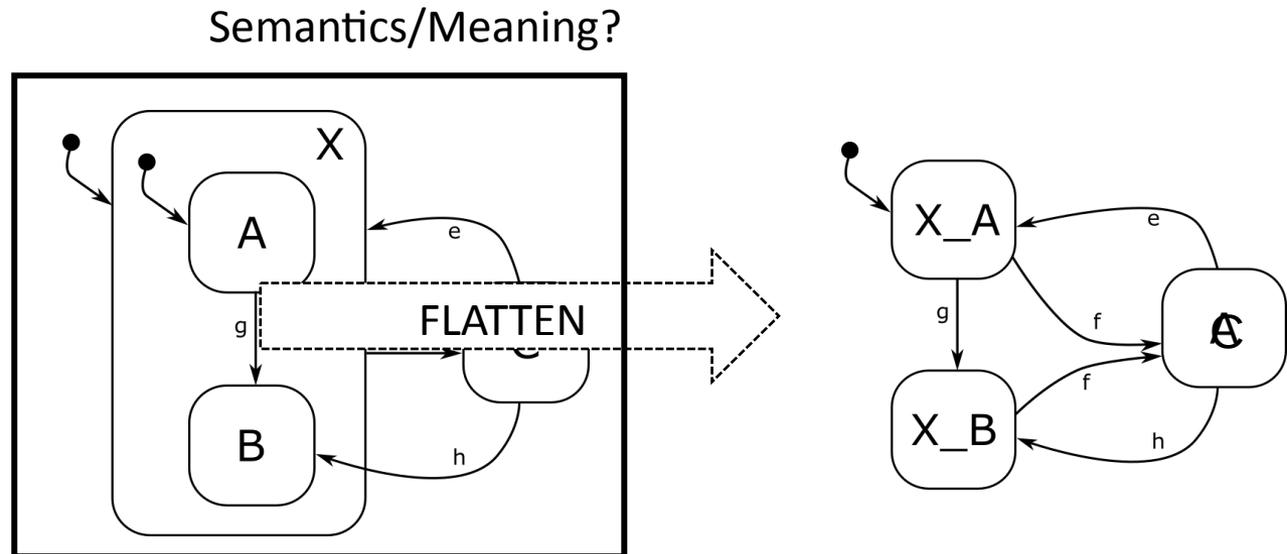
- Model the **dynamics** of the system:
 - *when*
 - the system is **in state A** and
 - **event is received** and
 - **guard** evaluates to **true**
 - *then*
 1. the **exit actions** of state A are evaluated and
 2. **output_action** is evaluated and
 3. the **enter actions** of state B are evaluated and
 4. the new **active state** becomes B

Entry/Exit Actions: Simulation Algorithm

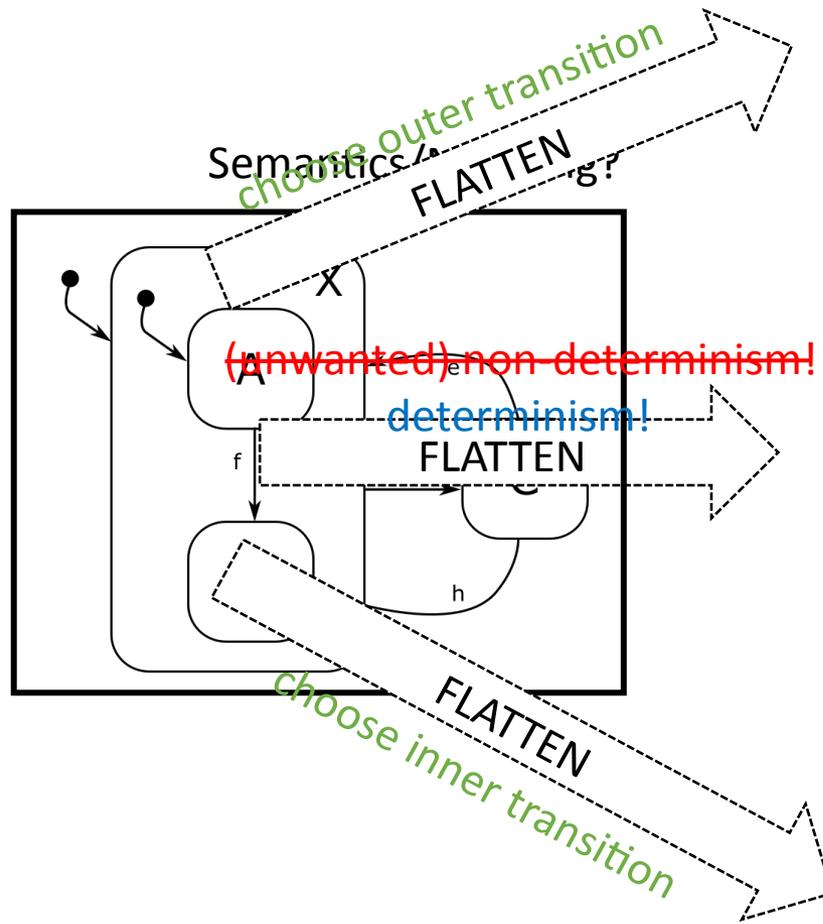
```
1 simulate(sc: Statechart) {
2     input_events = initialize_queue()
3     output_events = initialize_queue()
4     timers = initialize_set()
5     curr_state = sc.initial_state
6     for (var in sc.variables) {
7         var.value = var.initial_value
8     }
9     while (not finished()) {
10        curr_event = input_events.get()
11        enabled_transitions = find_enabled_transitions(curr_state, curr_event, sc.variables)
12        while (not quiescent()) {
13            chosen_transition = choose_one_transition(enabled_transitions)
14            cancel_timers(curr_state, timers)
15            execute_exit_actions(curr_state)
16            curr_state = chosen_transition.target
17            chosen_transition.action.execute(sc.variables, output_events)
18            execute_enter_actions(curr_state)
19            start_timers(curr_state, timers)
20            enabled_transitions = find_enabled_transitions(curr_state, sc.variables)
21        }
22    }
23 }
24
```

Hierarchy

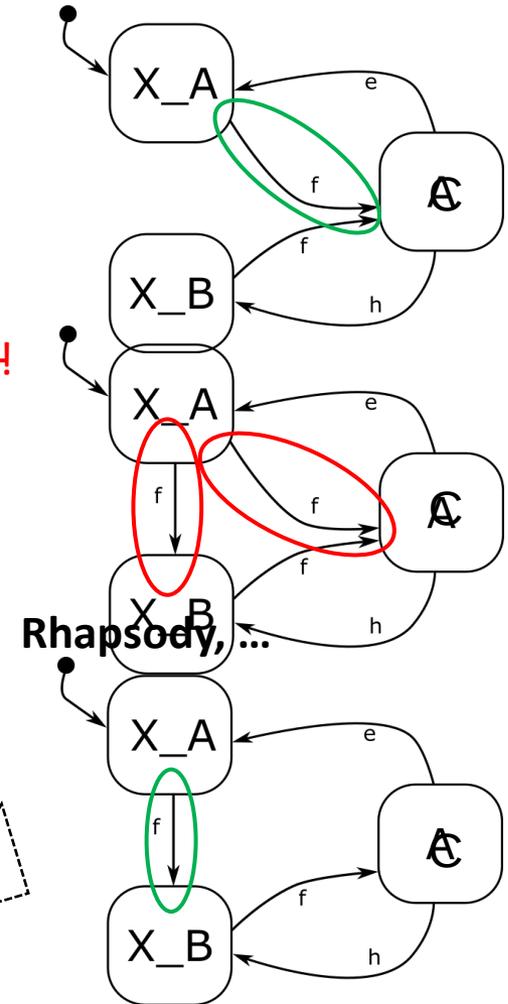
- Statechart states can be hierarchically (de-) **composed**
- Each hierarchical state has exactly one **initial/default state**
- An active hierarchical state has **exactly one active child** (down to leaf/atomic state)



Hierarchy: Modified Example



Statemate, Yakindu, ...

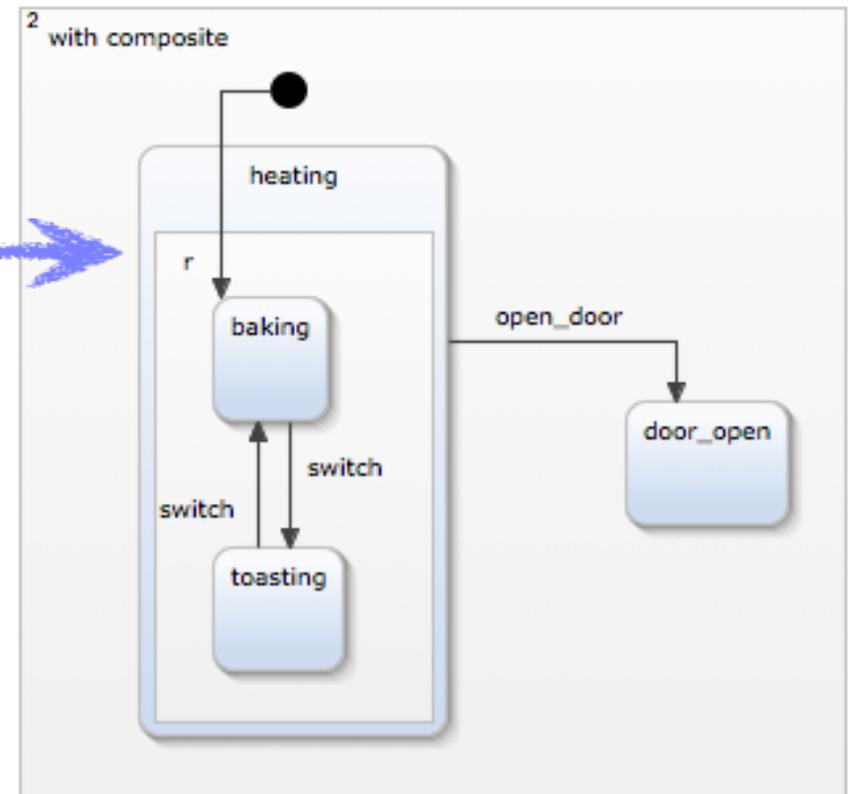
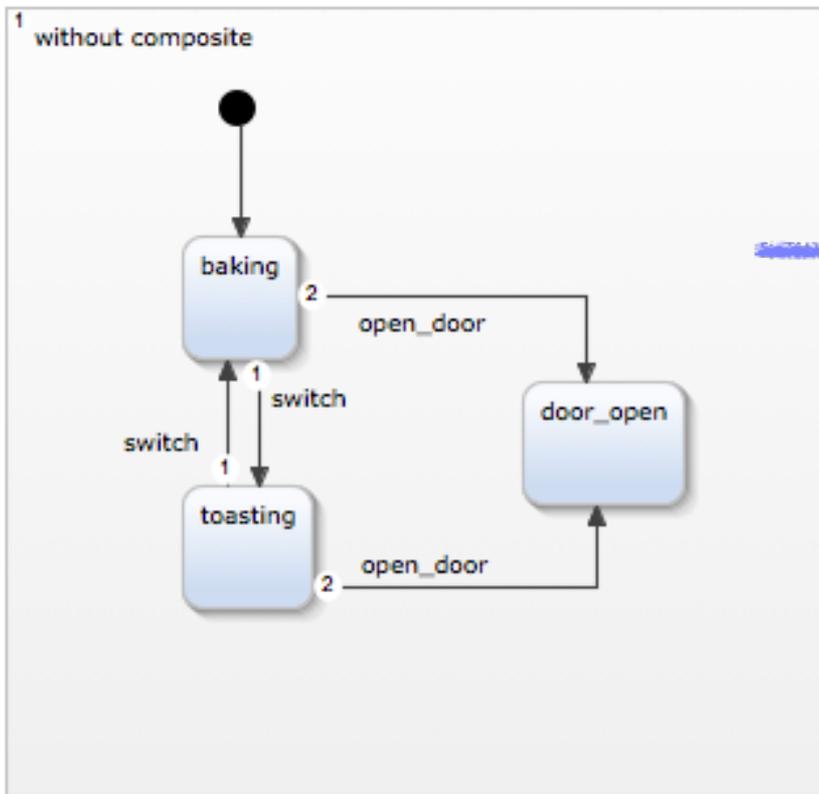


Hierarchy: why inner? ... see Code Generation

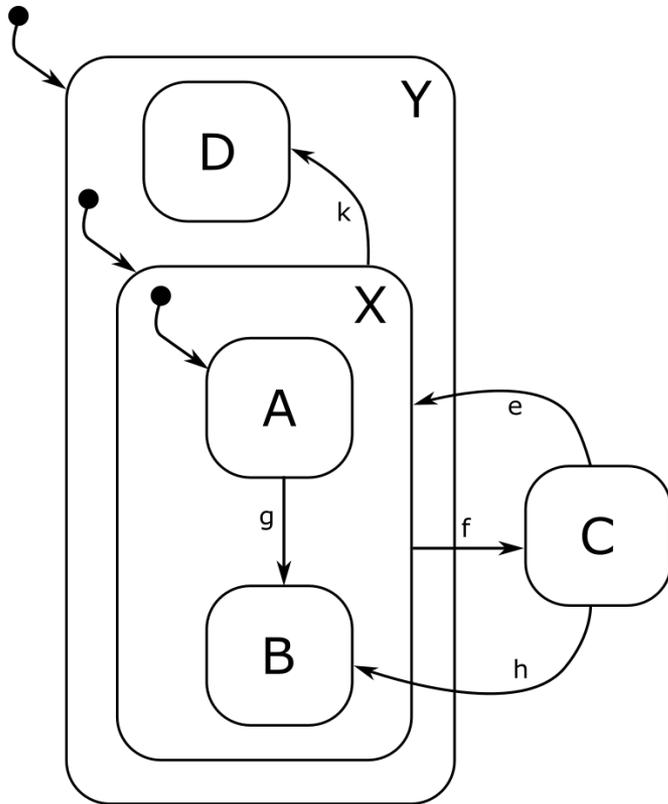


Composite States

- Hierarchical states are an ideal mechanism for hiding complexity
- Parent states can implement common behaviour for their substates
- Hierarchical event processing **reduces the number of transitions**
- Refactoring support: group states into a composite state

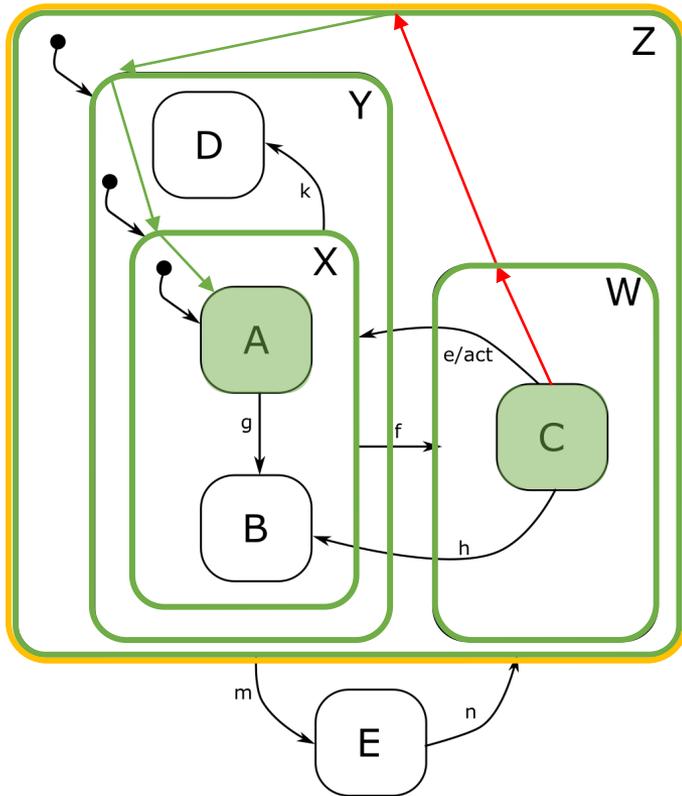


Hierarchy: Initialization



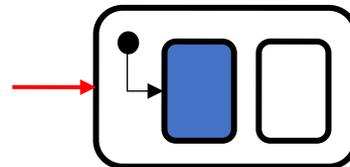
- Concept of **effective target state**
 - Recursive: the effective target state of a composite state is its initial state
- Effective target state of initial transition is $Y/X/A$
- Initialization:
 1. Enter Y, execute enter action
 2. Enter X, execute enter action
 3. Enter A, execute enter action

Hierarchy: Transitions



- Assume $Z/W/C$ is active and e is processed.
- Semantics:
 1. Find LCA, collect states to leave
 2. Leave states up the hierarchy
 3. Execute action act
 4. Find effective target state set, enter states down the hierarchy

Effective target states:

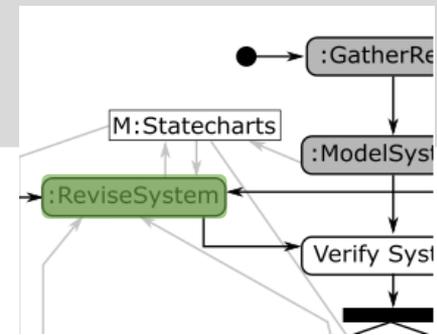
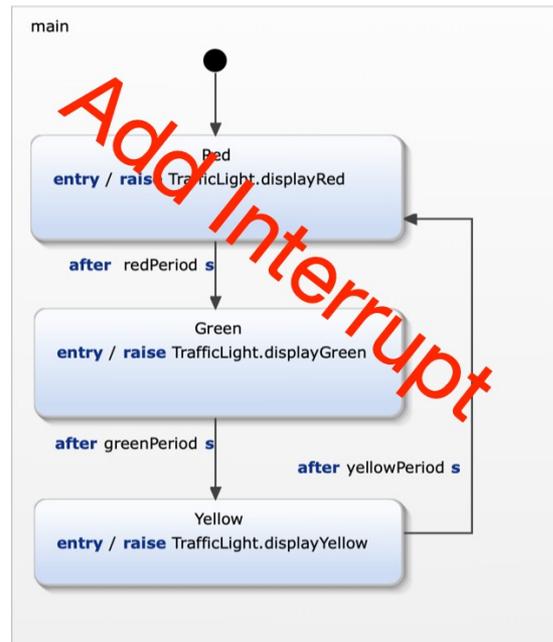


RECURSIVE!

Exercise 5

Model an
interruptible traffic light

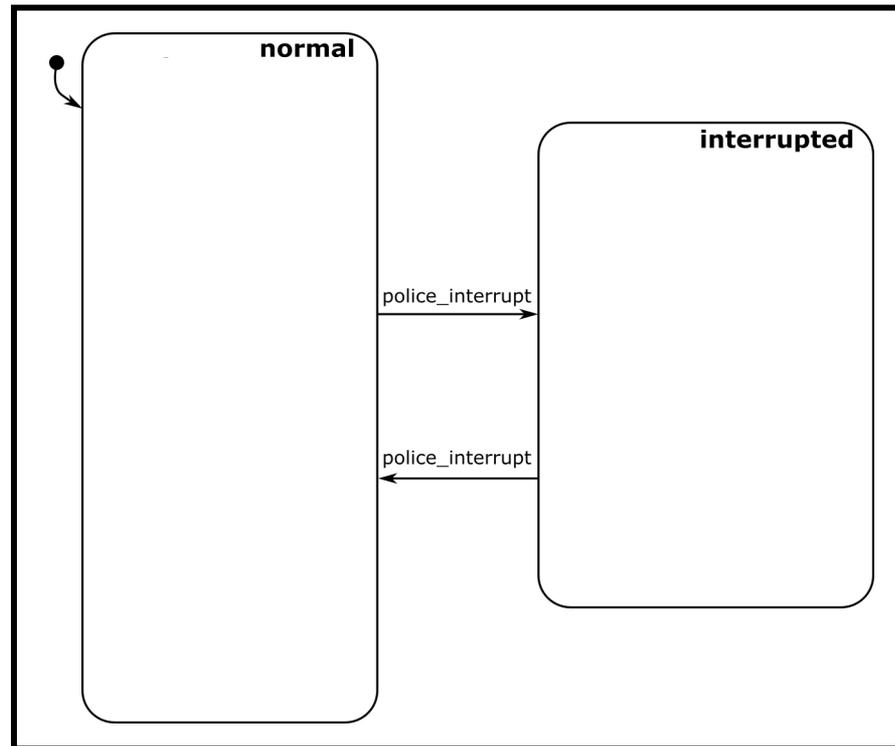
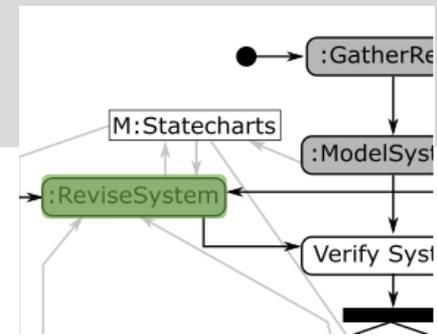
Exercise 5 - Requirements



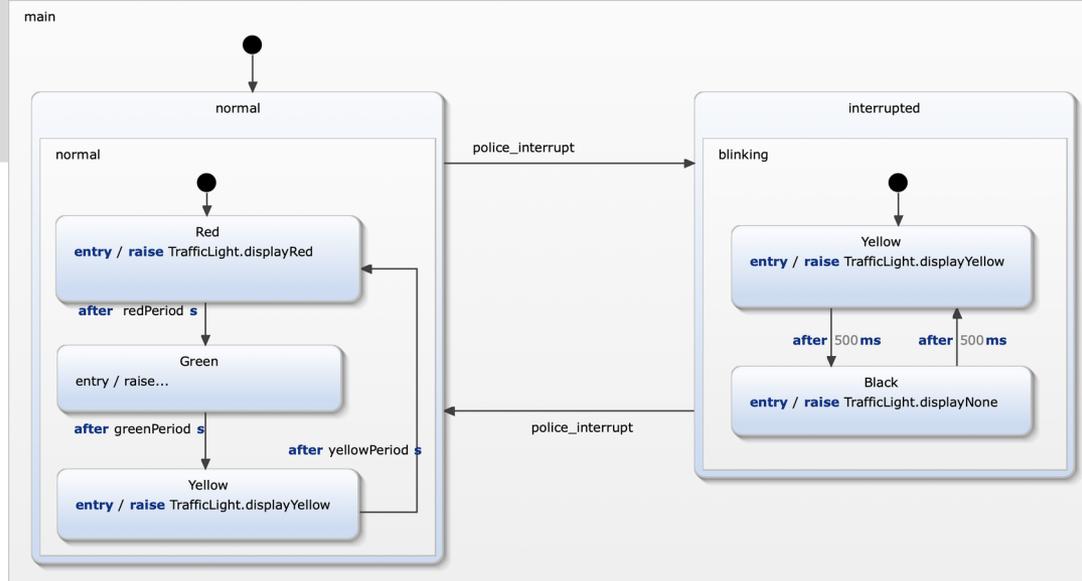
- R7a: police can interrupt autonomous operation .
- R7b: autonomous operation can be interrupted during any phase of constant red, yellow and green lights.
- R7c: in interrupted mode the yellow light blinks with a constant frequency of 1 Hz (on 0.5s, off 0.5s).
- R8a: police can resume to regular autonomous operation.
- R8b: when regular operation is resumed, the traffic light restarts with red (R) light on.

Exercise 5: Solution

- R7a: police can interrupt autonomous operation .
- R7b: autonomous operation can be interrupted during any phase of constant red, yellow and green lights.
- R7c: in interrupted mode the yellow light blinks with a constant frequency of 1 Hz (on 0.5s, off 0.5s).
- R8a: police can resume to regular autonomous operation.
- R8b: when regular operation is resumed, the traffic light restarts with red (R) light on.



Exercise 5 - Solution

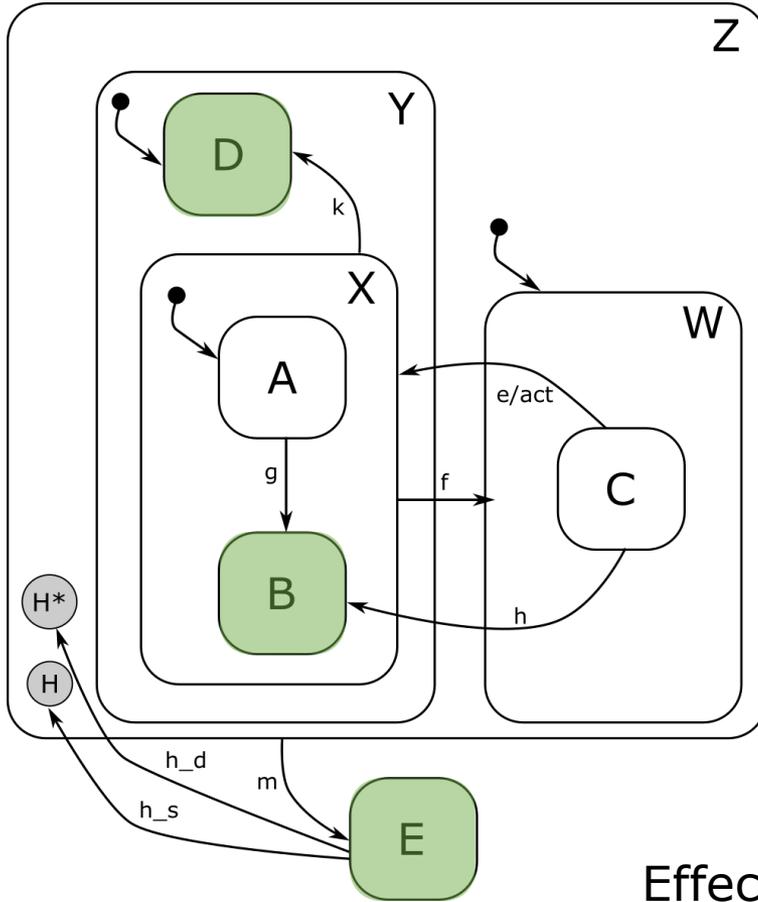


requirement	modelling approach
R6: police can interrupt autonomous operation.	An new incoming event <code>police_interrupt</code> triggers a transition to a new state <code>interrupted</code> .
R6a: autonomous operation can be interrupted during any phase of constant red, yellow and green lights.	The states <code>Red</code> , <code>Green</code> , and <code>Yellow</code> are grouped within a new composite state <code>normal</code> . This state is the source state of the transition to state <code>interrupted</code> and thus also applies to all substates.
R7: in interrupted mode the yellow light blinks with a constant frequency of 1 Hz. (on 0.5s, off 0.5s).	State <code>interrupted</code> is a composite state with two substates <code>Yellow</code> and <code>Black</code> . These switch the yellow light on and off. Timed transitions between these states ensure correct timing for blinking.
R8: police can resume to regular autonomous operation.	A transition triggered by <code>police_interrupt</code> leads from state <code>interrupted</code> to state <code>normal</code> .
R8a: when regular operation is resumed the traffic light restarts with red (R) light on.	When activating state <code>normal</code> its substate <code>Red</code> is activated by default.

History

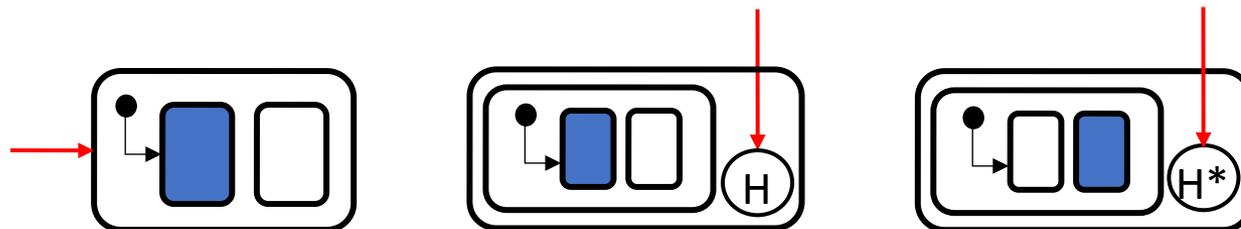
History: pseudo-states

H shallow history
 H* deep history



- Assume $Z/Y/X/B$ is active, and m is processed
 - Effective target state: E
- If h_s is processed
 - Effective target state: $Z/Y/D$
- If h_d is processed
 - Effective target state: $Z/Y/X/B$

Effective target states:

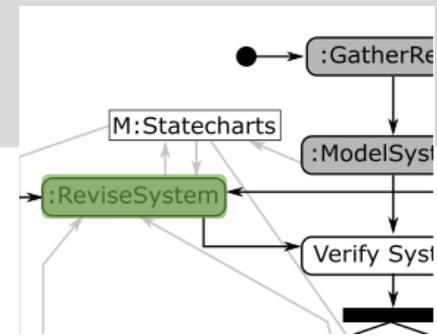


RECURSIVE!

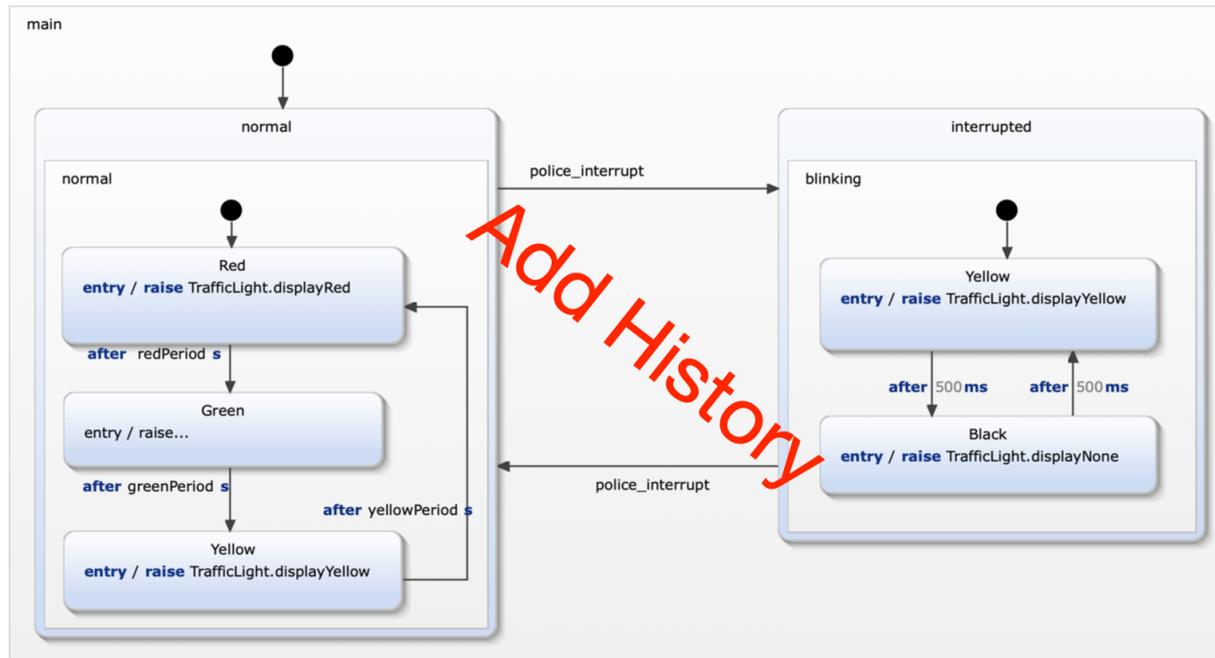
Exercise 6

Model an interruptible traffic light that restores its state

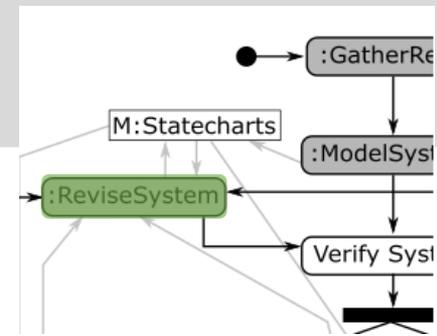
Exercise 6: Requirements



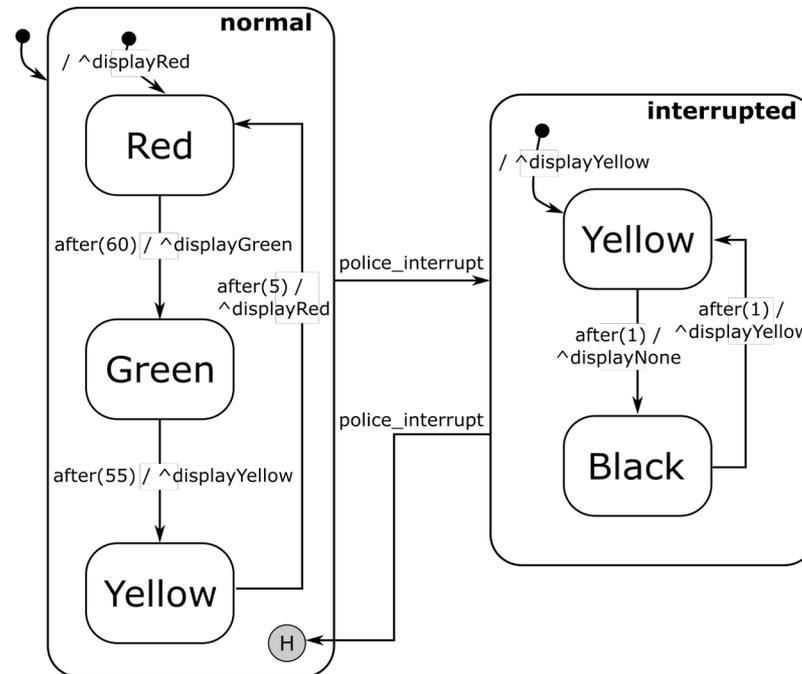
- R8b: when regular operation is resumed the traffic light restarts with the last active light color red (R), green (G), or yellow (Y) on.



Exercise 6: Solution



- R8b: when regular operation is resumed the traffic light restarts with the last active light color red (R), green (G), or yellow (Y) on.

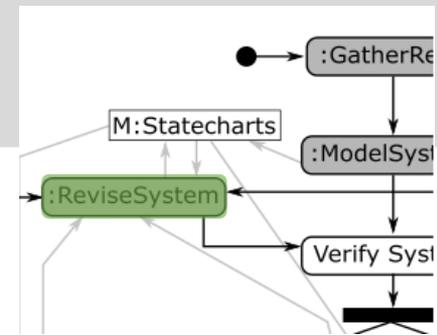


Exercise 7

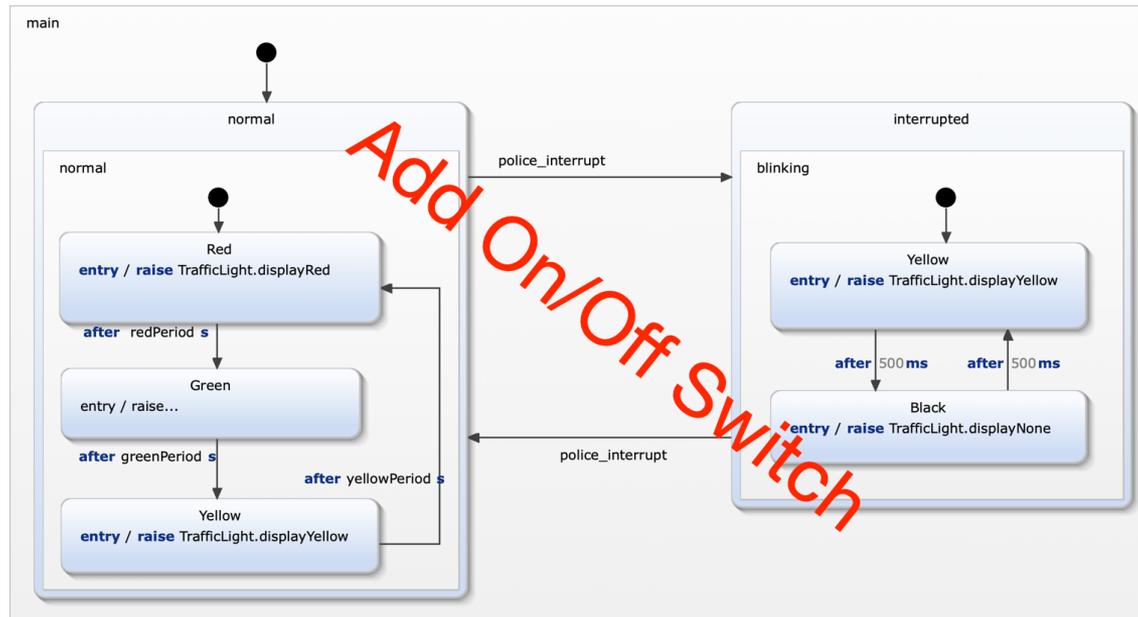
Model an interruptible traffic light that restores its state and can be switched on/off

Exercise 7: Requirements

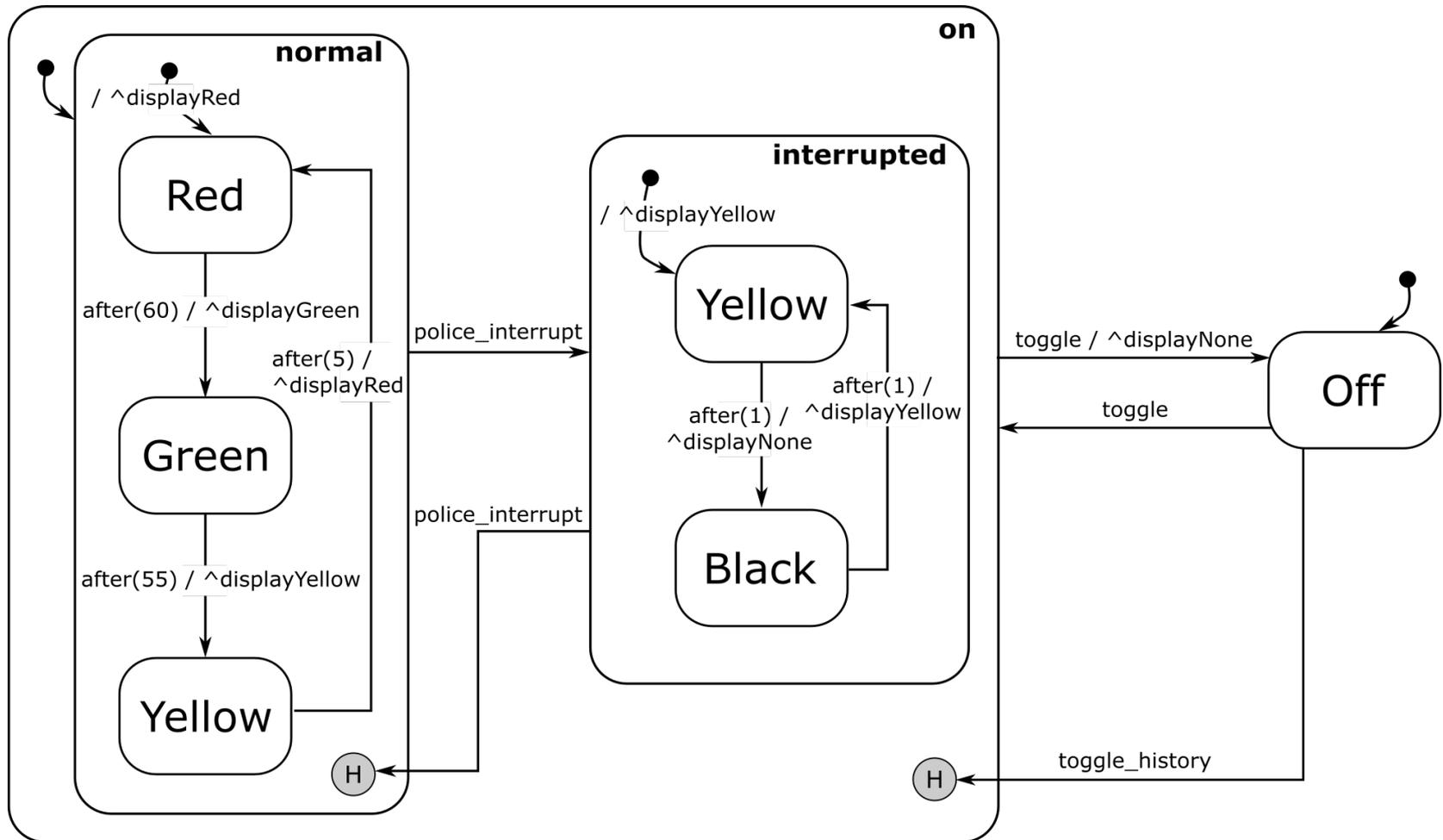
Add another level of hierarchy that supports switching on and off the entire traffic light. Go into detail with shallow and deep histories.



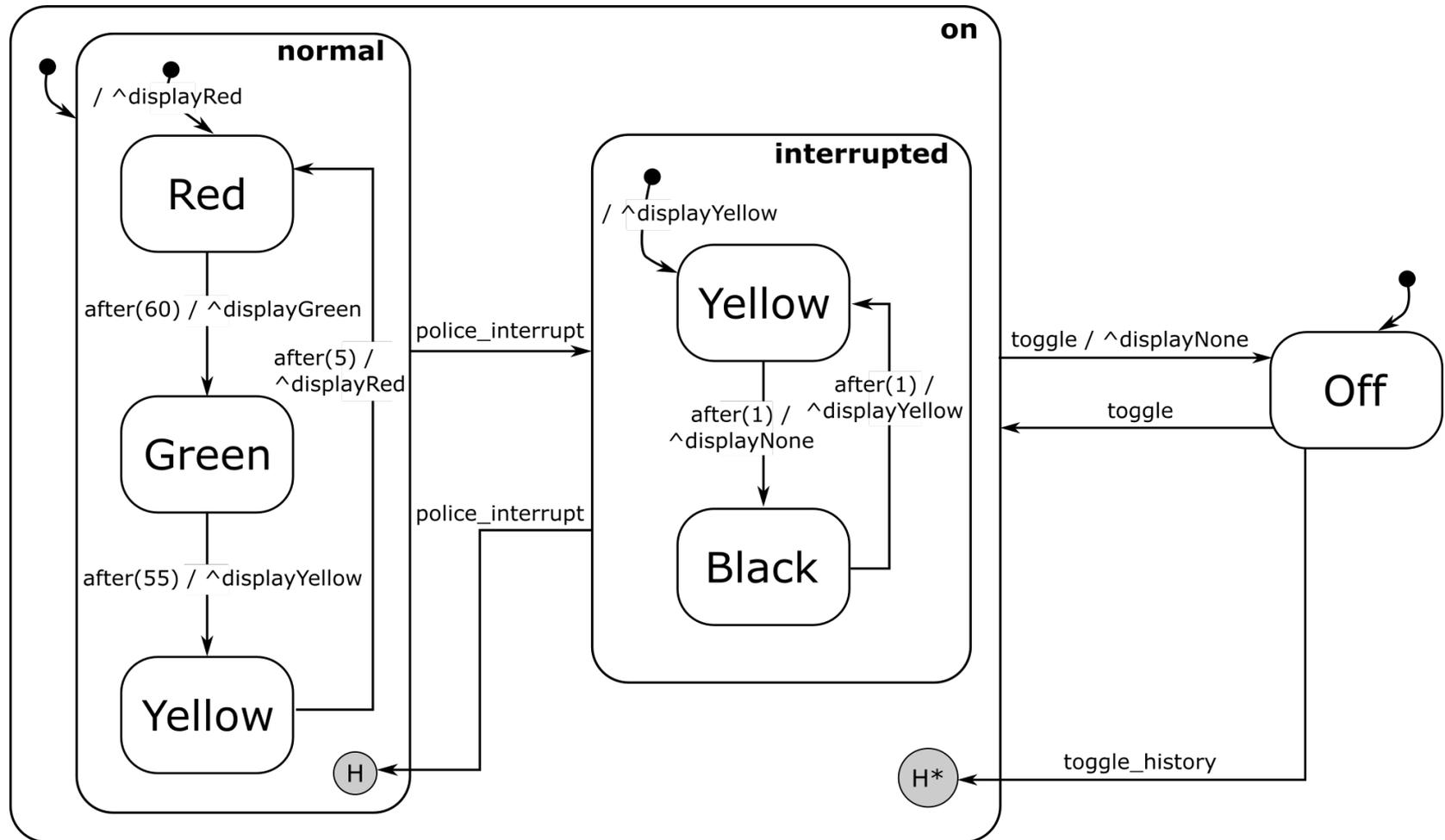
- R9: The traffic light can be switched on and off.
- R9a: The traffic light is initially off.
- R9b: If the traffic light is off none of its lights (R/G/Y) are on.
- R9c: After switching off and on again the traffic light must switch on the light that was on before the switching off.



Exercise 7: Solution



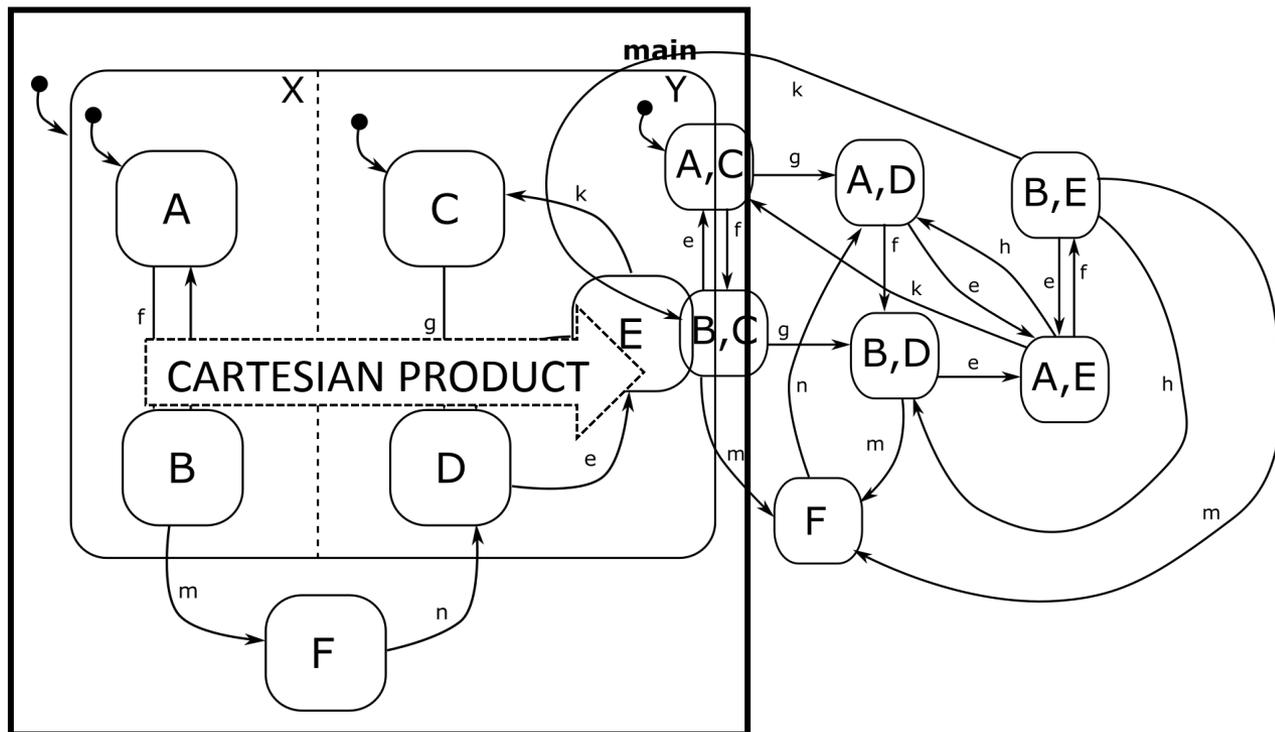
Exercise 7: Alternative Solution



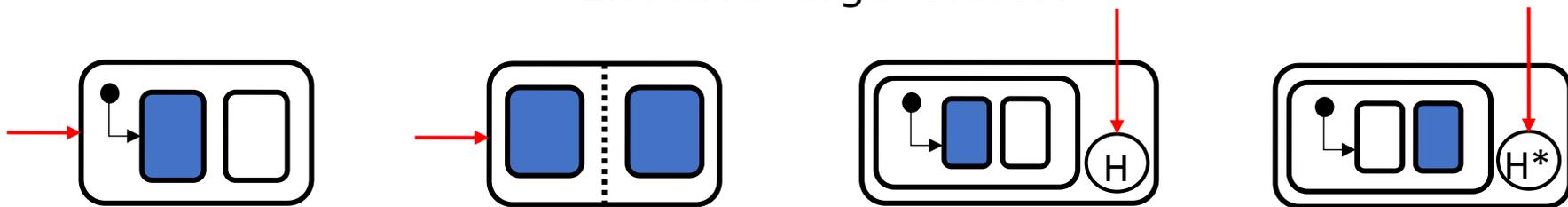
Orthogonality

Orthogonal Components/Regions: "and" states

Semantics/Meaning?

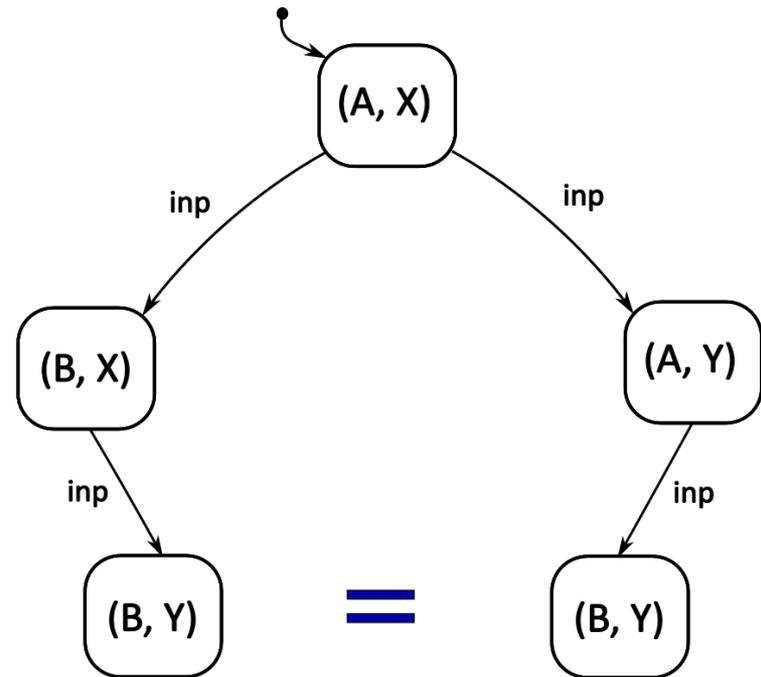
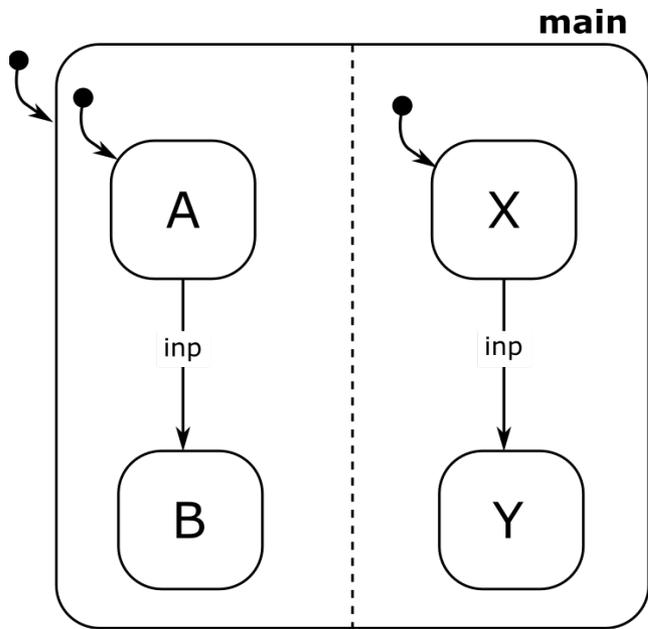


Effective target states:

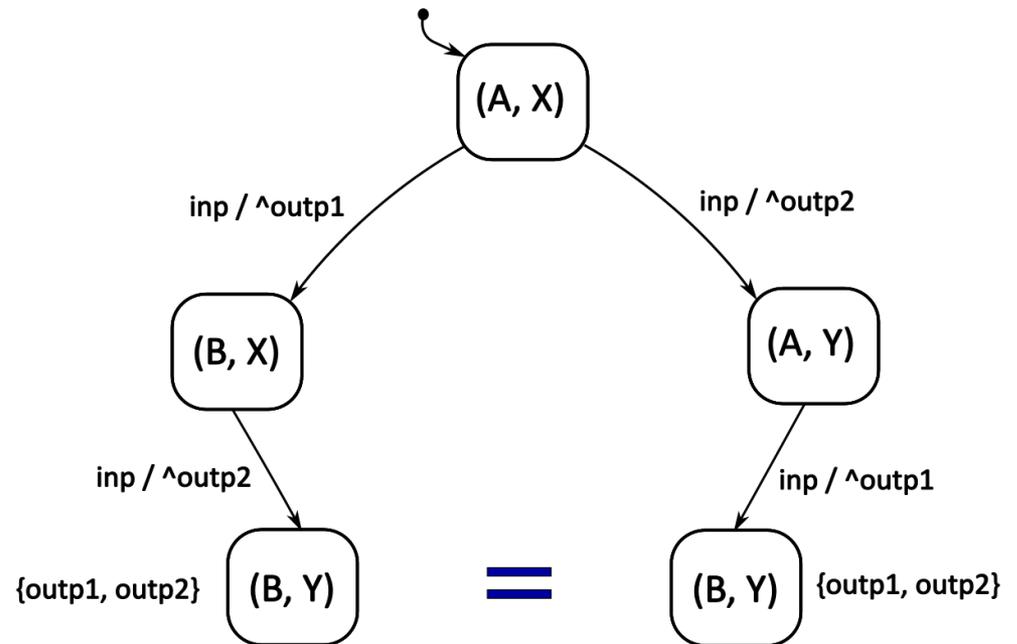
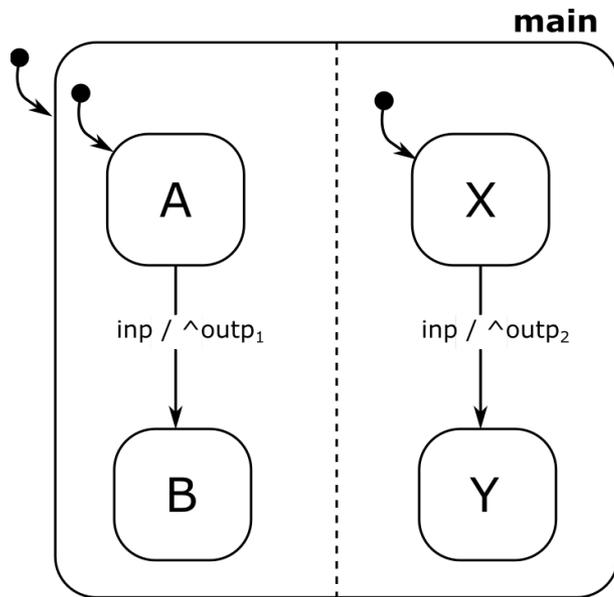


RECURSIVE!

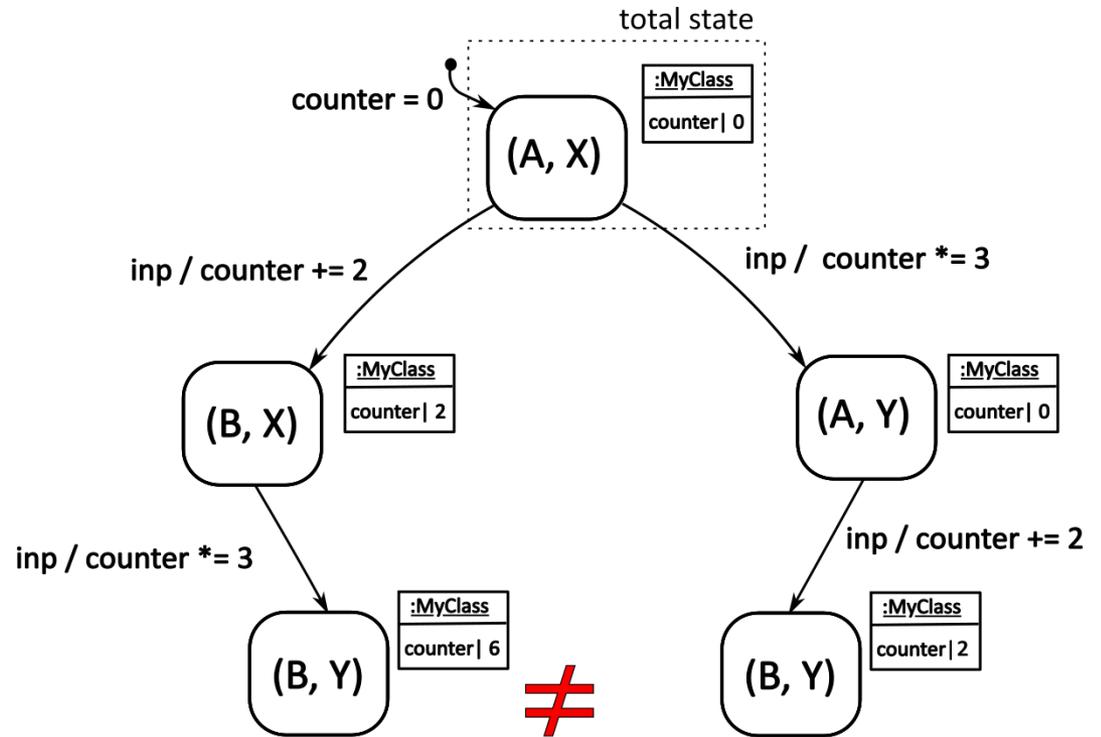
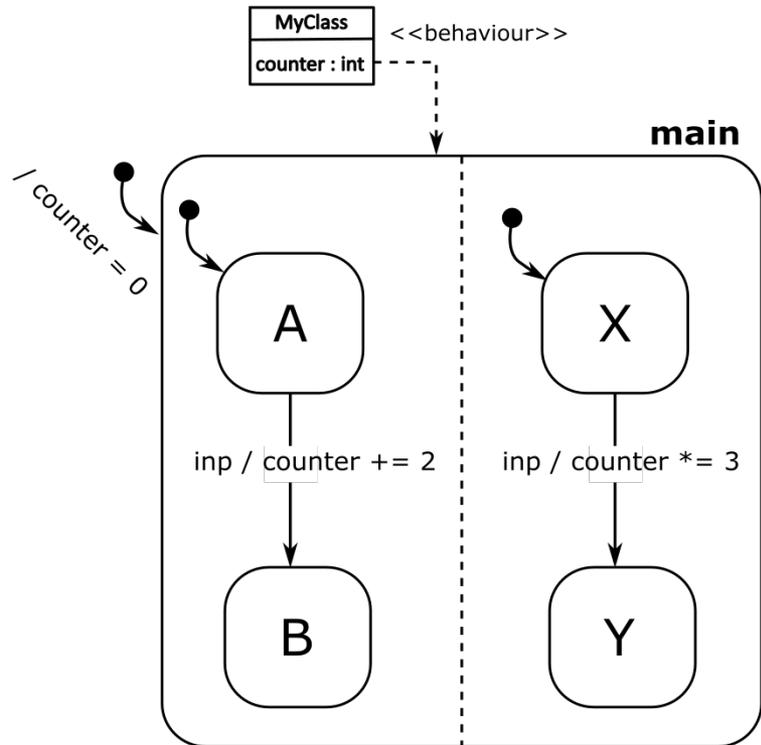
Parallel (In)Dependence



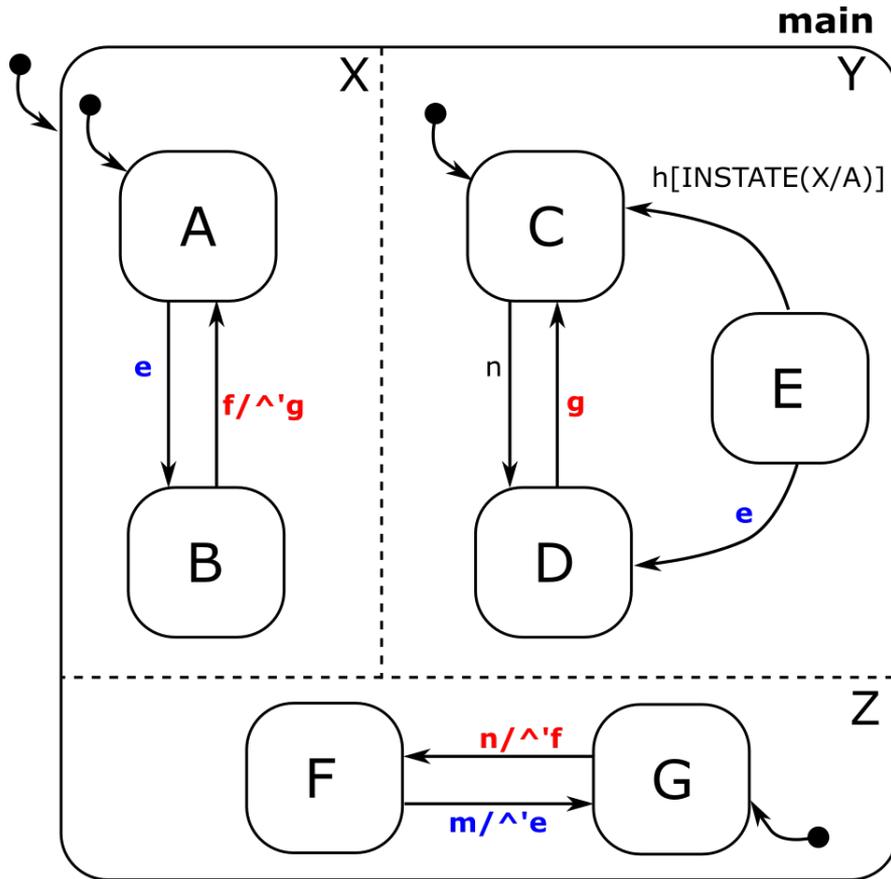
Parallel (In)Dependence



Parallel (In)Dependence



Orthogonality: Communication



Input Segment: **nmnn**

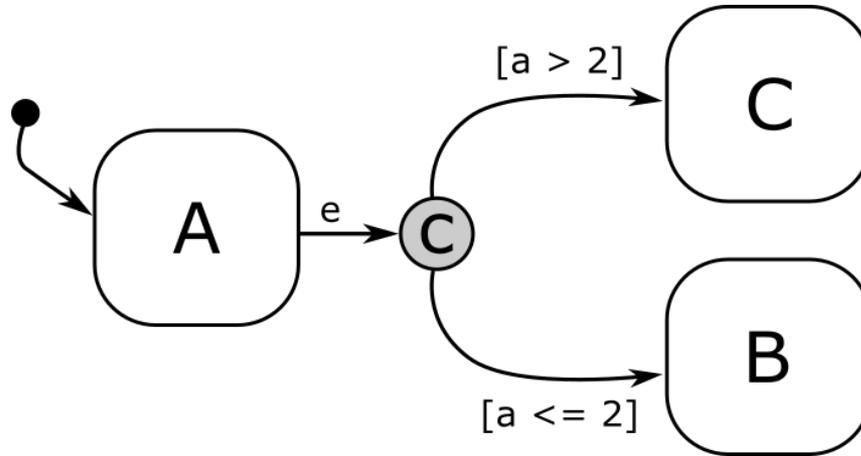
Orthogonal Components can communicate:

- raising/broadcasting local events:
^'<<event name>>
- state reference is a transition guard:
INSTATE(<<state location>>)

Simulation Algorithm

```
1  simulate(sc: Statechart) {
2      input_events = initialize_queue()
3      output_events = initialize_queue()
4      local_events = initialize_queue()
5      timers = initialize_set()
6      curr_state = get_effective_target_states(sc.initial_state)
7      for (var in sc.variables) {
8          var.value = var.initial_value
9      }
10     while (not finished()) {
11         curr_event = input_events.get()
12         for (region in sc.orthogonal_regions) {
13             enabled_transitions[region] = find_enabled_transitions(curr_state, curr_event, sc.variables)
14         }
15         while (not quiescent()) {
16             chosen_region = choose_one_region(sc.orthogonal_regions)
17             chosen_transition = choose_one_transition(enabled_transitions[chosen_region])
18             states_to_exit = get_states_to_exit(get_lca(curr_state, chosen_transition))
19             for (state_to_exit in states_to_exit) {
20                 cancel_timers(state_to_exit, timers)
21                 execute_exit_actions(state_to_exit)
22                 remove_state_from_curr_state(state_to_exit)
23             }
24             chosen_transition.action.execute(sc.variables, output_events, local_events)
25             states_to_enter = get_effective_target_states(chosen_transition)
26             for (state_to_enter in states_to_enter) {
27                 add_state_to_curr_state(state_to_enter)
28                 execute_enter_actions(state_to_enter)
29                 start_timers(state_to_enter, timers)
30             }
31             enabled_transitions = find_enabled_transitions(curr_state, sc.variables, local_events)
32         }
33     }
34 }
```

Conditional Transitions



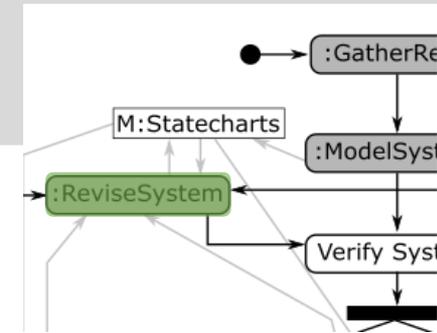
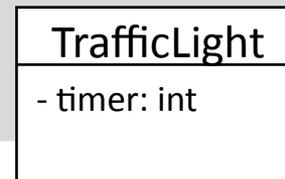
- `getEffectiveTargetStates()`: select one *True*-branch
- Conditions should not overlap to avoid non-determinism
 - in Yakindu, priority makes deterministic
 - “else” branch is required
- Equivalent (in this case) to two transitions:
 - $A - e[a > 2] \rightarrow C$
 - $A - e[a \leq 2] \rightarrow B$

Exercise 8

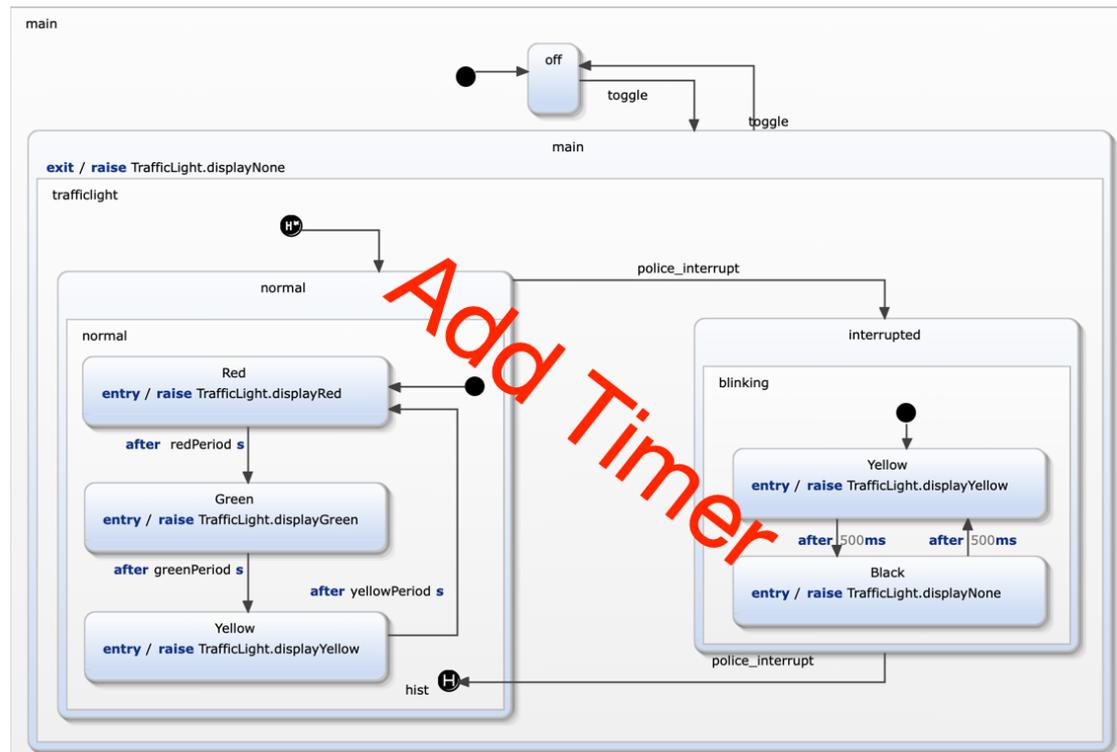
Add a timer
to the traffic light

Exercise 8: Requirements

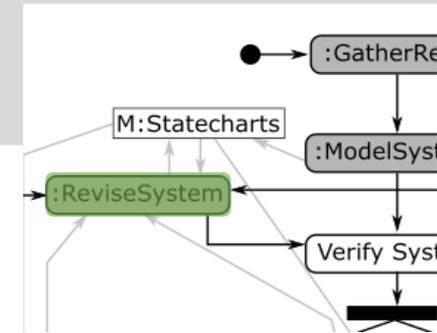
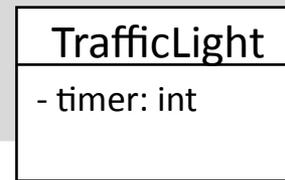
In this exercise a timer must be modelled. It introduces the use of orthogonal regions.



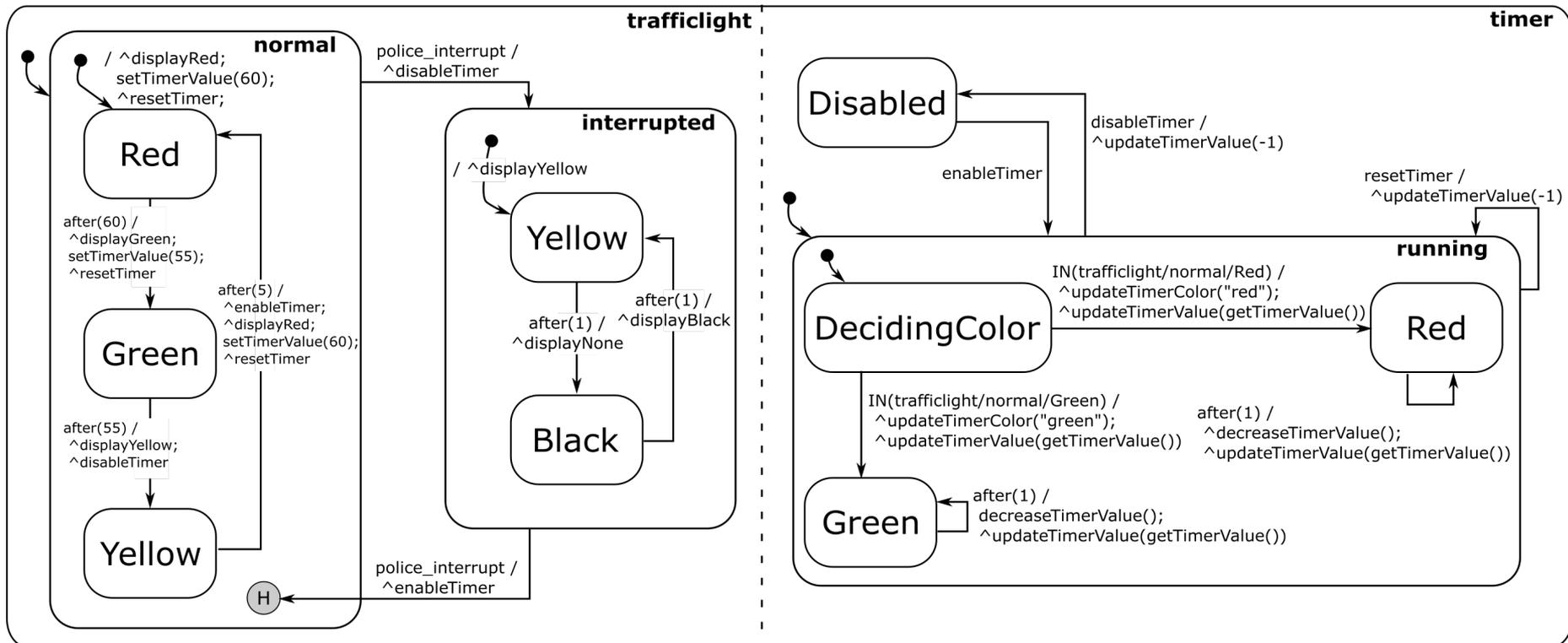
- R10a: A timer displays the remaining time while the light is red or green.
- R10b: This timer decreases and displays its value every second.
- R10c: The colour of the timer reflects the colour of the traffic light.



Exercise 8: Solution

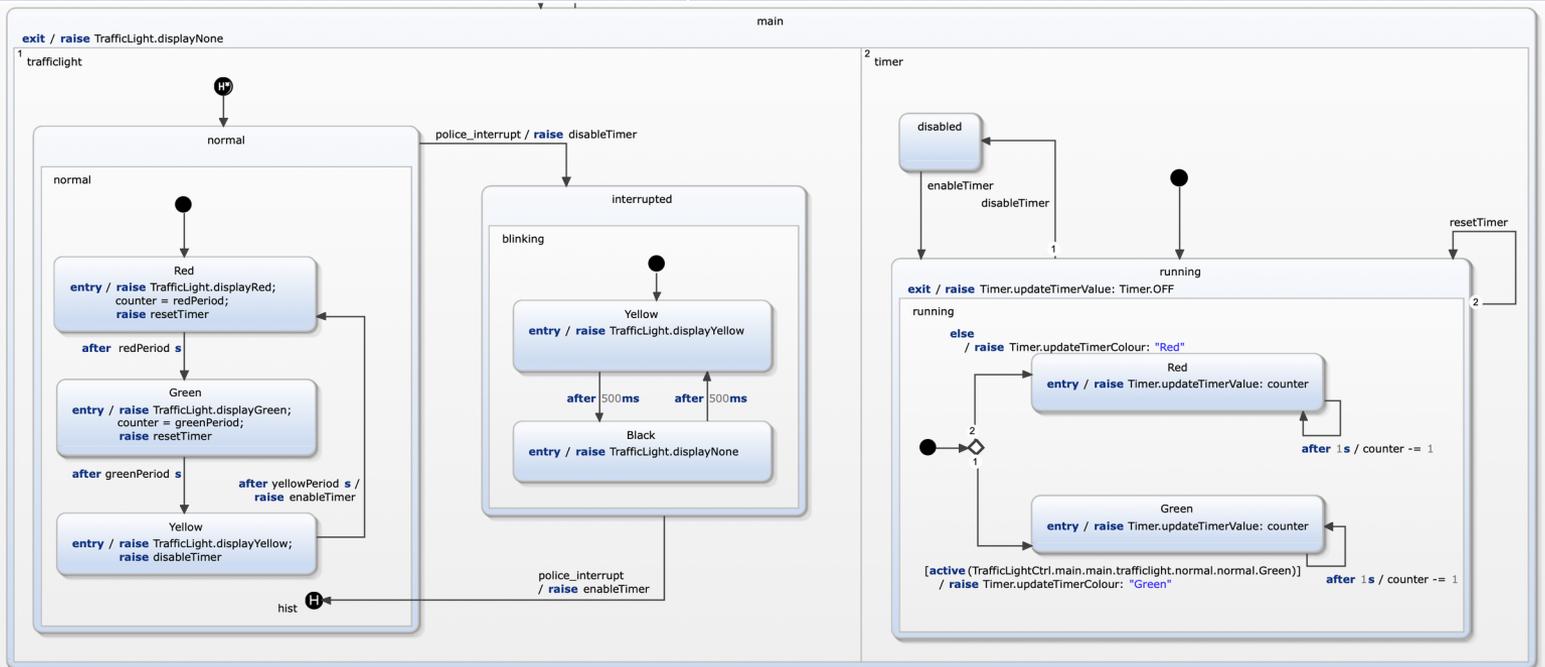


- R10a: A timer displays the remaining time while the light is red or green.
- R10b: This timer decreases and displays its value every second.
- R10c: The colour of the timer reflects the colour of the traffic light.



Solution 8

requirement	modelling approach
<p>R10: a timer displays the remaining time while the light is red or green</p>	<p>The timer is defined in a second region within state on (main in the Yakindu model).</p>
<p>R10a: This timer decreases and displays its value every second.</p>	<p>An internal variable for the counter is introduced. When switching the traffic light phase, the counter value is set to how long the light has been in that phase. Additionally, the local events <code>resetTimer</code>, <code>enableTimer</code>, and <code>disableTimer</code> are used to synchronize traffic light phase switches with the timer.</p>
<p>R10b: The colour of the timer reflects the colour of the traffic light.</p>	<p>When the timer is enabled it checks the active traffic light phase using the <code>active()</code> function.</p>



Yakindu syntax

Yakindu:

- **raise** e == ^e

- strict alternation between “or” and “and” states

TrafficLightCtrl.**main.main**.trafficlight.**normal.normal**.Green

- **active()** == INSTATE() == IN()

Code Generation



- Code generators for C, C++, Java, Python, Swift, Typescript, SCXML
- Plain-code approach by default
- Very efficient code
- Easy integration of custom generators



Python



JavaScript



Code Generation

- Various different approaches for implementing a state machine:
 - switch / case
 - state transition table
 - state pattern
- Which one is the best depends on
 - Runtime (performance, predictability) requirements
 - ROM vs. RAM memory
 - Debugging capabilities
 - Clarity and maintainability
(of generated code ~ certification, round-trip)

Switch / Case

- Each state corresponds to one “case”
- Each case executes state-specific statements and state transitions

```
public void stateMachine() {  
    while (true) {  
        switch (activeState) {  
            case RED: {  
                activeState = State.RED_YELLOW;  
                break;  
            }  
            case RED_YELLOW: {  
                activeState = State.GREEN;  
                break;  
            }  
            case GREEN: {  
                activeState = State.YELLOW;  
                break;  
            }  
            case YELLOW: {  
                activeState = State.RED;  
                break;  
            }  
        }  
    }  
}
```

State Transition Table

- Specifies the state machine purely declaratively.
- One of the dimensions indicates current states, while the other indicates events.

```
enum columns {
    SOURCE_STATE,
    USER_UP, USER_DOWN, POSSENSOR_UPPER_POSITION, POSSENSOR_LOWER_POSITION,
    TARGET_STATE
};

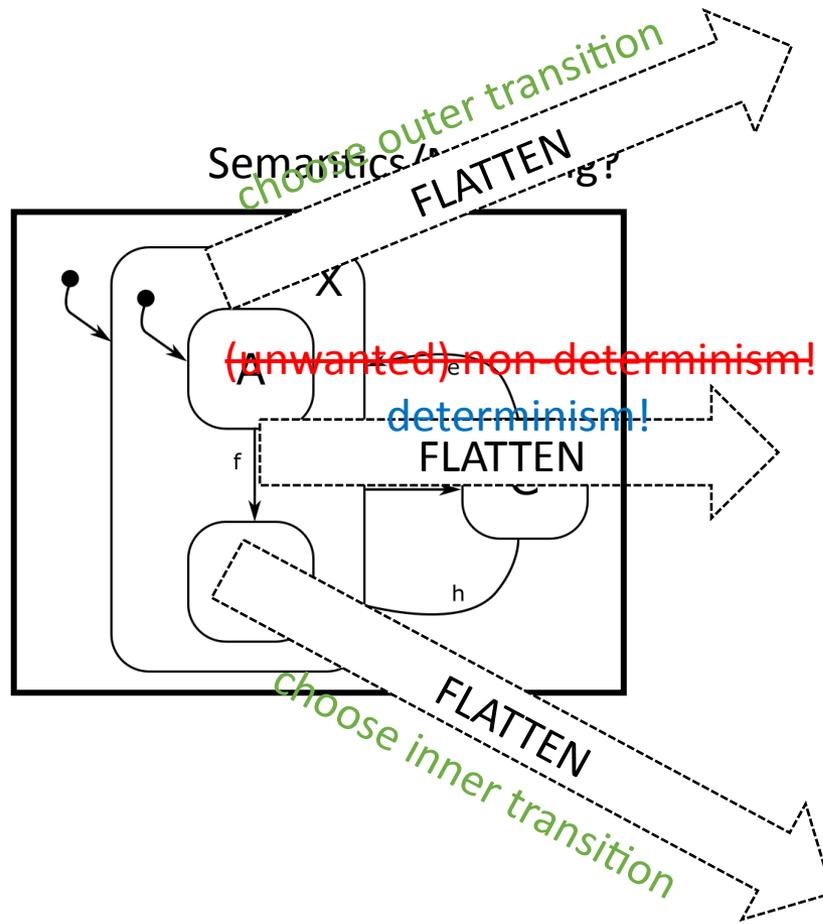
#define ROWS 7
#define COLS 6
int state_table[ROWS][COLS] = {
/*      source,      up,      down, upper, lower, target */
    { INITIAL,      false, false, false, false, IDLE },
    { IDLE,         true,  false, false, false, MOVING_UP },
    { IDLE,         false, true,  false, false, MOVING_DOWN },
    { MOVING_UP,    false, true,  false, false, IDLE },
    { MOVING_UP,    false, false, true,  false, IDLE },
};
```

State Pattern

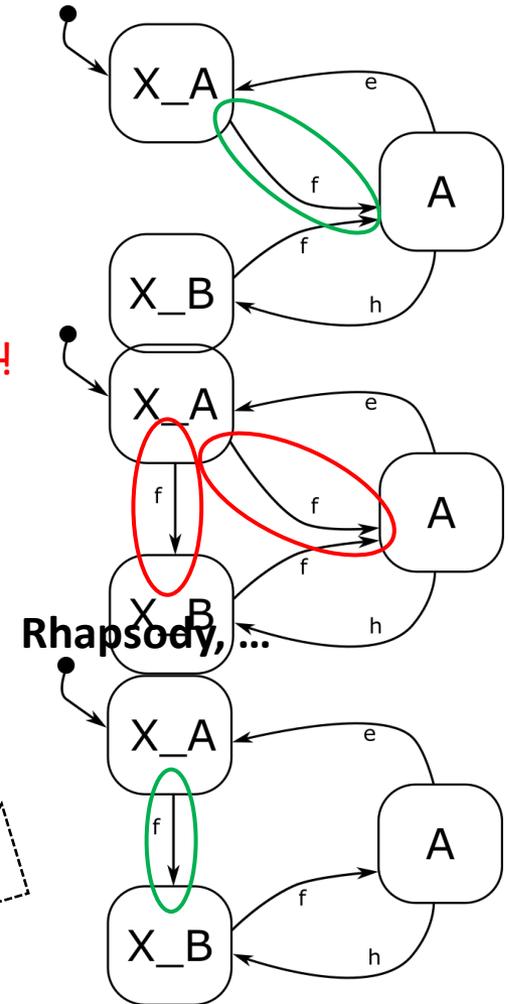
- Object-oriented implementation, behavioural design pattern
- Used by frameworks such as Spring Statemachine, Boost MSM or Qt State Machine Framework
- Each State becomes a class, events become methods
- All classes derive from a common interface

```
public class MovingUp extends AbstractState {  
  
    public MovingUp(StateMachine stateMachine) {  
        super(stateMachine);  
    }  
  
    @Override  
    public void raiseUserDown() {  
        stateMachine.activateState(new Idle(stateMachine));  
    }  
  
    @Override  
    public void raisePosSensorUpperPosition() {  
        stateMachine.activateState(new Idle(stateMachine));  
    }  
  
    @Override  
    public String getName() {  
        return "Moving up";  
    }  
  
}
```

Hierarchy: outer vs. inner

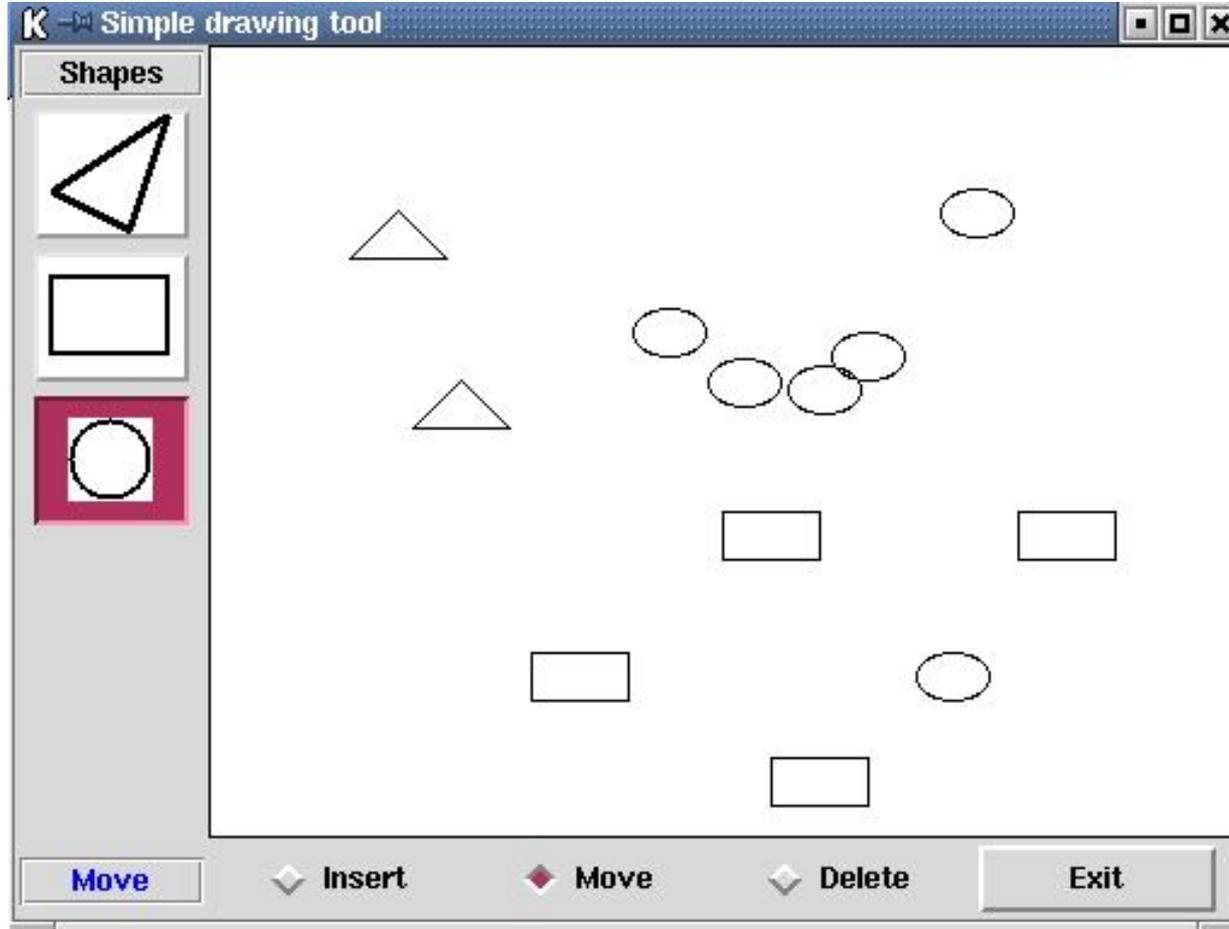


Statemate, Yakindu, ...



Hierarchy: why inner?

Manual Object-Oriented implementation (no Statechart compiler)
using State Pattern



Code Generation

	Fast	Memory efficient	easy to debug	Easy to understand
 Switch / Case				
State Transition Table				
State Pattern				



very simplified illustration



Code Generator Model

```
GeneratorModel for yakindu::java {  
  statechart exercise5 {  
    feature Outlet {  
      targetProject = "5_sctunit"  
      targetFolder = "src-gen"  
      libraryTargetFolder = "src"  
    }  
  }  
}
```

- Has a generator ID
- Has a generator entry
- Each generator entry contains 1..n feature-configurations
- Each feature-configuration contains 1..n properties

Generated Code

Sample

Files

- src-gen
 - traffic.light
 - trafficlightctrl
 - ITrafficLightCtrlStateMachine.java
 - SynchronizedTrafficLightCtrlStateMachine.java
 - TrafficLightCtrlStateMachine.java
 - IStateMachine.java
 - ITimer.java
 - ITimerCallback.java
 - RuntimeService.java
 - TimerService.java

- 8 files
- 1311 lines of code
- 302 manual (UI) code

```
TrafficLightCtrl.sct TrafficLightCtrlStateMachine.java
```

```
        break;
    case main_main_trafficlight_interrupted_blinking_Yellow:
        exitSequence_main_main_trafficlight_interrupted_blinking_Yellow();
        break;
    case main_main_trafficlight_normal_normal_Red:
        exitSequence_main_main_trafficlight_normal_normal_Red();
        break;
    case main_main_trafficlight_normal_normal_Yellow:
        exitSequence_main_main_trafficlight_normal_normal_Yellow();
        break;
    case main_main_trafficlight_normal_normal_Green:
        exitSequence_main_main_trafficlight_normal_normal_Green();
        break;
    default:
        break;
    }
}

/* Default exit sequence for region blinking */
private void exitSequence_main_main_trafficlight_interrupted_blinking() {
    switch (stateVector[0]) {
        case main_main_trafficlight_interrupted_blinking_Black:
            exitSequence_main_main_trafficlight_interrupted_blinking_Black();
            break;
        case main_main_trafficlight_interrupted_blinking_Yellow:
            exitSequence_main_main_trafficlight_interrupted_blinking_Yellow();
            break;
        default:
            break;
    }
}

/* Default exit sequence for region normal */
private void exitSequence_main_main_trafficlight_normal_normal() {
    switch (stateVector[0]) {
        case main_main_trafficlight_normal_normal_Red:
            exitSequence_main_main_trafficlight_normal_normal_Red();
            break;
        case main_main_trafficlight_normal_normal_Yellow:
            exitSequence_main_main_trafficlight_normal_normal_Yellow();
            break;
        case main_main_trafficlight_normal_normal_Green:
            exitSequence_main_main_trafficlight_normal_normal_Green();
            break;
        default:
            break;
    }
}

/* Default exit sequence for region timer */
private void exitSequence_main_main_timer() {
    switch (stateVector[1]) {
```

Interface

```
        TrafficLightCtrl
interface:
  in event police_interrupt
  in event toggle

interface TrafficLight:
  out event displayRed
  out event displayGreen
  out event displayYellow
  out event displayNone

interface Timer:
  out event updateTimerColour: string
  out event updateTimerValue: integer

internal:
  event resetTimer
  event disableTimer
  event enableTimer
  var counter: integer
```

Setup Code (Excerpt)

```
protected void setupStatemachine() {
    statemachine = new SynchronizedTrafficLightCtrlStatemachine();
    timer = new MyTimerService(10.0);
    statemachine.setTimer(timer);

    statemachine.getSCITrafficLight().getListeners().add(new ITrafficLightCtrlStatemachine.SCITrafficLightListener() {
        @Override
        public void onDisplayYellowRaised() {
            setLights(false, true, false);
        }

        public void onDisplayRedRaised() {}

        public void onDisplayNoneRaised() {}

        public void onDisplayGreenRaised() {}
    });

    statemachine.getSCITimer().getListeners().add(new ITrafficLightCtrlStatemachine.SCITimerListener() {
        @Override
        public void onUpdateTimerValueRaised(long value) {
            crossing.getCounterVis().setCounterValue(value);
            repaint();
        }

        @Override
        public void onUpdateTimerColourRaised(String value) {
            crossing.getCounterVis().setColor(value == "Red" ? Color.RED : Color.GREEN);
        }
    });

    buttonPanel.getPoliceInterrupt()
        .addActionListener(e -> statemachine.getSCInterface().raisePolice_interrupt());

    buttonPanel.getSwitchOnOff()
        .addActionListener(e -> statemachine.getSCInterface().raiseToggle());

    statemachine.init();
}

private void setLights(boolean red, boolean yellow, boolean green) {
    crossing.getTrafficLightVis().setRed(red);
    crossing.getTrafficLightVis().setYellow(yellow);
    crossing.getTrafficLightVis().setGreen(green);
    repaint();
}

protected void run() {
    statemachine.enter();
    RuntimeService.getInstance().registerStatemachine(statemachine, 100);
}
```

Generator

```
GeneratorModel for yakindu::java {
```

```
    statechart TrafficLightCtrl {

        feature Outlet {
            targetProject = "traffic_light_history"
            targetFolder = "src-gen"
        }

        feature Naming {
            basePackage = "traffic.light"
            implementationSuffix = ""
        }

        feature GeneralFeatures {
            RuntimeService = true
            TimerService = true
            InterfaceObserverSupport = true
        }

        feature SynchronizedWrapper {
            namePrefix = "Synchronized"
            nameSuffix = ""
        }
    }
}
```

Runner

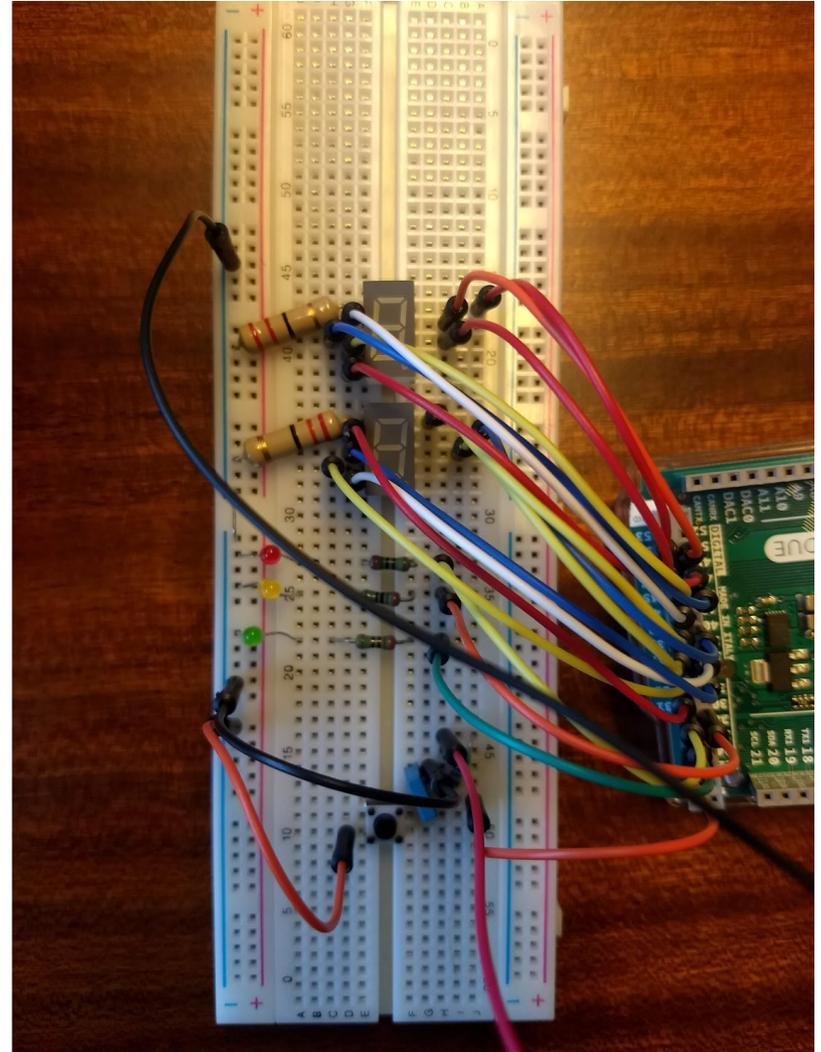
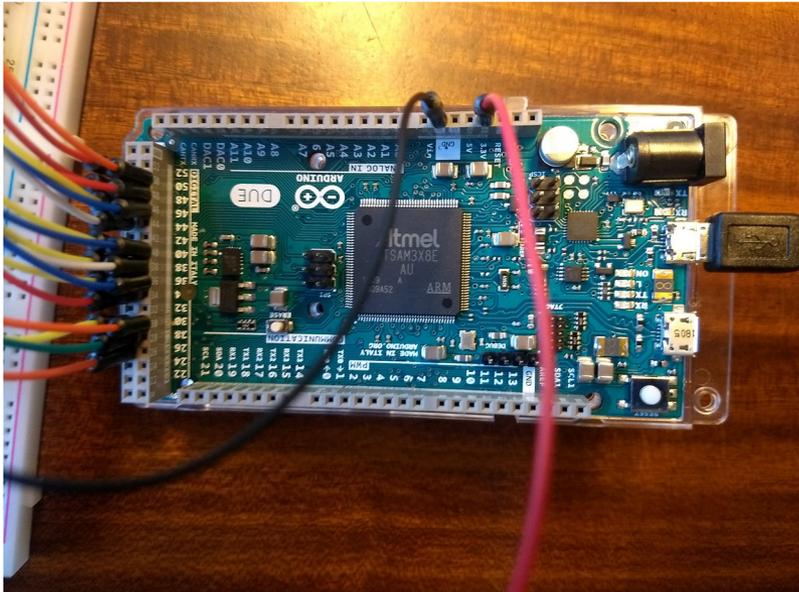
Deployed Application (Scaled Real-Time)



Deploying onto Hardware

Interface:

- `pinMode(pin_nr, mode)`
- `digitalWrite(pin_nr, {0, 1})`
- `digitalRead(pin_nr): {0, 1}`



Deploying onto Hardware

Runner

```
#define CYCLE_PERIOD (10)
static unsigned long cycle_count = 0L;
static unsigned long last_cycle_time = 0L;
```

```
void loop() {
  unsigned long read_pushbut
  if ( cycle_c
    sc_timer_s
  synchroniz
  trafficLig
  last_cycle
  cycle_coun
}
```

Generator

```
GeneratorModel for yakindu::c {
```

```
statechart TrafficLightCtrl {
```

```
feature Outlet {
  targetProject = "traffic_light_arduino"
  targetFolder = "src-gen"
  libraryTargetFolder = "src-gen"
}
```

```
feature FunctionInlining {
  inlineReactions = true
  inlineEntryActions = true
  inlineExitActions = true
  inlineEnterSequences = true
  inlineExitSequences = true
  inlineChoices = true
  inlineEnterRegion = true
  inlineExitRegion = true
  inlineEntries = true
}
```

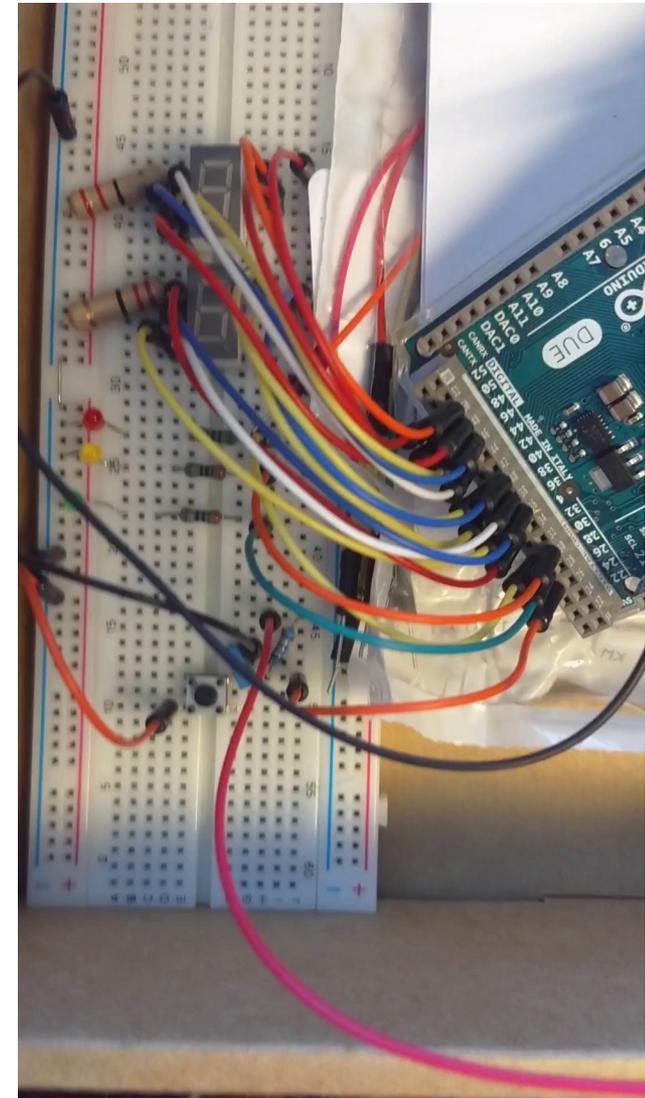
```
_PERIOD) ) {
  e_time);
```

Button Co

```
void read_pushk
int pin_value
if (pin_value
  button->las
}
if ((millis()
  if (pin_val
    button->e
    button->x}
}
}
button->debounce_state = pin_value;
}
```

```
ay) {
```

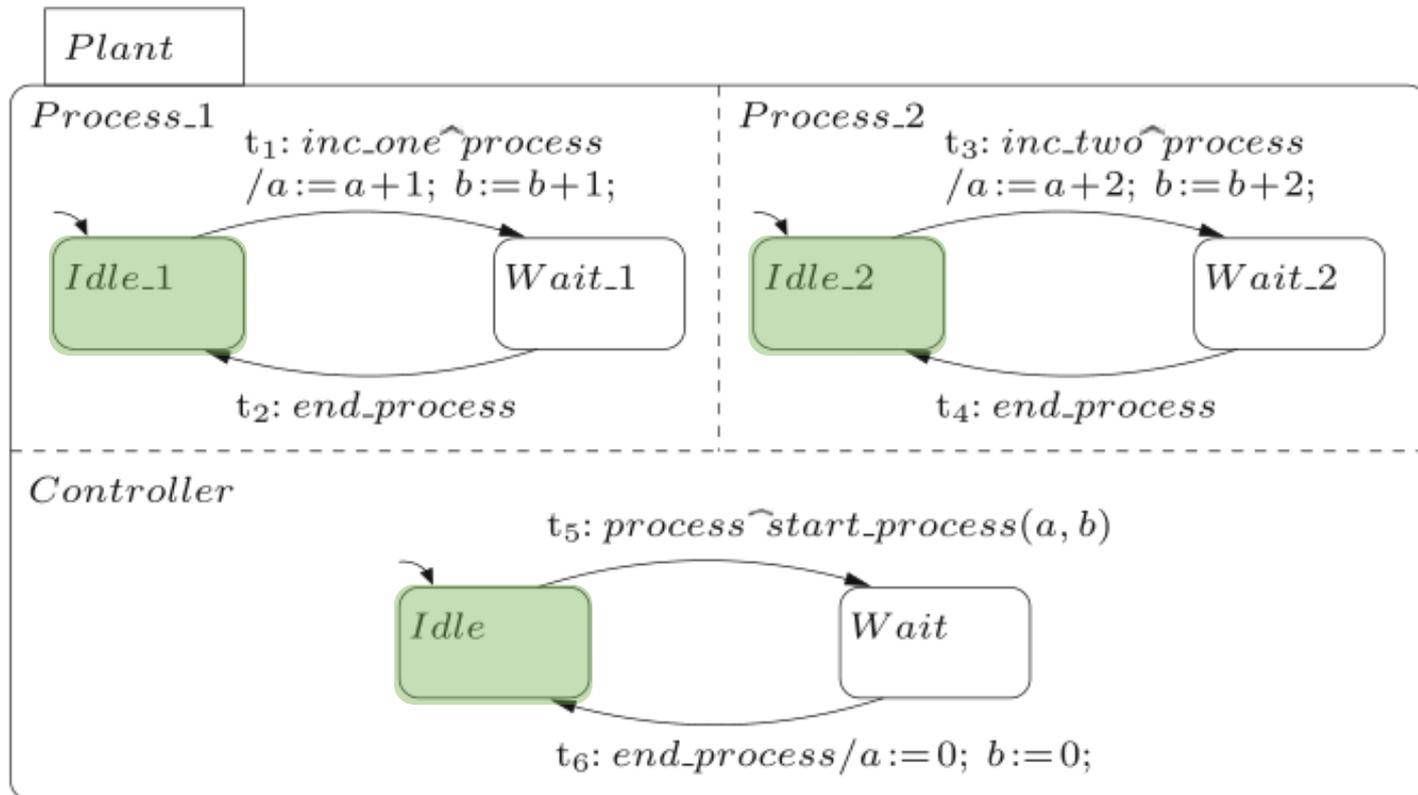
Deployed Application



Semantic Choices

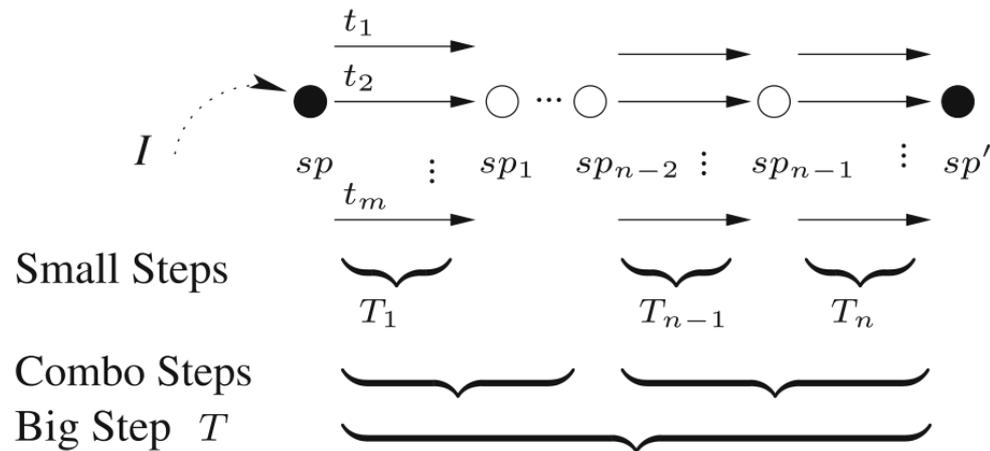
Semantic Choices

enabled events: $[inc_one, inc_two]$

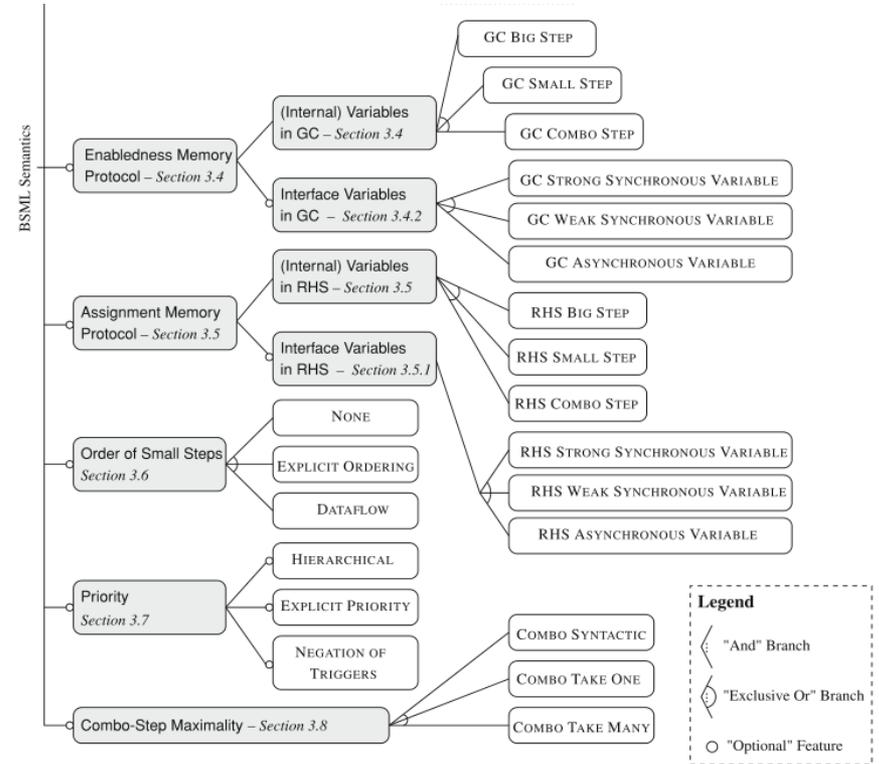
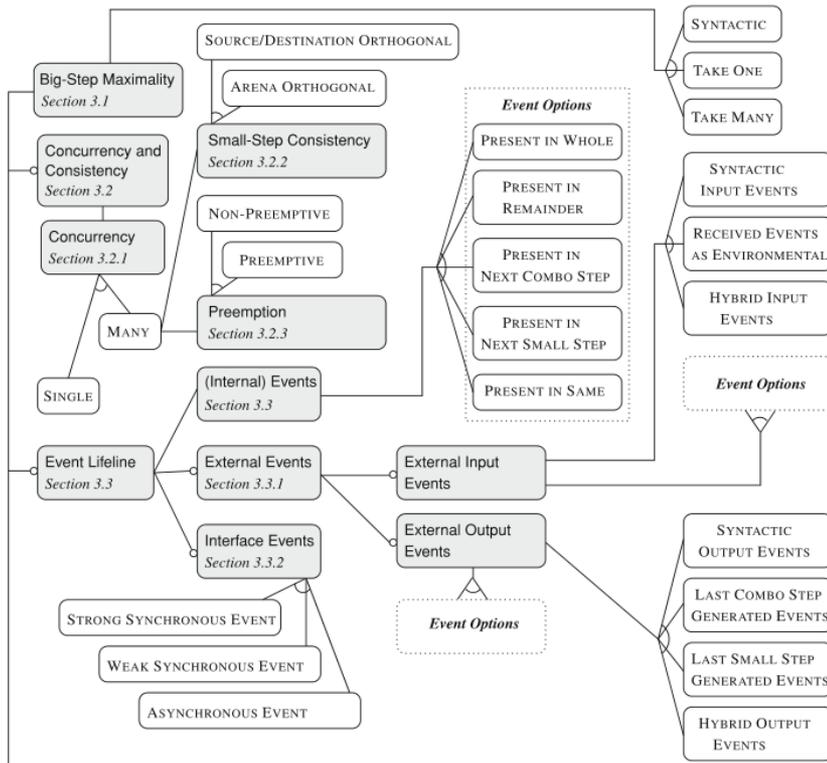


Big Step, Small Step

- A “big step” takes the system from one “quiescent state” to the next.
- A “small step” takes the system from one “snapshot” to the next (execution of a set of enabled transitions).
- A “combo step” groups multiple small steps.

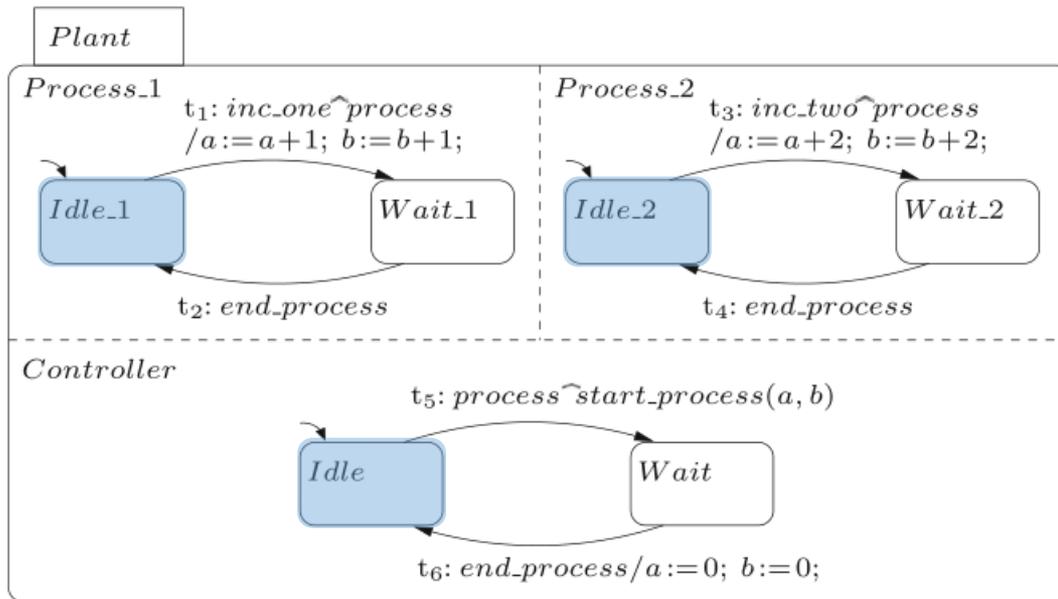


Semantic Options



Revisiting the Example

enabled events: $[inc_one, inc_two]$



concurrency: single

event lifeline: next combo step

assignment: RHS small step

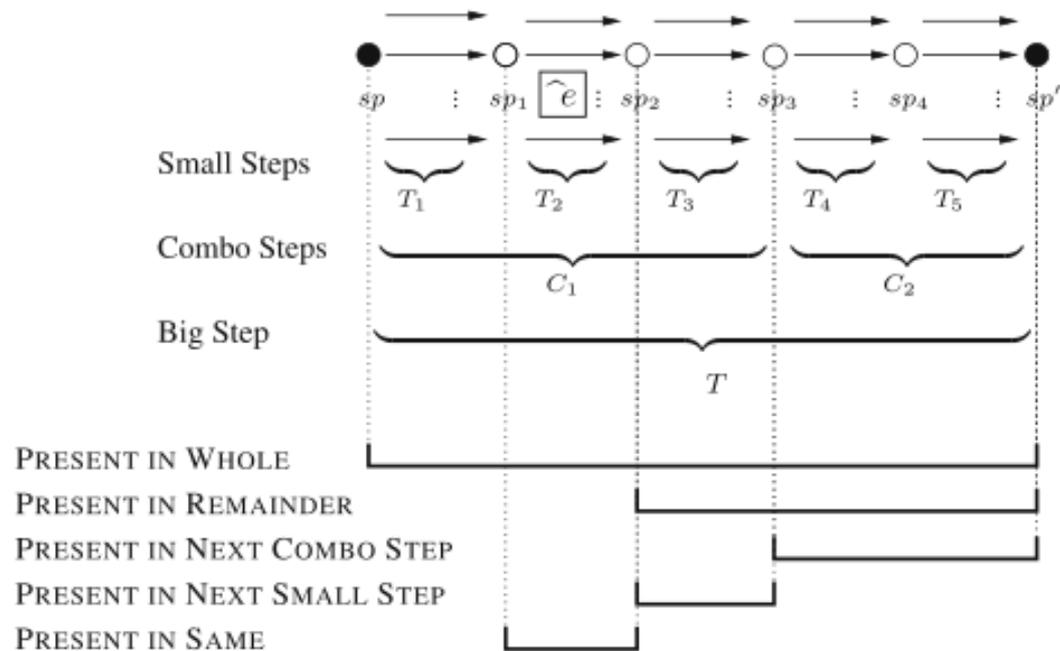
-> $\langle \{t1\}, \{t3\}, \{t5\} \rangle$ and

$\langle \{t3\}, \{t1\}, \{t5\} \rangle$

event lifeline: present in remainder

-> $\langle \{t1\}, \{t5\}, \{t3\} \rangle$ becomes possible

Event Lifeline



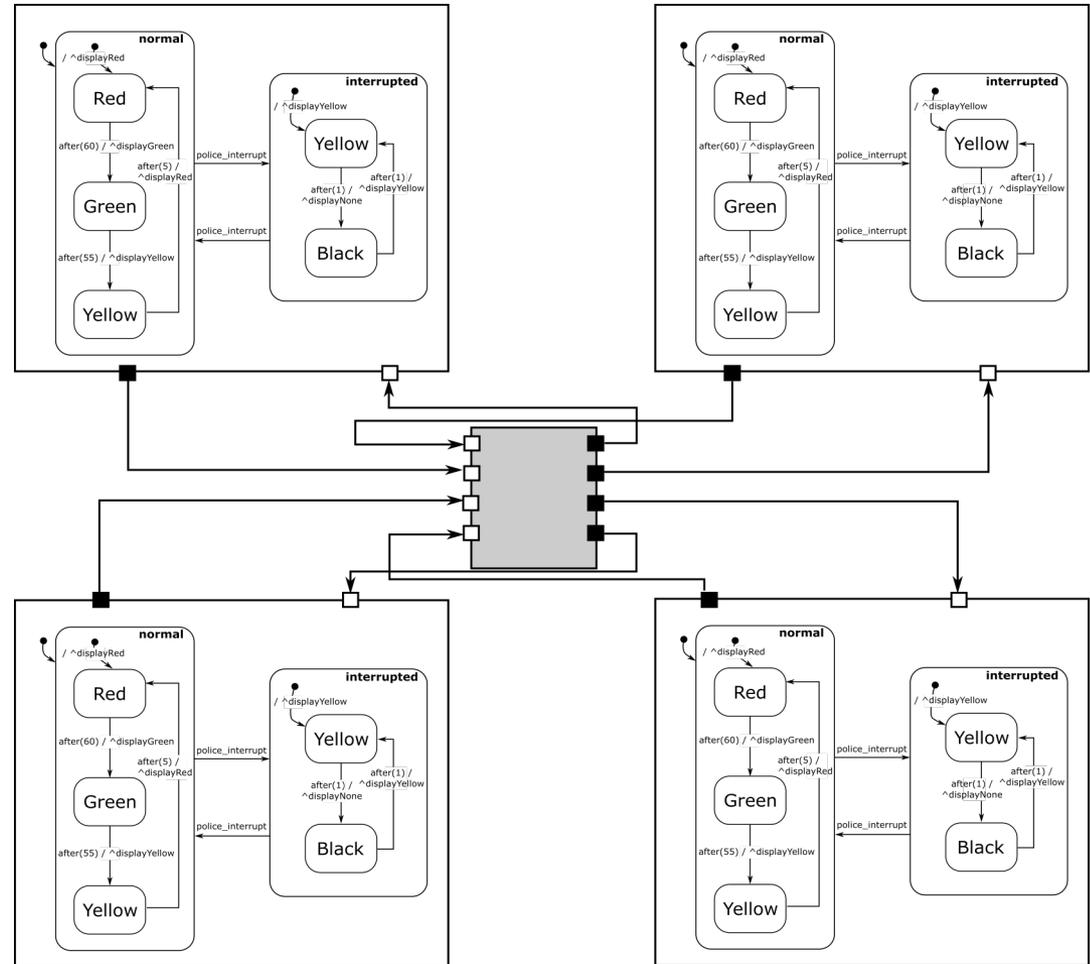
Semantic Options: Examples

	Rhapsody	Statemate	(Default) SCCD
Big Step Maximality	Take Many	Take Many	Take Many
Internal Event Lifeline	Queue	Next Combo Step	Queue
Input Event Lifeline	First Combo Step	First Combo Step	First Combo Step
Priority	Source-Child	Source-Parent	Source-Parent
Concurrency	Single	Single	Single

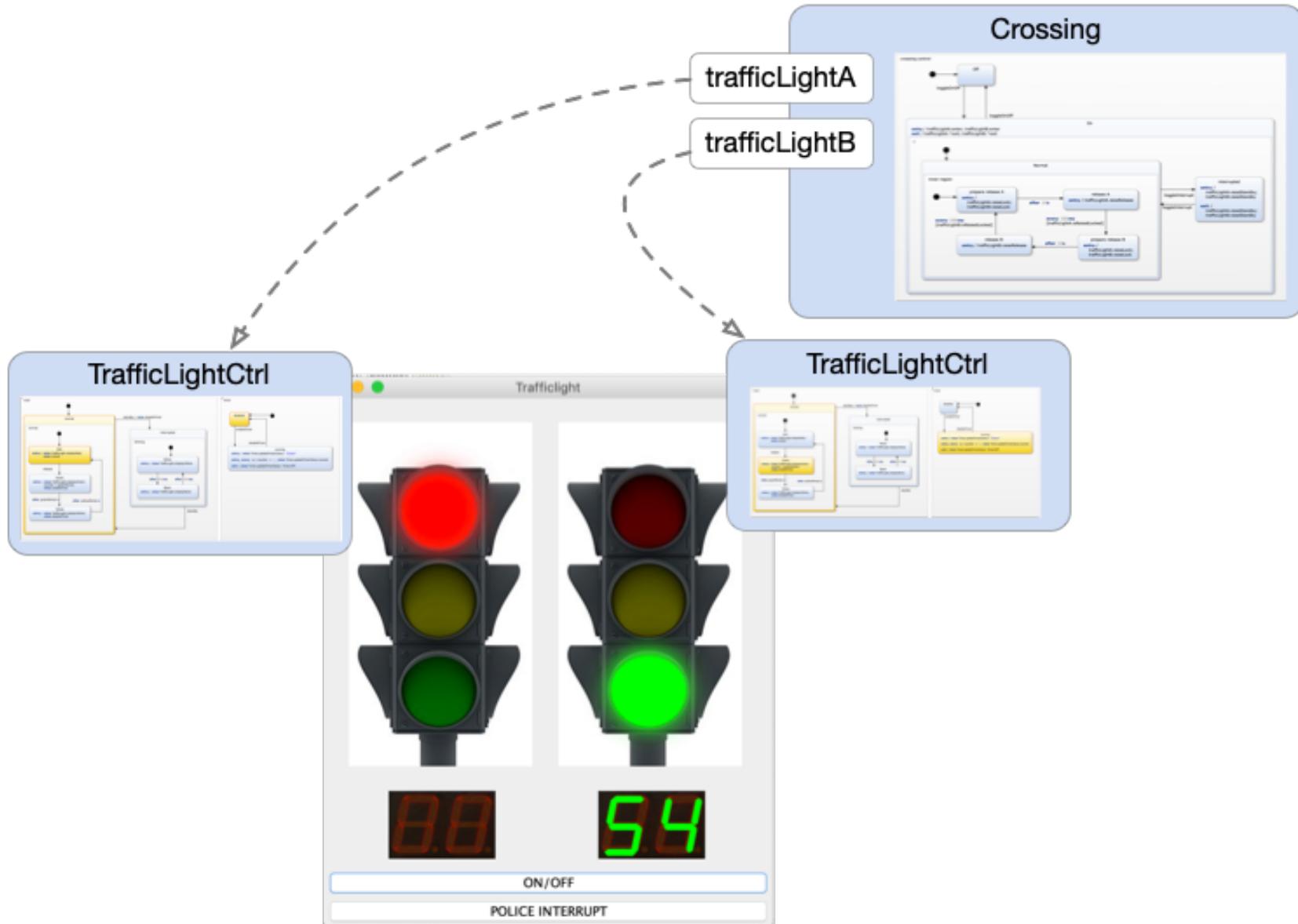
Composition

Composition of Statecharts

- Composition of multiple Statechart models
 - Instantiation
 - Communication
 - Semantics
- Often solved in code...

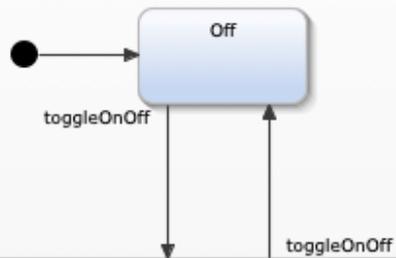


Composition Example



Composition Example

crossing control



On

entry / trafficLightA.enter; trafficLightB.enter
exit / trafficLightA.^exit; trafficLightB.^exit

r

Normal

inner region



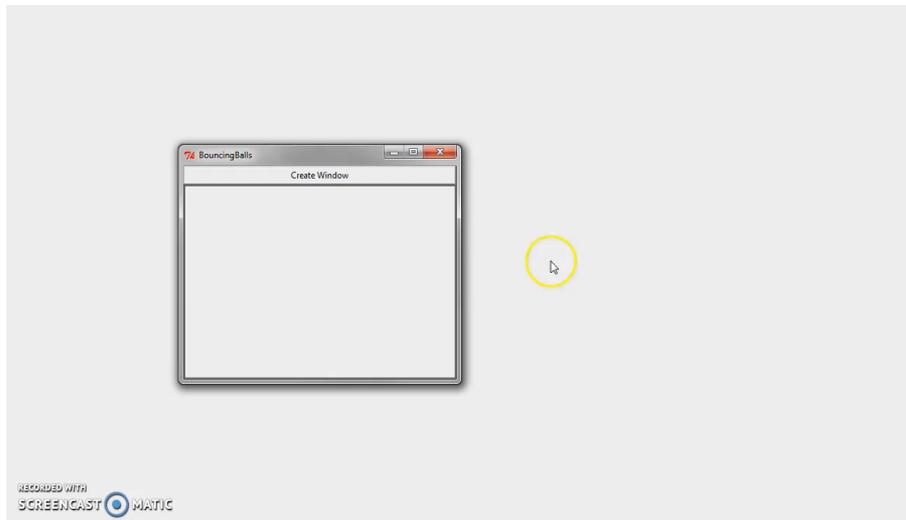
interrupted

entry /
trafficLightA.raiseStandby;
trafficLightB.raiseStandby
exit /
trafficLightA.raiseStandby;
trafficLightB.raiseStandby

toggleInterrupt

toggleInterrupt

Dynamic Structure: SCCD



Behavior

- Timed
- Autonomous
- Interactive
- Hierarchical

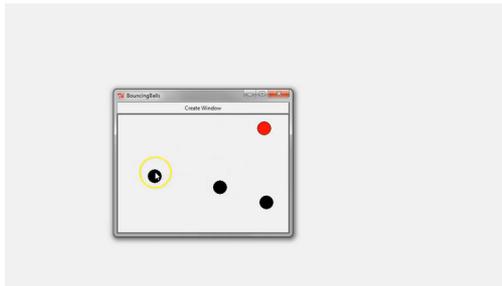
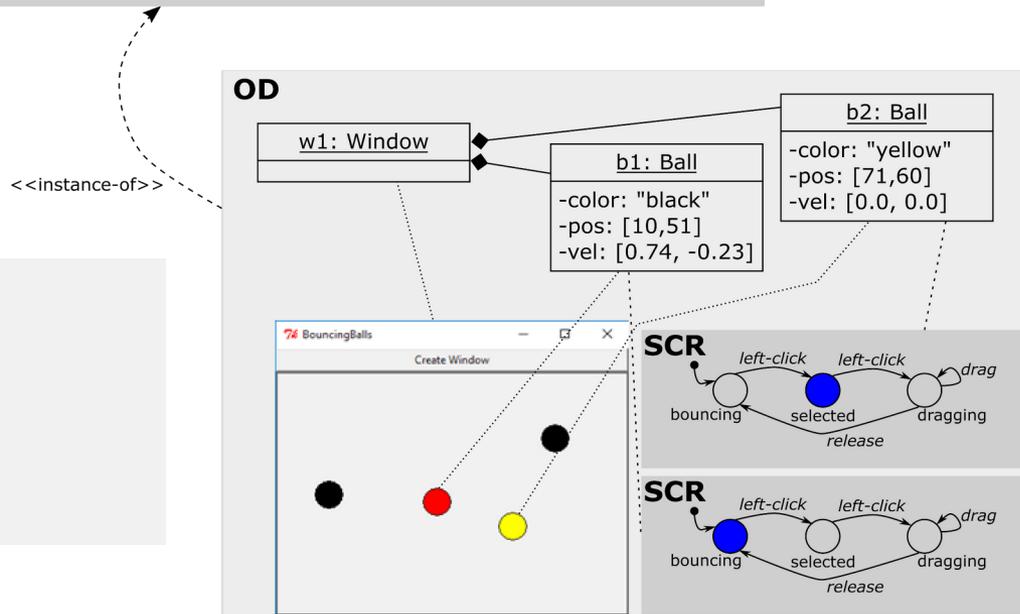
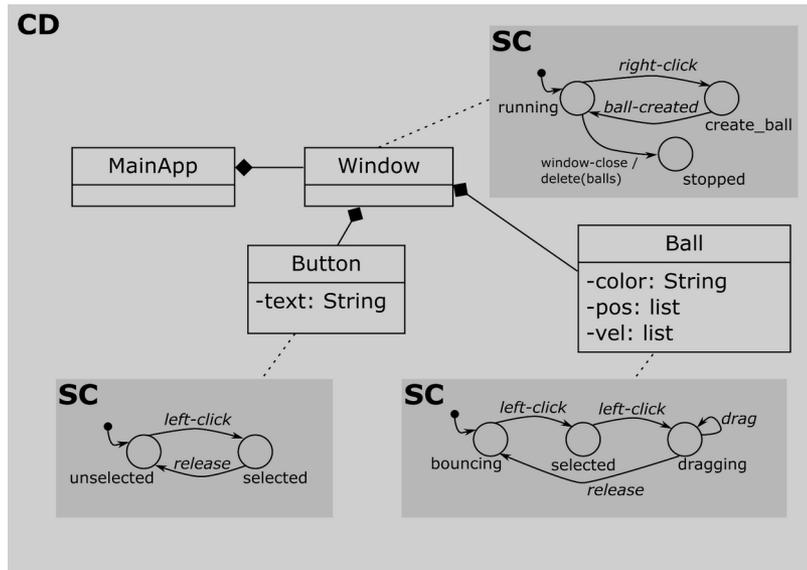
Structure

- Dynamic
- Hierarchical

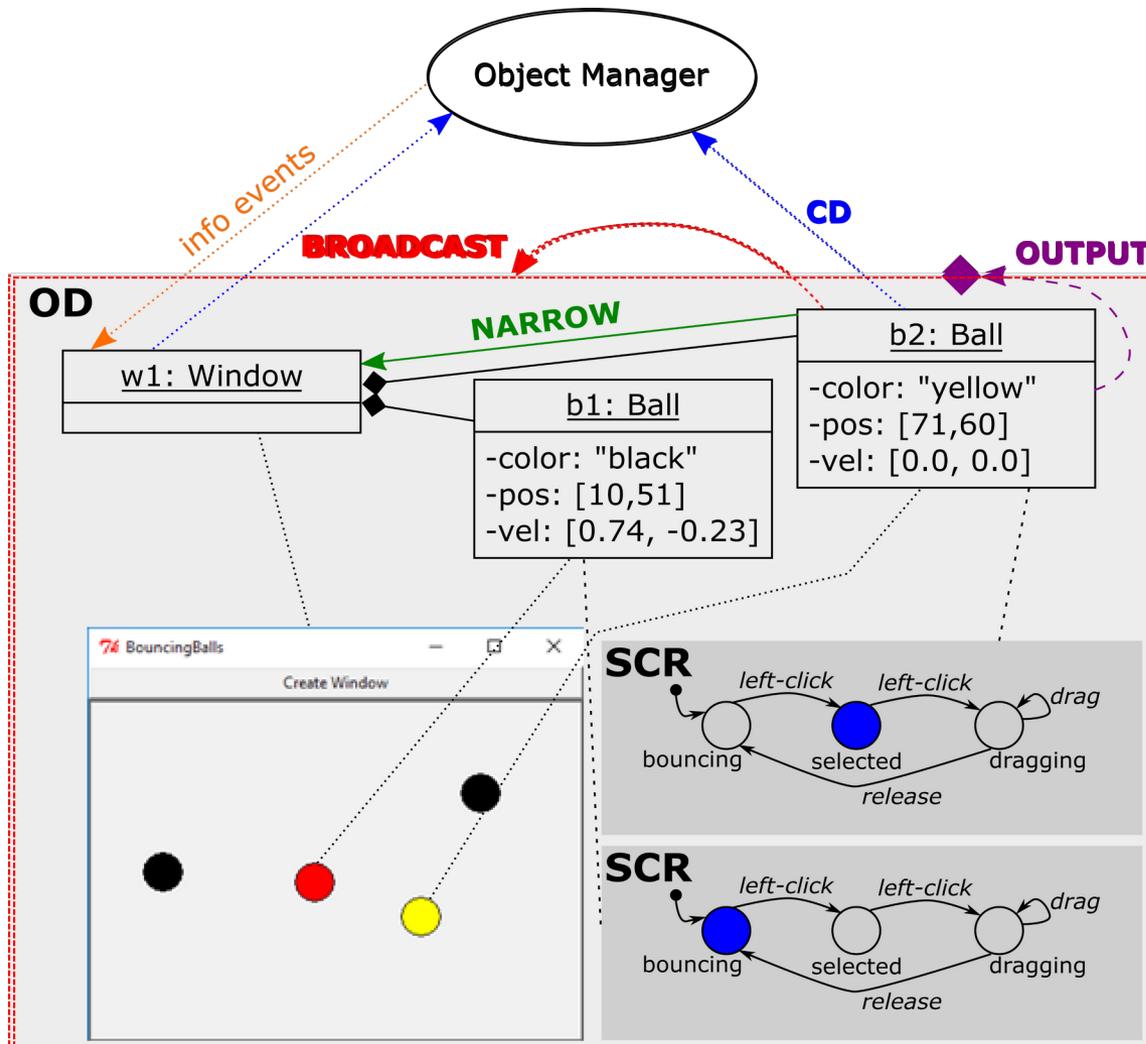
Design? Statecharts + ??? Coordination/Communication/Dynamic Structure often implemented in code...

Simon Van Mierlo, Yentl Van Tendeloo, Bart Meyers, Joeri Exelmans, and Hans Vangheluwe. SCCD: SCXML extended with class diagrams. In 3rd Workshop on Engineering Interactive Systems with SCXML, part of EICS 2016, 2016.

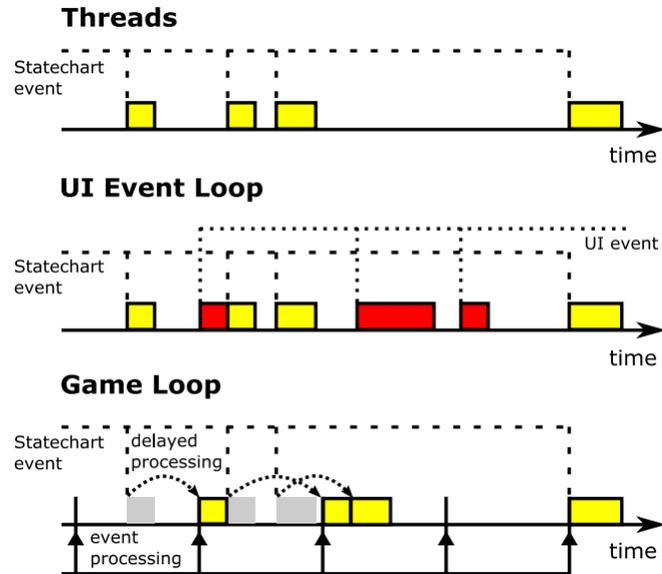
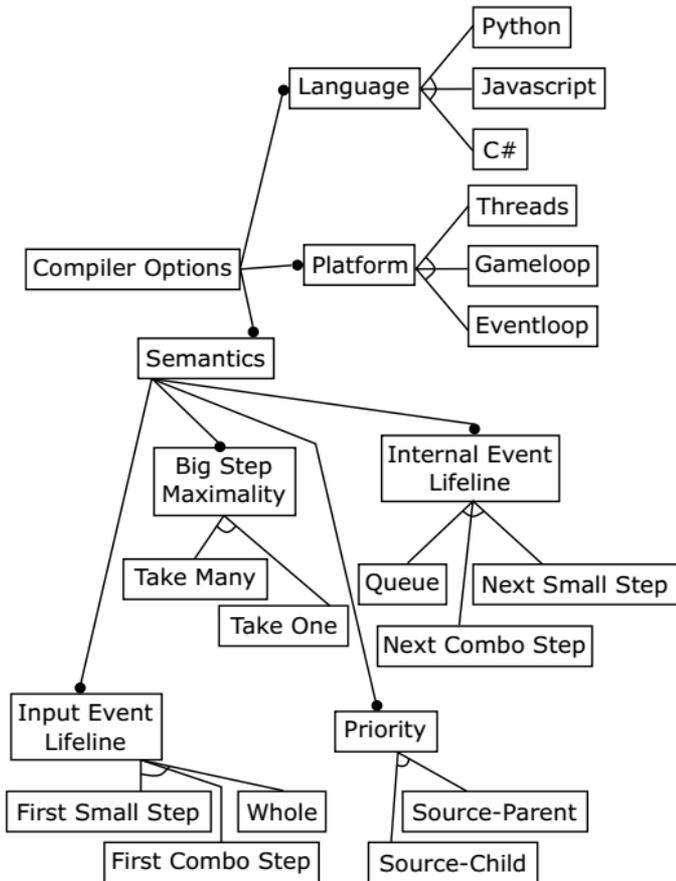
SCCD: Conformance



Communication: Event Scopes



SCCD Compiler



Simon Van Mierlo, Yentl Van Tendeloo, Bart Meyers, Joeri Exelmans, and Hans Vangheluwe. **SCCD: SCXML extended with class diagrams**. In *3rd Workshop on Engineering Interactive Systems with SCXML, part of EICS 2016*, 2016

<https://msdl.uantwerpen.be/documentation/SCCD/>

SCCD Documentation

SCCD [SCCD] is a language that combines the Statecharts [Statecharts] language with Class Diagrams. It allows users to model complex, timed, autonomous, reactive, dynamic-structure systems.

The concrete syntax of SCCD is an XML-format loosely based on the W3C SCXML recommendation. A conforming model can be compiled to a number of programming languages, as well as a number of runtime platforms implemented in those languages. This maximizes the number of applications that can be modelled using SCCD, such as user interfaces, the artificial intelligence of game characters, controller software, and much more.

This documentation serves as an introduction to the SCCD language, its compiler, and the different supported runtime platforms.

Contents

- Installation
 - Download
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 - SCCD Installation
- Language Features
 - Top-Level Elements
 - Class Diagram Concepts
 - Statechart Concepts
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- Compiler
- Runtime Platforms
 - Threads
 - Eventloop
 - Gameloop
- Examples
 - Timer
 - Traffic Lights
- Semantic Options
 - Big Step Maximality
 - Internal Event Lifeline
 - Input Event Lifeline
 - Priority
 - Concurrency
- Socket Communication
 - Initialization
 - Input Events
 - Output Events
 - HTTP client/server
- Internal Documentation
 - Statecharts Core

References

- [SCCD] Simon Van Mierlo, Yentl Van Tendeloo, Bart Meyers, Joeri Exelmans, and Hans Vangheluwe. SCCD: SCXML extended with class diagrams. In *3rd Workshop on Engineering Interactive Systems with SCXML, part of EICS 2016*, 2016. [LINK]
- [Statecharts] David Harel. Statecharts: A visual formalism for complex systems. *Sci. Comput. Program.* 8, 3 (1987), 231–274. [LINK]

Recap

- Model the behaviour of complex, timed, reactive, autonomous systems
 - “What” instead of “How”
(= implemented by Statecharts compiler)
- Abstractions:
 - States (composite, orthogonal)
 - Transitions
 - Timeouts
 - Events
- Tool support:
 - Yakindu
 - SCCD